



# **IS25RLP256E      IS25RWP256E**

**256Mb**

**SERIAL FLASH MEMORY 166MHZ MULTI I/O SPI &  
QUAD I/O QPI DTR INTERFACE WITH RPMC**

**PRELIMINARY DATA SHEET**

## 256Mb

**SERIAL FLASH MEMORY 166MHZ MULTI I/O SPI &  
QUAD I/O QPI DTR INTERFACE WITH RPMC**

PRELIMINARY INFORMATION

## FEATURES

- **Industry Standard Serial Interface**
  - IS25RLP256E: 256Mbit/32Mbyte
  - IS25RWP256E: 256Mbit/32Mbyte
  - 3 or 4 Byte Addressing Mode
  - Supports Standard SPI, Fast, Dual, Dual I/O, Quad, Quad I/O, SPI DTR, Dual I/O DTR, Quad I/O DTR, and QPI
  - QE=1 from the factory
  - Software & Hardware Reset
  - Supports Serial Flash Discoverable Parameters (SFDP)
- **High Performance Serial Flash (SPI)**
  - 80MHz Normal Read<sup>(3)</sup>
  - Up to 166MHz Fast Read
  - Up to 80MHz DTR (Dual Transfer Rate)
  - Equivalent Throughput of 664 Mb/s
  - Up to 104MHz RPMC Operation
  - Selectable Dummy Cycles
  - Configurable Drive Strength
  - Supports SPI Modes 0 and 3
  - More than 100,000 Erase/Program Cycles
  - More than 20-year Data Retention
- **Flexible & Efficient Memory Architecture**
  - Chip Erase with Uniform Sector/Block Erase (4/32/64KB or 4/32/256 KB)<sup>(2)</sup>
  - Program 1 to 256 or 512 Byte per Page<sup>(2)</sup>
  - Program/Erase Suspend & Resume
- **Efficient Read and Program modes**
  - Low Instruction Overhead Operations
  - Continuous Read 8/16/32/64 Byte Burst Wrap
  - Selectable Burst Length
  - QPI for Reduced Instruction Overhead
  - AutoBoot Operation
  - Data Learning Pattern for training in DTR operation
- **Low Power with Wide Temp. Ranges**
  - Single Voltage Supply
  - IS25RLP: 2.30V to 3.60V
  - IS25RWP: 1.70V to 1.95V
  - 7 mA Active Read Current
  - Temp Grades:
  - Industrial: -40°C to +85°C
- **Advanced Security Protection**
  - Software and Hardware Write Protection
  - Advanced Sector/Block Protection
  - Top/Bottom Block Protection
  - Power Supply Lock Protection
  - 4x256 Byte Dedicated Security Area with OTP User-lockable Bits
  - 128 bit Unique ID for Each Device (Call Factory)
- **RPMC Features**
  - Support Reply Protection Monotonic Counter
  - Four 32-bit Monotonic counters
  - Volatile HMAC Key register
  - Non-volatile Root Key register
- **Industry Standard Pin-out & Packages<sup>(1)</sup>**
  - M = 16-pin SOIC 300mil
  - J = 8-contact WSON 8x6mm

Notes:

1. Call Factory for other package options available.
2. For optional 512 Byte Page size with 256 KB Block size, see the Ordering Information.
3. 80MHz for 3.0V device and 50MHz for 1.8V device.

## GENERAL DESCRIPTION

The IS25RLP256E and IS25RWP256E Serial Flash memory offers a versatile storage solution with high flexibility and performance in a simplified pin count package. ISSI's "Industry Standard Serial Interface" Flash is for systems that require limited space, a low pin count, and low power consumption. The device is accessed through a 4-wire SPI Interface consisting of a Serial Data Input (SI), Serial Data Output (SO), Serial Clock (SCK), and Chip Enable (CE#) pins, which can also be configured to serve as multi-I/O (see pin descriptions).

The device supports Dual and Quad I/O as well as standard, Dual Output, and Quad Output SPI. Clock frequencies of up to 166MHz allow for equivalent clock rates of up to 664MHz (166MHz x 4) which equates to 83Mbytes/s of data throughput. The IS25xE series of Flash adds support for DTR (Double Transfer Rate) commands that transfer addresses and read data on both edges of the clock. These transfer rates can outperform 16-bit Parallel Flash memories allowing for efficient memory access to support XIP (execute in place) operation.

Initial state of the memory array is erased (all bits are set to 1) when shipped from the factory.

QPI (Quad Peripheral Interface) supports 2-cycle instruction further reducing instruction times. Pages can be erased in groups of 4Kbyte sectors, 32Kbyte blocks, 64K/256Kbyte blocks, and/or the entire chip. The uniform sector and block architecture allows for a high degree of flexibility so that the device can be utilized for a broad variety of applications requiring solid data retention.

The device supports RPMC (Replay Protection Monotonic Counter) function as an additional security, which could detect Hardware attack.

### Supply Voltage Range vs. Maximum Speed

		Voltage	Speed
IS25RLP (3.0V typ.)	SPI NOR protocol	2.30~3.6V	133MHz
		2.70~3.6V	166MHz
	RPMC protocol	2.30~3.6V	104MHz
IS25RWP (1.8V typ.)	SPI NOR protocol	1.70~1.95V	166MHz
	RPMC protocol	1.70~1.95V	104MHz

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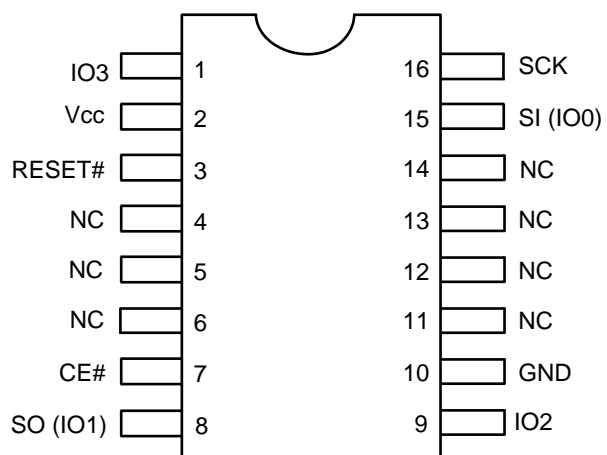
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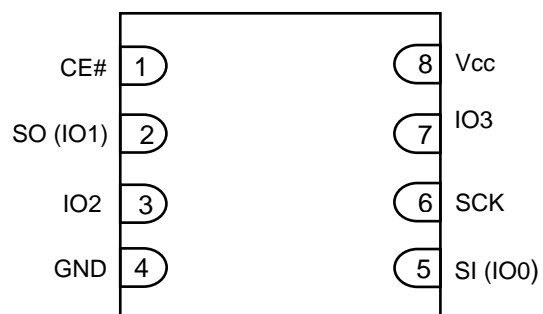
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## 1. PIN CONFIGURATION



**16-pin SOIC 300mil**



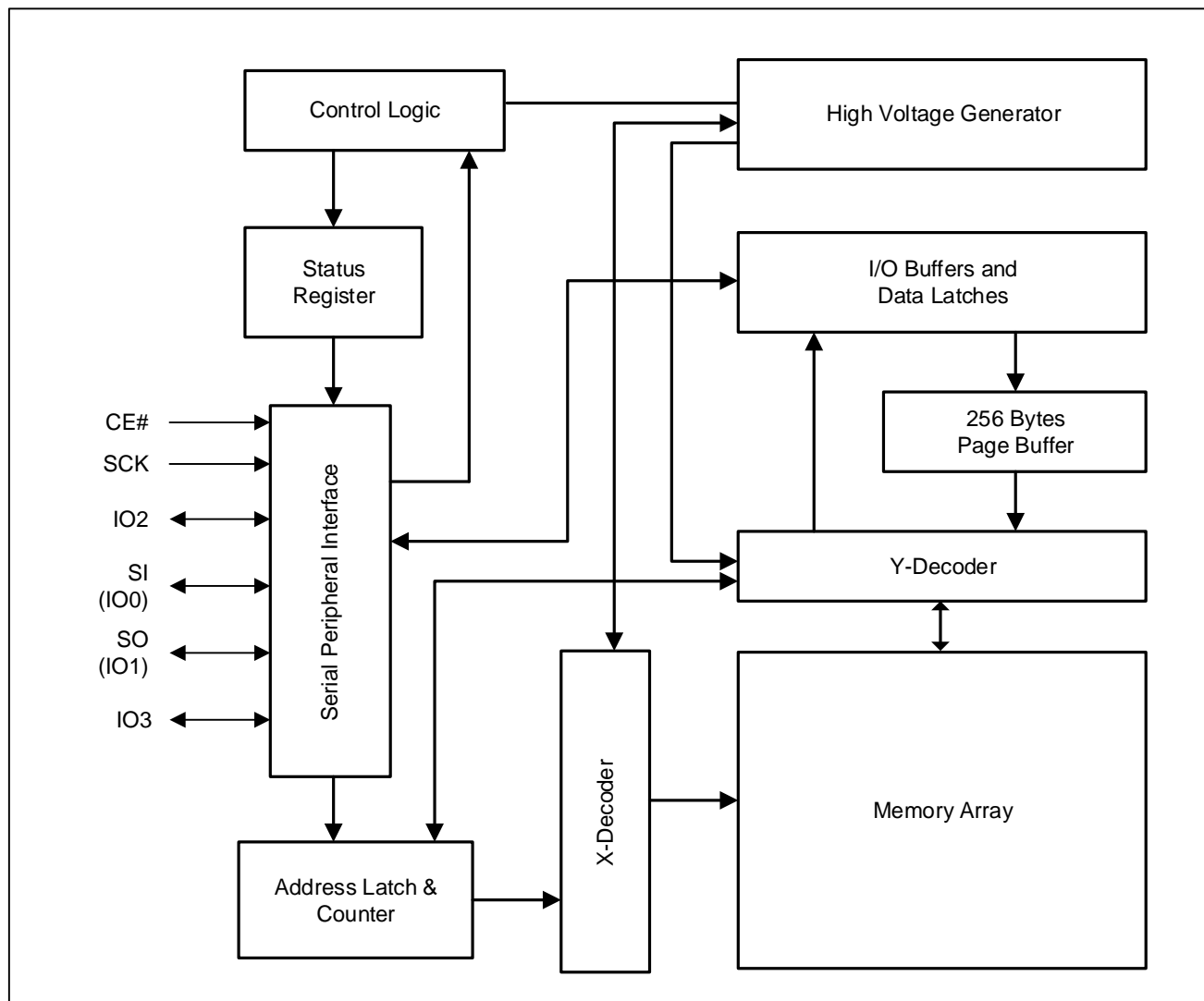
**8-contact WSON 8x6mm**



## 2. PIN DESCRIPTIONS

SYMBOL	TYPE	DESCRIPTION
CE#	INPUT	<p><b>Chip Enable:</b> The Chip Enable (CE#) pin enables and disables the devices operation. When CE# is high the device is deselected and output pins are in a high impedance state. When deselected the devices non-critical internal circuitry power down to allow minimal levels of power consumption while in a standby state.</p> <p>When CE# is pulled low the device will be selected and brought out of standby mode. The device is considered active and instructions can be written to, data read, and written to the device. After power-up, CE# must transition from high to low before a new instruction will be accepted.</p> <p>Keeping CE# in a high state deselects the device and switches it into its low power state. Data will not be accepted when CE# is high.</p>
SI (IO0), SO (IO1)	INPUT/OUTPUT	<p><b>Serial Data Input, Serial Output, and IOs (SI, SO, IO0, and IO1):</b></p> <p>This device supports standard SPI, Dual SPI, and Quad SPI operation. Standard SPI instructions use the unidirectional SI (Serial Input) pin to write instructions, addresses, or data to the device on the rising edge of the Serial Clock (SCK). Standard SPI also uses the unidirectional SO (Serial Output) to read data or status from the device on the falling edge of the serial clock (SCK).</p> <p>In Dual and Quad SPI mode, SI and SO become bidirectional IO pins to write instructions, addresses or data to the device on the rising edge of the Serial Clock (SCK) and read data or status from the device on the falling edge of SCK. Quad SPI instructions use the WP# and HOLD# pins as IO2 and IO3 respectively.</p>
IO2, IO3	INPUT/OUTPUT	<b>Serial Data IO :</b>
RESET#	INPUT	<p><b>Dedicated RESET#:</b> This dedicated RESET# is available in 16-pin SOIC package.</p> <p>The RESET# pin is a hardware RESET signal. When RESET# is driven HIGH, the memory is in the normal operating mode. When RESET# is driven LOW, the memory enters reset mode and output is High-Z. If RESET# is driven LOW while an internal WRITE, PROGRAM, or ERASE operation is in progress, data may be lost.</p> <p>Dedicated RESET# function can be Disabled when bit 0 of Function Register = 1. It has an internal pull-up resistor and may be left floating if not used.</p>
SCK	INPUT	<b>Serial Data Clock:</b> Synchronized Clock for input and output timing operations.
Vcc	POWER	<b>Power:</b> Device Core Power Supply
GND	GROUND	<b>Ground:</b> Connect to ground when referenced to Vcc
NC	Unused	<b>NC:</b> Pins labeled "NC" stand for "No Connect" and should be left uncommitted.

### 3. BLOCK DIAGRAM



#### 4. SPI MODES DESCRIPTION

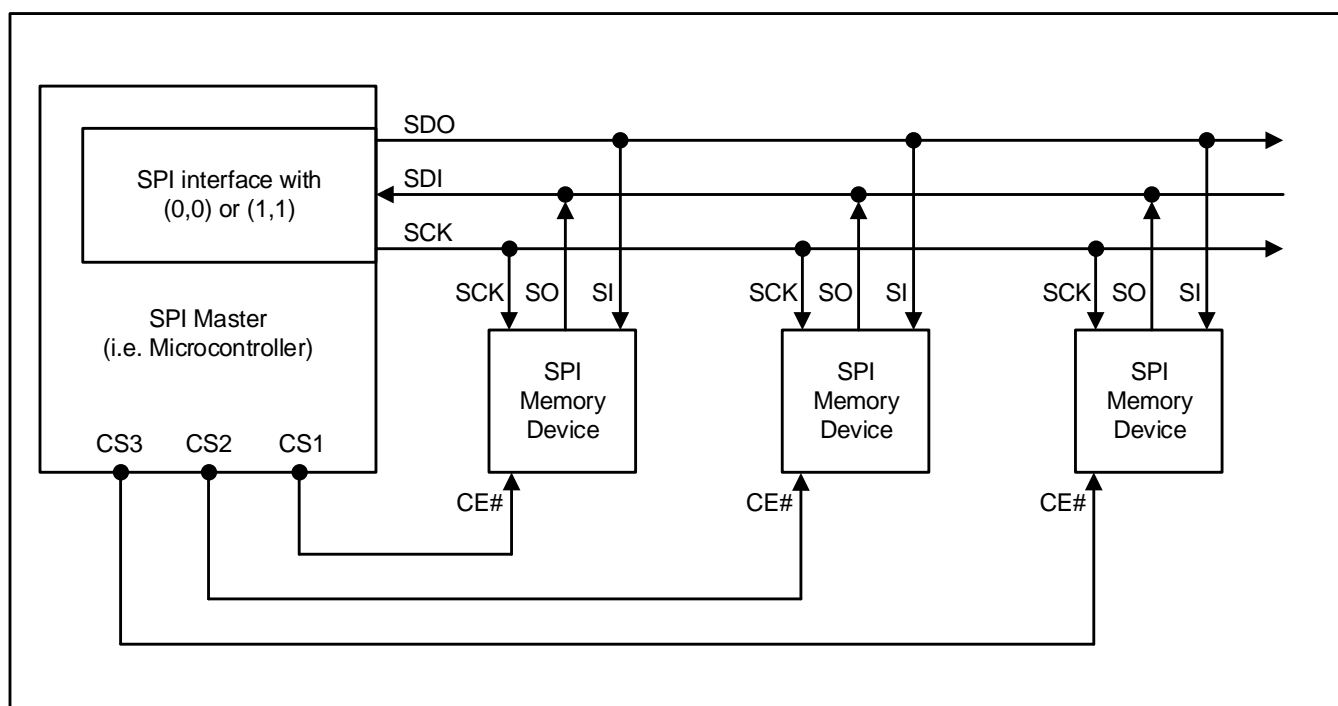
Multiple IS25RLP256E devices or multiple IS25RWP256E devices can be connected on the SPI serial bus and controlled by a SPI Master, i.e. microcontroller, as shown in Figure 4.1. The devices support either of two SPI modes:

Mode 0 (0, 0)

Mode 3 (1, 1)

The difference between these two modes is the clock polarity. When the SPI master is in stand-by mode, the serial clock remains at “0” (SCK = 0) for Mode 0 and the clock remains at “1” (SCK = 1) for Mode 3. Please refer to Figure 4.2 and Figure 4.3 for SPI and QPI mode. In both modes, the input data is latched on the rising edge of Serial Clock (SCK), and the output data is available from the falling edge of SCK.

**Figure 4.1 Connection Diagram among SPI Master and SPI Slaves (Memory Devices)**



**Note:**

1. SI and SO pins become bidirectional IO0 and IO1 respectively during Dual I/O mode and SI, SO.

The diagram illustrates the timing for SPI Mode 0 (0,0). It consists of three horizontal signal traces:

- SCK (Serial Clock):** A square wave signal. The label "SCK" is at the start, and "Mode 0 (0,0)" is centered below the trace.
- SI (Serial Input):** A signal trace that is high during the first two clock cycles, then transitions to low for the next two cycles, and continues with a dashed line indicating further cycles. The label "SI" is at the start.
- SO (Serial Output):** A signal trace that is low throughout the first four clock cycles, then transitions to high for the next two cycles, and continues with a dashed line. The label "SO" is at the start.

Hexagonal markers labeled "MSB" (Most Significant Bit) are placed on the SI and SO traces to indicate the timing of the first data bit transfer. The SI MSB occurs during the first clock cycle, and the SO MSB occurs during the fifth clock cycle.

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## 5. SYSTEM CONFIGURATION

The memory array is divided into uniform 4 Kbyte sectors or uniform 32/64 Kbyte blocks (a block consists of eight/sixteen adjacent sectors respectively).

Also optional uniform 32/256 Kbyte blocks are available with 512B page size. See ordering table for more detail information.

Table 5.1 illustrates the memory map of the device. The Status Register controls how the memory is protected.

### 5.1 BLOCK/SECTOR ADDRESSES

**Table 5.1 Block/Sector Addresses (Block Size = 32KB/64KB)**

Memory Density	Block No. (64Kbyte)	Block No. (32Kbyte)	Sector No.	Sector Size (Kbyte)	Address Range
256Mb	Block 0	Block 0	Sector 0	4	000000h - 000FFFh
			:	:	:
		Block 1	:	:	:
			Sector 15	4	00F000h - 00FFFFh
	Block 1	Block 2	Sector 16	4	010000h - 010FFFh
			:	:	:
		Block 3	:	:	:
			Sector 31	4	01F000h - 01FFFFh
	Block 2	Block 4	Sector 32	4	020000h - 020FFFh
			:	:	:
		Block 5	:	:	:
			Sector 47	4	02F000h - 02FFFFh
	:	:	:	:	:
	Block 254	Block 508	Sector 4064	4	FE0000h - FE0FFFh
			:	:	:
		Block 509	:	:	:
			Sector 4079	4	FEF000h - FFFFFFFh
	Block 255	Block 510	Sector 4080	4	FF0000h - FF0FFFh
			:	:	:
		Block 511	:	:	:
			Sector 4095	4	FFF000h - FFFFFFFh
	:	:	:	:	:
	Block 510	Block 1020	Sector 8160	4	1FE0000h - 1FE0FFFh
			:	:	:
		Block 1021	:	:	:
			Sector 8175	4	1FEF000h - 1FEFFFFh
	Block 511	Block 1022	Sector 8176	4	1FF0000h - 1FF0FFFh
			:	:	:
		Block 1023	:	:	:
			Sector 8191	4	1FFF000h - 1FFFFFFh

**Table 5.2 Block/Sector Addresses (Block Size = 32KB/256KB)**

Memory Density	Block No. (256Kbyte)	Block No. (32Kbyte)	Sector No.	Sector Size (Kbyte)	Address Range
256Mb	Block 0	Block 0	Sector 0	4	000000h - 000FFFh
			:	:	:
		Block 1	:	:	:
			Sector 15	4	00F000h - 00FFFFh
		Block 2	Sector 16	4	010000h - 010FFFh
			:	:	:
		Block 3	:	:	:
			Sector 31	4	01F000h - 01FFFFh
		Block 4	Sector 32	4	020000h - 020FFFh
			:	:	:
		Block 5	:	:	:
			Sector 47	4	02F000h - 02FFFFh
		Block 6	Sector 48		030000h - 030FFFh
					:
		Block 7			:
			Sector 63		03F000h - 03FFFFh
	:	:	:	:	:
	Block 63	Block 504	Sector 4032	4	FC0000h - FC0FFFh
			:	:	:
		Block 505	:	:	:
			Sector 4047	4	FCF000h - FCFFFFh
		:	:	:	:
		Block 510	Sector 4080	4	FF0000h - FF0FFFh
			:	:	:
		Block 511	:	:	:
			Sector 4095	4	FFF000h - FFFFFFh
	:	:	:	:	:
	Block 127	Block 1016	Sector 8128	4	1FC0000h - 1FC0FFFh
			:	:	:
		Block 1017	:	:	:
			Sector 8143	4	1FCF000h - 1FCFFFFh
		Block 1018	Sector 8144		1FD0000h - 1FD0FFFh
			:		:
		Block 1019	:		:
			Sector 8159		1FDF000h - 1FDFFFFh
		Block 1020	Sector 8160	4	1FE0000h - 1FE0FFFh
			:	:	:
		Block 1021	:	:	:
			Sector 8175	4	1FEF000h - 1FEFFFFh
		Block 1022	Sector 8176	4	1FF0000h - 1FF0FFFh
			:	:	:
		Block 1023	:	:	:
			Sector 8191	4	1FFF000h - 1FFFFFFh

## 5.2 SERIAL FLASH DISCOVERABLE PARAMETERS

The Serial Flash Discoverable Parameters (SFDP) standard defines the structure of the SFDP database within the memory device. SFDP is the standard of JEDEC JESD216.

The JEDEC-defined header with Parameter ID FF00h and related Basic Parameter Table is mandatory. Additional parameter headers and tables are optional.

**Table 5.3 Signature and Parameter Identification Data Values**

Description		Address (Byte)	Address (Bit)	Data
SFDP Signature		00h	7:0	53h
		01h	15:8	46h
		02h	23:16	44h
		03h	31:24	50h
SFDP Revision	Minor	04h	7:0	06h
	Major	05h	15:8	01h
Number of Parameter Headers (NPH)		06h	23:16	01h
Unused		07h	31:24	FFh
Parameter ID LSB		08h	7:0	00h
Parameter Minor Revision		09h	15:8	06h
Parameter Major Revision		0Ah	23:16	01h
Parameter Table Length (in DWPRDs)		0Bh	31:24	10h
Basic Flash Parameter Table Pointer (PTP)		0Ch	7:0	30h
		0Dh	15:8	00h
		0Eh	23:16	00h
Parameter ID MSB		0Fh	31:24	FFh
Parameter ID LSB		10h	7:0	84h
Parameter Minor Revision		11h	15:8	00h
Parameter Major Revision		12h	23:16	01h
Parameter Table Length (in DWPRDs)		13h	31:24	02h
Parameter Table Pointer (PTP)		14h	7:0	80h
		15h	15:8	00h
		16h	23:16	00h
Parameter ID MSB		17h	31:24	FFh
ID number ID LSB (RPMC)		18h	7:0	03h
Parameter Minor Revision		19h	15:8	00h
Parameter Major Revision		1Ah	23:16	01h
Parameter Table Length (in DWPRDs)		1Bh	31:24	02h
Parameter Table Pointer (PTP), First address of RPMC table		1Ch	7:0	90h
		1Dh	15:8	00h
		1Eh	23:16	00h
Parameter ID MSB		1Fh	31:24	FFh

**Table 5.4 JEDEC Basic Flash Parameter Table**

Description	Address (Byte)	Address (Bit)	Data
Minimum Sector Erase Sizes	30h	1:0	01b
Write Granularity		2	1b
Volatile Status Register Block Protect bits		3	0b
Write Enable Instruction Select for writing to Volatile Status Register		4	0b
Unused		7:5	111b
4KB Erase Instruction	31h	15:8	20h
Supports (1-1-2) Fast Read	32h	16	1b
Address Bytes		18:17	01b
Supports Double Transfer Rate (DTR) Clocking		19	1b
Supports (1-2-2) Fast Read		20	1b
Supports (1-4-4) Fast Read		21	1b
Supports (1-1-4) Fast Read		22	1b
Unused		23	1b
Reserved	33h	31:24	FFh
Flash memory Density (bits)	34h	7:0	FFh
	35h	15:8	FFh
	36h	23:16	FFh
Flash memory Density(bits)	37h	31:24	0Fh
1-4-4 Fast Read Wait Cycle Count	38h	4:0	00100b
1-4-4 Fast Read Mode bit Cycle Count		7:5	010b
1-4-4 Fast Read Instruction	39h	15:8	EBh
1-1-4 Fast Read Wait Cycle Count	3Ah	20:16	01000b
1-1-4 Fast Read Mode bit Cycle Count		23:21	000b
1-1-4 Fast Read Instruction	3Bh	31:24	6Bh
1-1-2 Fast Read Wait Cycle Count	3Ch	4:0	01000b
1-1-2 Fast Read Mode bit Cycle Count		7:5	000b
1-1-2 Fast Read Instruction	3Dh	15:8	3Bh
1-2-2 Fast Read Wait Cycle Count	3Eh	20:16	00000b
1-2-2 Fast Read Mode bit Cycle Count		23:21	100b
1-2-2 Fast Read Instruction	3Fh	31:24	BBh



**Table 5.4 JEDEC Basic Flash Parameter Table (Continued)**

Description	Address (Byte)	Address (Bit)	Data
Supports (2-2-2) Fast Read	40h	0	0b
Reserved		3:1	111b
Supports (4-4-4) Fast Read		4	1b
Reserved		7:5	111b
Reserved	43:41h	31:8	FFFFFFh
Reserved	45:44h	15:0	FFFFh
2-2-2 Fast Read Wait Cycle Count	46h	20:16	00000b
2-2-2 Fast Read Mode bit Cycle Count		23:21	000b
2-2-2 Fast Read Instruction	47h	31:24	FFh
Reserved	49:48h	15:0	FFFFh
4-4-4 Fast Read Wait Cycle Count	4Ah	20:16	00100b
4-4-4 Fast Read Mode bit Cycle Count		23:21	010b
4-4-4 Fast Read Instruction	4Bh	31:24	EBh
Erase Type 1 Size (4KB)	4Ch	7:0	0Ch
Erase Type 1 Instruction	4Dh	15:8	20h
Erase Type 2 Size (32KB)	4Eh	23:16	0Fh
Erase Type 2 Instruction	4Fh	31:24	52h
Erase Type 3 Size (64KB)	50h	7:0	10h(00h <sup>(1)</sup> )
Erase Type 3 Instruction	51h	15:8	D8h(FFh <sup>(1)</sup> )
Erase Type 4 Size (256KB)	52h	23:16	00h(12h <sup>(1)</sup> )
Erase Type 4 Instruction	53h	31:24	FFh(D8h <sup>(1)</sup> )
Multiplier from typical erase time to maximum erase time	57:54h	3:0	0010b
Sector Type 1 ERASE time (typ)		8:4	00110b
		10:9	01b
Sector Type 2 ERASE time (typ)		15:11	01000b
		17:16	01b
Sector Type 3 ERASE time (typ)		22:18	01010b (00000b <sup>(1)</sup> )
		24:23	01b(00b <sup>(1)</sup> )
Sector Type 4 ERASE time (typ)		29:25	00000b (00101b <sup>(1)</sup> )
	31:30	00b(10b <sup>(1)</sup> )	

**Note:**

1. Only for option K (256KB Block Size instead of 64KB)

**Table 5.4 JEDEC Basic Flash Parameter Table (Continued)**

Description	Address (Byte)	Address (Bit)	Data
Multiplier from typical time to maximum time for page or byte PROGRAM	58h	3:0	0010b
Page size		7:4	1000b (1001b <sup>(1)</sup> )
Page Program Typical time	5Ah:59h	12:8	00100b (01001b <sup>(1)</sup> )
		13	1b
Byte Program Typical time, first byte		17:14	1001b
		18	0b
Byte Program Typical time, additional byte		22:19	0000b
		23	0b
Chip Erase, Typical time	5Bh	28:24	10001b
Units		30:29	10b
Reserved		31	1b
Prohibited Operations During Program Suspend	5Ch	3:0	1100b
Prohibited Operations During Erase Suspend		7:4	1110b
Reserved	5Eh:5Dh	8	1b
Program Resume to Suspend Interval		12:9	0110b
Suspend in-progress program max latency		17:13	01100b
		19:18	10b
Erase Resume to Suspend Interval		23:20	0110b
Suspend in-progress erase max latency	5Fh	28:24	01100b
		30:29	10b
Suspend /Resume supported		31	0b
Program Resume Instruction	60h	7:0	7Ah
Program Suspend Instruction	61h	15:8	75h
Resume Instruction	62h	23:16	7Ah
Suspend Instruction	63h	31:24	75h
Reserved	64h	1:0	11b
Status Register Polling Device Busy		7:2	111101b

**Note:**

1. Only for option K (512B Page Size instead of 256B)

**Table 5.4 JEDEC Basic Flash Parameter Table (Continued)**

Description		Address (Byte)	Address (Bit)	Data
Exit Deep Power-down to next operation delay	3V	67h:65h	12:8	00010b
	1.8V			00010b
Exit Deep Power-down to next operation delay Units			14:13	10b
Exit Deep Power-down Instruction			22:15	ABh
Enter Deep Power-down Instruction			30:23	B9h
Deep Power-down Supported			31	0b
4-4-4 mode disable sequences (QPIDI)		69h:68h	3:0	1010b
4-4-4 mode enable sequences (QPIEN)			8:4	00100b
0-4-4 Mode Supported			9	1b
0-4-4 Mode Exit Method			15:10	110000b
0-4-4 Mode Entry Method:		6Ah	19:16	1100b
Quad Enable Requirements (QER)			22:20	010b
Hold or RESET Disable			23	0b
Reserved		6Bh	31:24	FFh
Volatile or Non-Volatile Register and Write Enable (WREN) Instruction for Status Register 1		6Ch	6:0	1100001b
Reserved			7	1b
Soft Reset and Rescue Sequence Support		6Eh:6Dh	13:8	110000b
Exit 4-Byte Addressing			23:14	1111101000b
Enter 4-Byte Addressing		6Fh	31:24	10101001b

**Table 5.5 Parameter Table (1): 4-byte Address Instruction Tables**

Description	Address (Byte)	Address (Bit)	Data
Support for (1-1-1) READ Command, Instruction = 13h	80h	0	1b
Support for (1-1-1) FAST_READ Command, Instruction = 0Ch		1	1b
Support for (1-1-2) FAST_READ Command, Instruction = 3Ch		2	1b
Support for (1-2-2) FAST_READ Command, Instruction = BCh		3	1b
Support for (1-1-4) FAST_READ Command, Instruction = 6Ch		4	1b
Support for (1-4-4) FAST_READ Command, Instruction = ECh		5	1b
Support for (1-1-1) Page Program Command, Instruction = 12h		6	1b
Support for (1-1-4) Page Program Command, Instruction = 34h		7	1b
Support for (1-4-4) Page Program Command, Instruction = 3Eh	81h	8	0b
Support for Erase Command-Type 1		9	1b
Support for Erase Command-Type 2		10	1b
Support for Erase Command-Type 3		11	1b (0b <sup>(1)</sup> )
Support for Erase Command-Type 4		12	0b (1b <sup>(1)</sup> )
Support for (1-1-1) DTR_READ Command, Instruction = 0Eh		13	1b
Support for (1-2-2) DTR_READ Command, Instruction = BEh		14	1b
Support for (1-4-4) DTR_READ Command, Instruction = EEh		15	1b
Support for volatile individual sector lock Read Command (E0h)	82h	16	1b
Support for volatile individual sector lock Write Command (E1h)		17	1b
Support for non-volatile individual sector lock Read Command (E2h)		18	1b
Support for non-volatile individual sector lock Write Command (E3h)		19	1b
Reserved		23:20	1111b
Reserved	83h	31:24	FFh
Sector Erase Instruction for Sector Type 1	84h	7:0	21h
Sector Erase Instruction for Sector Type 2	85h	15:8	5Ch
Sector Erase Instruction for Sector Type 3	86h	23:16	DCh (FFh <sup>(1)</sup> )
Sector Erase Instruction for Sector Type 4	87h	31:24	FFh (DCh <sup>(1)</sup> )

**Table 5.6 Parameter Table (2): RPMC Parameter Table**

Description		Address (Byte)	Address (Bit)	Data
Flash_Hardneing, 0= Flash Hardening is supported, 1=Flash Hardening is not supported.		90h	0	0b
MC_Size, 0=Monotonic counter size is 32 bit, 1=Reserved			1	0b
Busy_Polling_Method, 0=Poll for OP1 busy using OP2 RPMC Status[0], 1=Poll for OP1 busy using Status			2	0b
Reserved			3	1b
Num_Counter – 1 (Number of supported counters =4)			7:4	0011b
OP1 Opcode		91h	15:8	9Bh
OP2 Opcode		92h	23:16	96h
Update_Rate, Rate of Update = 5*(2**Update_Rate) seconds		93h	27:24	0000b
Reserved: Must be 0Fh			31:28	Fh
Read Counter Polling Delay Typical case to calculate HMAC two times (128 us)	polling delay_read counter	94h	4:0	01000b
	Units (00=1us, 01=16us, 10=128us, 11=1ms)		6:5	01b
	Reserved		7	1b
Write Counter Polling Short Delay (Worst Case Write+ one HMAC operation ), (160 us)	polling short_delay_write counter	95h	12:8	01010b
	Units (00=1us, 01=16us, 10=128us, 11=1ms)		14:13	01b
	Reserved		15	1b
Write Counter Polling Long Delay (Write+ one HMAC operation + Typical Sub Sector Erase Time ), (300 ms)	polling_long_delay_write counter	96h	20:16	10100b
	Units (00=1ms, 01=16ms, 10=128ms, 11=1)		22:21	11b
	Reserved		23	1b

## 6. REGISTERS

The device has various sets of Registers: Status, Function, Read, Extended Read and Autoboot. When the register is read continuously, the same data is output repeatedly until CE# goes HIGH.

### 6.1 STATUS REGISTER

Status Register Format and Status Register Bit Definitions are described in Tables 6.1 & 6.2.

**Table 6.1 Status Register Format**

	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
	SRWD	QE	BP3	BP2	BP1	BP0	WEL	WIP
Default	0	0	0	0	0	0	0	0

**Table 6.2 Status Register Bit Definition**

Bit	Name	Definition	Read-/Write	Type
Bit 0	WIP	Write In Progress Bit: "0" indicates the device is ready (default) "1" indicates a write cycle is in progress and the device is busy	R	Volatile
Bit 1	WEL	Write Enable Latch: "0" indicates the device is not write enabled (default) "1" indicates the device is write enabled	R/W <sup>1</sup>	Volatile
Bit 2	BP0	Block Protection Bit: (See Tables 6.4 for details) "0" indicates the specific blocks are not write-protected (default) "1" indicates the specific blocks are write-protected	R/W	Non-Volatile
Bit 3	BP1			
Bit 4	BP2			
Bit 5	BP3			
Bit 6	QE	Quad Enable bit: "0" indicates the Quad output function disable (default) "1" indicates the Quad output function enable	R/W	Non-Volatile
Bit 7	SRWD	Status Register Write Disable: (See Table 7.1 for details) "0" indicates the Status Register is not write-protected (default) "1" indicates the Status Register is write-protected	R/W	Non-Volatile

**Note:** WEL bit can be written by WREN and WRDI commands, but cannot by WRSR command.

The BP0, BP1, BP2, BP3, QE, and SRWD are non-volatile and volatile memory cells that can be written by a Write Status Register (WRSR) instruction. The default value of the BP0, BP1, BP2, BP3, QE, and SRWD bits were set to "0" at factory.

The function of Status Register bits are described as follows:

**WIP bit:** Write In Progress (WIP) is read-only, and can be used to detect the progress or completion of a Program, Erase, Write/Set Non-Volatile/OTP Register, or Gang Sector/Block Lock/Unlock operation. WIP is set to "1" (busy state) when the device is executing the operation. During this time the device will ignore further instructions except for Read Status/Function/Extended Read Register and Software/Hardware Reset instructions. In addition to the instructions, an Erase/Program Suspend instruction also can be executed during a Program or Erase operation. When an operation has completed, WIP is cleared to "0" (ready state) whether the operation is successful or not and the device is ready for further instructions.

**WEL bit:** Write Enable Latch (WEL) bit indicates the status of the internal write enable latch. When WEL bit is “0”, the internal write enable latch is disabled and the Write operations described in Table 6.3 are inhibited. When WEL bit is “1”, the Write operations are allowed. WEL bit is set by a Write Enable (WREN, 06h) instruction. Most of Write Non-Volatile/Volatile Register, Program and Erase instruction must be preceded by a WREN instruction.

WEL bit can be reset by a Write Disable (WRDI) instruction. It will automatically reset after the completion of any Write Non-Volatile Register, Program and Erase operation.

**Table 6.3 Instructions requiring WREN instruction ahead**

Instructions must be preceded by the WREN instruction		
Name	Hex Code	Operation
PP	02h	Serial Input Page Program (3-byte or 4-byte Address)
4PP	12h	Serial Input Page Program (4-byte Address)
PPQ	32h/38h	Quad Input Page Program (3-byte or 4-byte Address)
4PPQ	34h/3Eh	Quad Input Page Program (4-byte Address)
SER	D7h/20h	Sector Erase 4KB (3-byte or 4-byte Address)
4SER	21h	Sector Erase 4KB (4-byte Address)
BER32 (32KB)	52h	Block Erase 32KB (3-byte or 4-byte Address)
4BER32 (32KB)	5Ch	Block Erase 32KB (4-byte Address)
BER64 (64KB)	D8h	Block Erase 64KB (3-byte or 4-byte Address)
4BER64 (64KB)	DC	Block Erase 64KB (4-byte Address)
CER	C7h/60h	Chip Erase
<b>WRSR</b>	01h	Write Status Register
WRFR	42h	Write Function Register
SRPNV	65h	Set Read Parameters (Non-Volatile)
<b>SRPV<sup>(1)</sup></b>	63h	Set Read Parameters (Volatile)
SERPNV	85h	Set Extended Read Parameters (Non-Volatile)
SERP	83h	Set Extended Read Parameters (Volatile)
IRER	64h	Erase Information Row
IRP	62h	Program Information Row
WRABR	15h	Write AutoBoot Register
WRBRNV	18h	Write Non-Volatile Bank Address Register
<b>WRBRV<sup>(1)</sup></b>	C5h	Write Volatile Bank Address Register
WRDYB	FBh	Write DYB Register (4-byte Address)
4WRDYB	E1h	Write DYB Register (3-byte or 4-byte Address)
PGPPB	FDh	Write PPB (3-byte or 4-byte Address)
4PGPPB	E3h	Write PPB (4-byte Address)
ERPPB	E4h	Erase PPB
PGASP	2Fh	Program ASP
WRPLB	A6h	Write PPB Lock Bit
SFRZ	91h	Set FREEZE bit
GBLK	7Eh	GANG Sector/Block Lock
GBUN	98h	GANG Sector/Block Unlock
PGPWD	E8h	Program Password
<b>PNVDLR</b>	<b>43h</b>	<b>Program Non-Volatile Data Learning Pattern Register</b>
<b>WRVDLR</b>	<b>4Ah</b>	<b>Program Volatile Data Learning Pattern Register</b>

**Note:**

1. C0h command for SRPV operation and 17h command for WRBRV operation do not require WREN command ahead.

**BP3, BP2, BP1, BP0 bits:** The Block Protection (BP3, BP2, BP1 and BP0) bits are used to define the portion of the memory area to be protected. Refer to Table 6.4 for the Block Write Protection (BP) bit settings. When a defined combination of BP3, BP2, BP1 and BP0 bits are set, the corresponding memory area is protected. Any program or erase operation to that area will be inhibited.



**Note:** A Chip Erase (CER) instruction will be ignored unless all the Block Protection Bits are “0”s.

**SRWD bit:** The Status Register Write Disable (SRWD) bit operates in conjunction with the Write Protection (WP#) signal to provide a Hardware Protection Mode. When the SRWD is set to “0”, the Status Register is not write-protected. When the SRWD is set to “1” and the WP# is pulled low ( $V_{IL}$ ), the bits of Status Register (SRWD, QE, BP3, BP2, BP1, BP0) become read-only, and a WRSR instruction will be ignored. If the SRWD is set to “1” and WP# is pulled high ( $V_{IH}$ ), the Status Register can be changed by a WRSR instruction.

**QE bit:** The Quad Enable (QE) is a non-volatile bit in the Status Register that allows quad operation. When the QE bit is set to “0”, the pin WP# and HOLD#/RESET# are enabled. When the QE bit is set to “1”, the IO2 and IO3 pins are enabled.

**WARNING:** The QE bit must be set to “0” if WP# or HOLD#/RESET# pin (or ball) is tied directly to the power supply.

**Table 6.4 Block assignment by Block Write Protect (BP) Bits**

Status Register Bits				Byte Protected	Block Size =64KB Protected Memory Area (256Mb, 512Blocks)	
BP3	BP2	BP1	BP0		TBS = 0, Top area	TBS = 1, Bottom area
0	0	0	0	0KB	0 ( None)	0 ( None)
0	0	0	1	64KB	1 (1 block: 511 <sup>st</sup> )	1 (1 block: 0 <sup>th</sup> )
0	0	1	0	128KB	2 (2 blocks: 510 <sup>th</sup> and 511 <sup>st</sup> )	2 (2 blocks: 0 <sup>th</sup> and 1 <sup>st</sup> )
0	0	1	1	256KB	3 (4 blocks: 508 <sup>th</sup> to 511 <sup>st</sup> )	3 (4 blocks: 0 <sup>th</sup> to 3 <sup>rd</sup> )
0	1	0	0	512KB	4 (8 blocks: 504 <sup>th</sup> to 511 <sup>st</sup> )	4 (8 blocks: 0 <sup>th</sup> to 7 <sup>th</sup> )
0	1	0	1	1MB	5 (16 blocks: 496 <sup>th</sup> to 511 <sup>st</sup> )	5 (16 blocks: 0 <sup>th</sup> to 15 <sup>th</sup> )
0	1	1	0	2MB	6 (32 blocks: 480 <sup>th</sup> to 511 <sup>st</sup> )	6 (32 blocks: 0 <sup>th</sup> to 31 <sup>st</sup> )
0	1	1	1	4MB	7 (64 blocks: 448 <sup>th</sup> to 511 <sup>st</sup> )	7 (64 blocks: 0 <sup>th</sup> to 63 <sup>rd</sup> )
1	0	0	0	8MB	8 (128 blocks: 384 <sup>th</sup> to 511 <sup>st</sup> )	8 (128 blocks: 0 <sup>th</sup> to 127 <sup>th</sup> )
1	0	0	1	16MB	9 (256 blocks: 256 <sup>th</sup> to 511 <sup>st</sup> )	9 (256 blocks: 0 <sup>th</sup> to 255 <sup>th</sup> )
1	0	1	0	24MB	10 (384 blocks : 128 <sup>th</sup> to 511 <sup>st</sup> )	10 (384 blocks : 0 <sup>th</sup> to 383 <sup>rd</sup> )
1	0	1	1	28MB	11 (448 blocks : 64 <sup>th</sup> to 511 <sup>st</sup> )	11 (448 blocks : 0 <sup>th</sup> to 447 <sup>th</sup> )
1	1	0	0	30MB	112 (480 blocks : 32 <sup>nd</sup> to 511 <sup>st</sup> )	12 (480 blocks : 0 <sup>th</sup> to 479 <sup>th</sup> )
1	1	0	1	31MB	13 (496 blocks : 16 <sup>th</sup> to 511 <sup>st</sup> )	13 (496 blocks : 0 <sup>th</sup> to 495 <sup>th</sup> )
1	1	1	0	32256KB	14 (504 blocks : 8 <sup>th</sup> to 511 <sup>st</sup> )	14 (504 blocks : 0 <sup>th</sup> to 503 <sup>rd</sup> )
1	1	1	1	32MB (All)	15 (512 blocks : 0 <sup>th</sup> to 511 <sup>st</sup> )	15 (512 blocks : 0 <sup>th</sup> to 511 <sup>st</sup> )

Status Register Bits				Block Size = 256KB <sup>(2)</sup> , Protected Memory Area (256Mb, 128Blocks)	
BP3	BP2	BP1	BP0	TBS(T/B selection) = 0, Top area	TBS(T/B selection) = 1, Bottom area
0	0	0	0	0( None)	0( None)
0	0	0	1	1(1 block : 127 <sup>th</sup> )	1(1 block : 0 <sup>th</sup> )
0	0	1	0	2(2 blocks : 126 <sup>th</sup> and 127 <sup>th</sup> )	2(2 blocks : 0 <sup>th</sup> and 1 <sup>st</sup> )
0	0	1	1	3(4 blocks : 124 <sup>th</sup> to 127 <sup>th</sup> )	3(4 blocks : 0 <sup>th</sup> to 3 <sup>rd</sup> )
0	1	0	0	4(8 blocks : 120 <sup>th</sup> to 127 <sup>th</sup> )	4(8 blocks : 0 <sup>th</sup> to 7 <sup>th</sup> )
0	1	0	1	5(16 blocks : 112 <sup>th</sup> to 127 <sup>th</sup> )	5(16 blocks : 0 <sup>th</sup> to 15 <sup>th</sup> )
0	1	1	0	6(32 blocks : 96 <sup>th</sup> to 127 <sup>th</sup> )	6(32 blocks : 0 <sup>th</sup> to 31 <sup>st</sup> )
0	1	1	1	7(64 blocks : 64 <sup>th</sup> to 127 <sup>th</sup> )	7(64 blocks : 0 <sup>th</sup> to 63 <sup>rd</sup> )
1	x	x	x	8-15(128 blocks : 0 <sup>th</sup> to 127 <sup>th</sup> ) All blocks	8-15(128 blocks : 0 <sup>th</sup> to 127 <sup>th</sup> ) All blocks

**Notes:**
**1. x is don't care**
**2. For Optional 256KB BP Table, see the Ordering Information (Option "K")**

## 6.2 FUNCTION REGISTER

Function Register Format and Bit definition are described in Table 6.5 and Table 6.6.

**Table 6.5 Function Register Format**

	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
	IRL3	IRL2	IRL1	IRL0	ESUS	PSUS	TBS	Reserved
Default	0	0	0	0	0	0	0	0 or 1

**Note:**

1. “0” in 16-pin SOIC package, and “1” in WSON package.

**Table 6.6 Function Register Bit Definition**

Bit	Name	Definition	Read /Write	Type
Bit 0	Reserved	Reserved	R only	Reserved
Bit 1	TBS	Top/Bottom Selection. (See Table 6.4 for details) “0” indicates Top area. “1” indicates Bottom area.	R/W	OTP
Bit 2	PSUS	Program suspend bit: “0” indicates program is not suspend “1” indicates program is suspend	R	Volatile
Bit 3	ESUS	Erase suspend bit: “0” indicates Erase is not suspend “1” indicates Erase is suspend	R	Volatile
Bit 4	IR Lock 0	Lock the Information Row 0: “0” indicates the Information Row can be programmed “1” indicates the Information Row cannot be programmed	R/W	OTP
Bit 5	IR Lock 1	Lock the Information Row 1: “0” indicates the Information Row can be programmed “1” indicates the Information Row cannot be programmed	R/W	OTP
Bit 6	IR Lock 2	Lock the Information Row 2: “0” indicates the Information Row can be programmed “1” indicates the Information Row cannot be programmed	R/W	OTP
Bit 7	IR Lock 3	Lock the Information Row 3: “0” indicates the Information Row can be programmed “1” indicates the Information Row cannot be programmed	R/W	OTP

**Note:** Once OTP bits of Function Register are written to “1”, it cannot be modified to “0” any more.

**TBS bit:** BP0~3 area assignment can be changed from Top (default) to Bottom by setting TBS bit to “1”. However, once Bottom is selected, it cannot be changed back to Top since TBS bit is OTP. See Table 6.4 for details.

**PSUS bit:** The Program Suspend Status bit indicates when a Program operation has been suspended. The PSUS changes to “1” after a suspend command is issued during the program operation. Once the suspended Program resumes, the PSUS bit is reset to “0”.

**ESUS bit:** The Erase Suspend Status bit indicates when an Erase operation has been suspended. The ESUS bit is “1” after a suspend command is issued during an Erase operation. Once the suspended Erase resumes, the ESUS bit is reset to “0”.

**IR Lock bit 0 ~ 3:** The default is “0” so that the Information Row can be programmed. If the bit is set to “1”, it cannot be changed back to “0” again since IR Lock bits are OTP.

### 6.3 READ REGISTER AND EXTENDED REGISTER

Read Register format and bit definitions are described below. Read Register and Extended Read Register consist of a pair of rewritable non-volatile register and volatile register, respectively. During power up sequence, volatile register will be loaded with the value of non-volatile value.

#### 6.3.1 READ REGISTER

Table 6.7 and Table 6.8 define all bits that control features in SPI/QPI modes.

The Dummy Cycle bits (P6, P5, P4, P3) define how many dummy cycles are used during various READ modes. The wrap selection bits (P2, P1, P0) define burst length with an enable bit.

The SET READ PARAMETERS Operations (SRPNV: 65h, SRPV: C0h or 63h) are used to set dummy cycles, and burst length with wrap around. SRPNV is used to set the non-volatile register and SRPV is used to set the volatile register.

**Table 6.7 Read Register Parameter Bit Table**

	P7	P6	P5	P4	P3	P2	P1	P0
	Reserved	Dummy Cycles	Dummy Cycles	Dummy Cycles	Dummy Cycles	Wrap Enable	Burst Length	Burst Length
Default	0	0	0	0	0	0	0	0

**Table 6.8 Read Register Bit Definition**

Bit	Name	Definition	Read-Write	Type
P0	Burst Length	Burst Length	R/W	Non-Volatile and Volatile
P1	Burst Length	Burst Length	R/W	Non-Volatile and Volatile
P2	Burst Length Set Enable	Burst Length Set Enable Bit: "0" indicates disable (default) "1" indicates enable	R/W	Non-Volatile and Volatile
P3	Dummy Cycles	Number of Dummy Cycles: Bits1 to Bit4 can be toggled to select the number of dummy cycles (1 to 15 cycles)	R/W	Non-Volatile and Volatile
P4	Dummy Cycles		R/W	Non-Volatile and Volatile
P5	Dummy Cycles		R/W	Non-Volatile and Volatile
P6	Dummy Cycles		R/W	Non-Volatile and Volatile
P7	Reserved	Reserved	R	Reserved

**Table 6.9 Burst Length Data**

	P1	P0
8 bytes	0	0
16 bytes	0	1
32 bytes	1	0
64 bytes	1	1

**Table 6.10 Wrap Function**

Wrap around boundary	<b>P2</b>
Whole array regardless of P1 and P0 value	0
Burst Length set by P1 and P0	1

**Table 6.11 Read Dummy Cycles vs Max Frequency**
**3.0V Device (IS25RLP)**

P[6:3]	Dummy Cycles <sup>2,3</sup>	Fast Read <sup>5</sup> 0Bh/0Ch		Fast Read Dual Output 3Bh/3Ch	Fast Read Dual IO BBh/BCh	Fast Read Quad Output 6Bh/6Ch	Fast Read Quad IO EBh/ECh	FRDTR 0Dh/0Eh	FRDDTR BDh/BEh	FRQDTR EDh/EEh
		SPI	QPI	SPI	SPI	SPI	SPI, QPI	SPI/QPI	SPI <sup>4</sup>	SPI, QPI
0	Default <sup>1</sup>	<b>166MHz</b>	<b>81MHz</b>	<b>166MHz</b>	<b>104MHz</b>	<b>145MHz</b>	<b>81MHz</b>	<b>80/69MHz</b>	<b>60MHz</b>	<b>69MHz</b>
1	1	98MHz	23MHz	75MHz	52MHz	63MHz	23MHz	50/11MHz	29MHz	11MHz
2	2	110MHz	34MHz	84MHz	80MHz	75MHz	34MHz	63/23MHz	40MHz	23MHz
3	3	122MHz	46MHz	98MHz	98MHz	87MHz	46MHz	75/34MHz	52MHz	34MHz
4	4	133MHz	58MHz	133MHz	<b>104MHz</b>	98MHz	58MHz	80/46MHz	<b>60MHz</b>	46MHz
5	5	145MHz	69MHz	140MHz	122MHz	110MHz	69MHz	80/58MHz	75MHz	58MHz
6	6	156MHz	<b>81MHz</b>	150MHz	133MHz	122MHz	<b>81MHz</b>	80/69MHz	80MHz	69MHz
7	7	166MHz	93MHz	166MHz	145MHz	133MHz	93MHz	80/80MHz	80MHz	80MHz
8	8	<b>166MHz</b>	104MHz	<b>166MHz</b>	156MHz	<b>145MHz</b>	104MHz	80/80MHz	80MHz	80MHz
9	9	166MHz	122MHz	166MHz	166MHz	156MHz	122MHz	80/80MHz	80MHz	80MHz
10	10	166MHz	127MHz	166MHz	166MHz	166MHz	127MHz	80/80MHz	80MHz	80MHz
11	11	166MHz	139MHz	166MHz	166MHz	166MHz	139MHz	80/80MHz	80MHz	80MHz
12	12	166MHz	151MHz	166MHz	166MHz	166MHz	151MHz	80/80MHz	80MHz	80MHz
13	13	166MHz	162MHz	166MHz	166MHz	166MHz	162MHz	80/80MHz	80MHz	80MHz
14	14	166MHz	166MHz	166MHz	166MHz	166MHz	166MHz	80/80MHz	80MHz	80MHz
15	15	166MHz	166MHz	166MHz	166MHz	166MHz	166MHz	80/80MHz	80MHz	80MHz

**1.8V Device (IS25RWP)**

P[6:3]	Dummy Cycles <sup>2,3</sup>	Fast Read <sup>5</sup> 0Bh/0Ch		Fast Read Dual Output 3Bh/3Ch	Fast Read Dual IO BBh/BCCh	Fast Read Quad Output 6Bh/6Ch	Fast Read Quad IO EBh/ECh	FRDTR 0Dh/0Eh	FRDDTR BDh/BEh	FRQDTR EDh/EEh
		SPI	QPI	SPI	SPI	SPI	SPI, QPI	SPI/QPI	SPI <sup>4</sup>	SPI, QPI
0	Default <sup>1</sup>	<b>166MHz</b>	<b>81MHz</b>	<b>166MHz</b>	<b>98MHz</b>	<b>138MHz</b>	<b>81MHz</b>	<b>80/58MHz</b>	<b>60MHz</b>	<b>58MHz</b>
1	1	98MHz	23MHz	75MHz	46MHz	63MHz	23MHz	46/10MHz	23MHz	10MHz
2	2	110MHz	34MHz	84MHz	75MHz	75MHz	34MHz	58/15MHz	40MHz	15MHz
3	3	122MHz	46MHz	98MHz	87MHz	87MHz	46MHz	69/20MHz	50MHz	20MHz
4	4	133MHz	58MHz	122MHz	<b>98MHz</b>	98MHz	58MHz	75/30MHz	<b>60MHz</b>	30MHz
5	5	145MHz	69MHz	133MHz	110MHz	110MHz	69MHz	80/40MHz	70MHz	40MHz
6	6	156MHz	<b>81MHz</b>	145MHz	122MHz	122MHz	<b>81MHz</b>	80/58MHz	75MHz	58MHz
7	7	166MHz	93MHz	156MHz	133MHz	133MHz	93MHz	80/69MHz	80MHz	69MHz
8	8	<b>166MHz</b>	104MHz	<b>166MHz</b>	145MHz	<b>138MHz</b>	104MHz	<b>80/80MHz</b>	80MHz	80MHz
9	9	166MHz	122MHz	166MHz	156MHz	140MHz	122MHz	80/80MHz	80MHz	80MHz
10	10	166MHz	127MHz	166MHz	166MHz	145MHz	127MHz	80/80MHz	80MHz	80MHz
11	11	166MHz	139MHz	166MHz	166MHz	151MHz	139MHz	80/80MHz	80MHz	80MHz
12	12	166MHz	151MHz	166MHz	166MHz	166MHz	151MHz	80/80MHz	80MHz	80MHz
13	13	166MHz	162MHz	166MHz	166MHz	166MHz	162MHz	80/80MHz	80MHz	80MHz
14	14	166MHz	166MHz	166MHz	166MHz	166MHz	166MHz	80/80MHz	80MHz	80MHz
15	15	166MHz	166MHz	166MHz	166MHz	166MHz	166MHz	80/80MHz	80MHz	80MHz

**Notes:**

- Default dummy cycles are as follows.

Operation	Command		Dummy Cycles		Comment
	SPI mode	QPI mode	SPI mode	QPI mode	
Fast Read	0Bh/0Ch	0Bh/0Ch	8	6	RDUID, IRRD instructions are also applied.
Fast Read Dual Output	3Bh/3Ch	-	8	-	
Fast Read Quad Output	6Bh/6Ch	-	8	-	
Fast Read Dual IO	BBh/BCCh	-	4	-	
Fast Read Quad IO	EBh/ECh	EBh/ECh	6	6	
Fast Read DTR	0Dh/0Eh	0Dh/0Eh	8	6	
Fast Read Dual IO DTR	BDh/BEh	-	4	-	
Fast Read Quad IO DTR	EDh/EEh	EDh/EEh	6	6	

- Enough number of dummy cycles must be applied to execute properly the AX read operation.
- Must satisfy bus I/O contention. For instance, if the number of dummy cycles and AX bits cycles are same, then X must be Hi-Z.
- QPI mode is not available for FRDDTR command.
- RDUID, IRRD instructions are also applied.
- RPMC OP2 operation frequency: 104MHz with fixed 8-dummy cycle.

### 6.3.2 EXTENDED READ REGISTER

Table 6.12 and Table 6.13 define all bits that control features in SPI/QPI modes. The ODS2, ODS1, ODS0 (EB7, EB6, EB5) bits provide a method to set and control driver strength.

The SET EXTENDED READ PARAMETERS Operations (SERPNV: 85h, SERPV: 83h) are used to set all the Extended Read Register bits, and can thereby define the output driver strength used during READ modes. SRPNV is used to set the non-volatile register and SRPV is used to set the volatile register.

**Table 6.12 Extended Read Register Bit Table**

Bit	EB7	EB6	EB5	EB4	EB3	EB2	EB1	EB0
Name	ODS2	ODS1	ODS0	DLPEN	E_ERR	P_ERR	PROT_E	Reserved
Default	1	1	1	0	0	0	0	0

**Table 6.13 Extended Read Register Bit Definition**

Bit	Name	Definition	Read-/Write	Type
EB0	Reserved	Reserved	R	Reserved
EB1	PROT_E	Protection Error Bit: "0" indicates no error "1" indicates protection error in an Erase or a Program operation	R	Volatile
EB2	P_ERR	Program Error Bit: "0" indicates no error "1" indicates a Program operation failure or protection error	R	Volatile
EB3	E_ERR	Erase Error Bit: "0" indicates no error "1" indicates an Erase operation failure or protection error	R	Volatile
EB4	DLPEN	DLP Enable Bit "0" indicates Disable "1" indicates Enable	R/W	Non-Volatile and Volatile
EB5	ODS0	Output Driver Strength: Output Drive Strength can be selected according to Table 6.14	R/W	Non-Volatile and Volatile
EB6	ODS1		R/W	Non-Volatile and Volatile
EB7	ODS2		R/W	Non-Volatile and Volatile

**Table 6.14 Driver Strength Table**

ODS2	ODS1	ODS0	Description	Remark
0	0	0	Reserved	
0	0	1	12.50%	
0	1	0	25%	
0	1	1	37.50%	
1	0	0	Reserved	
1	0	1	75%	
1	1	0	100%	
1	1	1	50%	Default

**PROT\_E bit:** The Protection Error bit indicates whether an Erase or Program operation has attempted to modify a protected array sector or block, or to access a locked Information Row region. When the bit is set to “1” it indicates that there was an error or errors in previous Erase or Program operations. See Table 6.15 for details.

**P\_ERR bit:** The Program Error bit indicates whether a Program operation has succeeded or failed, or whether a Program operation has attempted to program a protected array sector/block or a locked Information Row region. When the bit is set to “1” it indicates that there was an error or errors in previous Program or Write/Set Non-Volatile Register operations. See Table 6.15 for details.

**E\_ERR bit:** The Erase Error bit indicates whether an Erase operation has succeeded or failed, or whether an Erase operation has attempted to erase a protected array sector/block or a locked Information Row region. When the bit is set to “1” it indicates that there was an error or errors in previous Erase or Write/Set Non-Volatile Register operations. See Table 6.15 for details.

**Table 6.15 Instructions to set PROT\_E, P\_ERR, or E\_ERR bit**

Instructions	Description
PP/4PP/PPQ/4PPQ/PGPPB/4PGPPB/PGPWD	The commands will set the P_ERR if there is a failure in the operation. Attempting to program within the protected array sector/block or within an erase suspended sector/block will result in a programming error with P_ERR and PROT_E set to “1”.
IRP	The command will set the P_ERR if there is a failure in the operation. In attempting to program within a locked Information Row region, the operation will fail with P_ERR and PROT_E set to 1.
PGASP	The command will set the P_ERR if there is a failure in the operation. Attempting to program ASPR[2:1] after the Protection Mode is selected or attempting to program ASPR[2:1] = 00b will result in a programming error with P_ERR and PROT_E set to “1”.
UNPWD	If the UNPWD command supplied password does not match the hidden internal password, the UNPWD operation fails in the same manner as a programming operation on a protected sector/block and sets the P_ERR and PROT_E to “1”.
WRSR/WRABR/SRPNV/SERPNV/WRBRNV/PNVDLR	The update process for the non-volatile register bits involves an erase and a program operation on the non-volatile register bits. If either the erase or program portion of the update fails, the related error bit (P_ERR or E_ERR) will be set to “1”. Only for WRSR command, when Status Register is write-protected by SRWD bit and WP# pin, attempting to write the register will set PROT_E and E_ERR to “1”.
WRFR	The commands will set the P_ERR if there is a failure in the operation.
SER/4SER/BER32K/4BER32K/BER64K/4BER64K/BER256K/4BER256K/CER/IRER/ERPPB	The commands will set the E_ERR if there is a failure in the operation. E_ERR and PROT_E will be set to “1” when the user attempts to erase a protected main memory sector/block or a locked Information Row region. Chip Erase (CER) command will set E_ERR and PROT_E if any blocks are protected by Block Protection bits (BP3~BP0). But Chip Erase (CER) command will not set E_ERR and PROT_E if sectors/blocks are protected by ASP (DYB bits or PPB bits) only.

**Notes:**

1. OTP bits in the Function Register and TBPARM (OTP bit) in the ASP Register may only be programmed to “1”. Writing of the bits back to “0” is ignored and no error is set.
2. Read only bits and partially protected bits by FREEZE bit in registers are never modified by a command so that the corresponding bits in the Write/Set Register command data byte are ignored without setting any error indication.
3. Once the PROT\_E, P\_ERR, and E\_ERR error bits are set to “1”, they remains set to “1” until they are cleared to “0” with a Clear Extended Read Register (CLERP) command. This means that those error bits must be cleared through the CLERP command. Alternatively, Hardware Reset, or Software Reset may be used to clear the bits.
4. Any further command will be executed even though the error bits are set to “1”.



## 6.4 AUTOBOOT REGISTER

AutoBoot Register Bit (32 bits) definitions are described in Table 6.16.

**Table 6.16 AutoBoot Register Parameter Bit Table**

Bits	Symbols	Function	Type	Default Value	Description
AB[31:5]	ABSA	AutoBoot Start Address	Non-Volatile	0000000h	32 byte boundary address for the start of boot code access
AB[4:1]	ABSD	AutoBoot Start Delay	Non-Volatile	0h	Number of initial delay cycles between CE# going low and the first bit of boot code being transferred, and it is the same as dummy cycles of FRD (QE=0) or FRQIO (QE=1). Example: The number of initial delay cycles is 8 (QE=0) or 6 (QE=1) when AB[4:1]=0h (Default setting).
AB0	ABE	AutoBoot Enable	Non-Volatile	0	1 = AutoBoot is enabled 0 = AutoBoot is not enabled

## 6.5 BANK ADDRESS REGISTER

Related Commands: Read Volatile Bank Address Register (RDBR 16h/C8h), Write Volatile Bank Address Register (WRBRV 17h/C5h), Write Non-Volatile Bank Address Register (WRBRNV 18h), Enter 4-byte Address Mode (EN4B B7h), and Exit 4-byte Address Mode (EX4B 29h).

Bank Address Register Bit (8 bits) definitions are described in Table 6.17 and Table 6.18.

**Table 6.17 Bank Address Register Bit Table**

	BA7	BA6	BA5	BA4	BA3	BA2	BA1	BA0
	EXTADD	Reserved	Reserved	Reserved	Reserved	Reserved	Reserved	BA24
Default	0	0	0	0	0	0	0	0

**Table 6.18 Bank Address Register Bit Definition**

Bit	Name	Definition	Read- /Write	Type
BA0	BA24	Enables 128Mb segment selection in 3-byte addressing "0" indicates lower 128Mb segment is selected. "1" indicates upper 128Mb segment is selected. "Reserved" in 128Mb.	R/W	Non-Volatile and Volatile
BA1	Reserved	Reserved	R	Reserved
BA2	Reserved	Reserved	R	Reserved
BA3	Reserved	Reserved	R	Reserved
BA4	Reserved	Reserved	R	Reserved
BA5	Reserved	Reserved	R	Reserved
BA6	Reserved	Reserved	R	Reserved
BA7	EXTADD	3-byte or 4-byte addressing mode selection Bit: "0" indicates 3-byte addressing mode "1" indicates 4-byte addressing mode	R/W	Non-Volatile and Volatile

**BA24:** The Bank Address Register supplies additional high order bits of the main flash array byte boundary address for legacy commands that supply only the low order 24 bits of address. The Bank Address is used as the high bits of address (above A23) for all 3-byte address commands when EXTADD=0.

The Bank Address is not used when EXTADD = 1 and traditional 3-byte address commands are instead required to provide all four bytes of address.

**EXTADD:** Extended Address (EXTADD) controls the address field size for legacy SPI commands. When shipped from factory, it is cleared to "0" for 3-byte addressing mode.

When set to "1" for 4-byte addressing mode, the legacy commands will require 4 bytes of address for the address field

## 6.6 ADVANCED SECTOR/BLOCK PROTECTION (ASP) RELATED REGISTER

### 6.6.1 ADVANCED SECTOR/BLOCK PROTECTION REGISTER (ASPR)

Related Commands: Read ASP (RDASP 2Bh) and Program ASP (PGASP 2Fh).

Advanced Sector/Block Protection (ASP) Register Bit (16 bits) definitions are described in Tables 6.19 and 6.20.

**Table 6.19 Advanced Sector/Block Protection Register (ASPR) Bit Table**

	15	14 to 7	6	5	4	3	2	1	0
	TBPARAM <sup>(1)</sup>	Reserved	Reserved	Reserved	Reserved	Reserved	PWDMLB	PSTMLB	Reserved
Default	1	1	1	1	1	1	1	1	1

**Table 6.20 Advanced Sector/Block Protection Register (ASPR) Bit Definition**

Bit	Name	Definition	Read-/Write	Type
0	Reserved	Reserved	R	Reserved
1	PSTMLB	Persistent Protection Mode Lock Bit "0" = Persistent Protection Mode permanently enabled. "1" = Persistent Protection Mode not permanently enabled.	R/W	OTP
2	PWDMLB	Password Protection Mode Lock Bit "0" = Password Protection Mode permanently enabled. "1" = Password Protection Mode not permanently enabled.	R/W	OTP
14:3	Reserved	Reserved	R	Reserved
15	TBPARAM <sup>(1)</sup>	Configures Parameter Sectors location "0" = 4KB physical sectors at top, (high address) "1" = 4KB physical sectors at bottom (Low address)	R/W	OTP <sup>(1)</sup>

**Note:**

1. TBPARAM bit is Reserved for optional block size of 256KB device

The Advanced Sector/Block Protection Register (ASPR) is used to permanently configure the behavior of Advanced Sector/Block Protection (ASP) features and parameter sectors location.

**PWDMLB (ASPR[2]) and PSTMLB (ASPR[1]) bits:** When shipped from the factory, all devices default ASP to the Persistent Protection Mode, with all sectors unprotected, when power is applied. The device programmer or host system must then choose which sector/block protection method to use. Programming either of the Protection Mode Lock Bits locks the part permanently in the selected mode:

- ASPR[2:1] = 11 = No ASP mode selected, Persistent Protection Mode is the default.
- ASPR[2:1] = 10 = Persistent Protection Mode permanently selected.
- ASPR[2:1] = 01 = Password Protection Mode permanently selected.
- ASPR[2:1] = 00 = Illegal condition, attempting to program both bits to zero results in a programming failure and the program operation will abort. It will result in a programming error with P\_ERR set to 1.

As a result, PWDMLB and PSTMLB are mutually exclusive, only one may be programmed to zero.

ASPR programming rules:

- If the Password Protection Mode is chosen, the password must be programmed prior to setting the corresponding bit.
- Once the Protection Mode is selected, the ASPR[2:1] bits are permanently protected from programming and no further change to the ASPR[2:1] is allowed. Attempting to program ASPR[2:1] after selected will result in a programming error with P\_ERR set to 1. The programming time of the ASPR is the same as the typical page programming time. The system can determine the status of the ASPR programming operation by reading the WIP bit in the Status Register or Extended Read Register.
- TBPARAM bit can be programmed even after ASPR[2:1] bits are programmed while the FREEZE bit in the PPB Lock Register is "0".

**TBPARM bit:** TBPARM defines the logical location of the parameter block. The parameter block consists of thirty-two 4KB sectors, which replace two 64KB blocks. When TBPARM is set to a “0” the parameter block is in the top of the memory array address space. When TBPARM is set to a “1” the parameter block is at the Bottom of the array. TBPARM is OTP and set to a “1” when it ships from Factory. If TBPARM is programmed to “0”, an attempt to change it back to “1” will fail and ignore the Program.

The desired state of TBPARM must be selected during the initial configuration of the device during system manufacture; before the first program or erase operation on the main flash array. TBPARM must not be programmed after programming or erasing is done in the main flash array.

TBS can be programmed independent of TBPARM. Therefore, the user can select to store parameter information from the bottom of the array and protect boot code starting at the top of the array, and vice versa. Or the user can select to store and protect the parameter information starting from the top or bottom together.

**Note:** For optional device with block size of 256KB, there is no parameter blocks, so TBPARM bit is reserved.

### 6.6.2 PASSWORD REGISTER

Related Commands: Read Password (RDPWD E7h), Program Password (PGPWD E8h), and Unlock Password (UNPWD, E9h).

**Table 6.21 Password Register Bit Definition**

Bit	Name	Definition	Default	Read- /Write	Type
63:0	PSWD	64 bit hidden password: The password is no longer readable after the password protection mode is selected by programming ASPR bit 2 to zero.	FFFFFFFF- FFFFFFFFh	R/W	OTP

### 6.6.3 PPB LOCK REGISTER

Related Commands: Read PPB Lock Bit (RDPLB A7h), Write PPB Lock Bit (WRPLB A6h), and Set FREEZE Bit (SFRZ 91h).

**Table 6.22 PPB Lock Register Bit Definition**

Bit	Name	Definition	Default	Read- /Write	Type
0	PPBLK	PPB Lock bit: Protect PPB Array “0” = PPB array protected until next power cycle or Hardware Reset “1” = PPB array may be programmed or erased.	Persistent: 1 Password: 0	R/W	Volatile
6:1	Reserved	Reserved	Reserved	R	Reserved
7	FREEZE	Lock current state of BP3-0 bits in Status Register, TBS in Function Register and TBPARM in ASPR, and Information Row (IR) regions. “1” = Locked “0” = Un-locked	0	R/W	Volatile

**PPBLK bit:** The PPB Lock bit is a volatile bit for protecting all PPB bits. When cleared to 0, it locks all PPBs, when set to 1, it allows the PPBs to be changed. The WRPLB command is used to clear the PPB Lock bit to 0. The PPB Lock bit must be cleared to 0 only after all the PPBs are configured to the desired settings.

In Persistent Protection mode, the PPB Lock bit is set to 1 during POR or Hardware Reset. When cleared to 0, no software command sequence can set the PPB Lock bit to 1, only another Hardware Reset or power-up can set the PPB Lock bit.

In the Password Protection mode, the PPB Lock bit is cleared to 0 during POR or Hardware Reset. The PPB Lock bit can only be set to 1 by the Unlock Password command.

**FREEZE bit:** FREEZE bit, when set to “1”, locks the current state of BP3-0 in Status Register, TBS in the Function Register, TBPARM in the Advanced Sector/Block Protection Register, and the Information Row. This prevents writing, programming, or erasing these areas. As long as FREEZE remains cleared to logic “0”, BP3-0 in Status Register, TBS in the Function Register, and TBPARM in the Advanced Sector/Block Protection Register are writable and the Information Row is programmable. Once FREEZE has been written to a logic “1” it can only be cleared to a logic “0” by a power-on cycle or a Hardware Reset. Software Reset will not affect the state of FREEZE. The FREEZE is volatile and the default state of FREEZE after power-on is “0”. The FREEZE can be set to “1” by a SFRZ command.

#### 6.6.4 PPB REGISTER

Related Commands: Read PPB (RDPPB FCh or 4RDPPB E2h)), Program PPB (PGPPB FDh or 4PGPPB E3h), and Erase PPB (ERPPB E4h).

**Table 6.23 PPB Register Bit Definition**

Bit	Name	Definition	Default	Read-Write	Type
7:0	PPB	Read or Program per sector/block PPB: 00h = PPB for the sector/block addressed by the RDPPB or PGPPB command is programmed to “0”, protecting that sector/block from program or erase operations. FFh = PPB for the sector/block addressed by the RDPPB or PGPPB command is erased to “1”, not protecting that sector/block from program or erase operations.	FFh	R/W	Non-Volatile

#### 6.6.5 DYB REGISTER

Related Commands: Read DYB (RDDYB FAh or 4RDDYB E0h) and Write DYB (WRDYB FBh or 4WRDYB E1h).

**Table 6.24 DYB Register Bit Definition**

Bit	Name	Definition	Default	Read-Write	Type
7:0	DYB	Read or Write per sector/block DYB: 00h = DYB for the sector/block addressed by the RDDYB or WRDYB command is cleared to “0”, protecting that sector/block from program or erase operations. FFh = DYB for the sector/block addressed by the RDDYB or WRDYB command is set to “1”, not protecting that sector/block from program or erase operations.	FFh	R/W	Volatile

**6.6.6 DATA LEARNING PATTERN REGISTER**

Related Commands: Read DLP (RDDLP 41h), Program Non-Volatile DLP Register (PNVDLR 43h), Write Volatile DLP Register (WRVDLR 4Ah).

The Data Learning Pattern (DLP) resides in an 8-bit NON-Volatile Data Learning Register (NVDLR) as well as an 8-bit Volatile Data Learning Register (VDLR). When shipped from the factory, default value is 00h.

A copy of the data pattern in NVDLR will also be written to the VDLR. The VDLR can be written to at any time, but on power cycles the data pattern will return back to data pattern in NVDLR.

During Data Training phase, DLP will come from VDLR.

**Table 6.25 Non-Volatile Data Learning Register (NVDLR)**

Bit	Name	Definition	Default	Read- /Write	Type
7:0	NVDLP	Non-Volatile Data Learning Pattern: The value that may be transferred to the host during DDR read command latency (dummy) cycles to provide a training pattern to help the host more accurately center the data capture point in the received data bits.	00h	R/W	Non-Volatile

**Table 6.26 Volatile Data Learning Register (VDLR)**

Bit	Name	Definition	Default	Read- /Write	Type
7:0	VDLP	Volatile Data Learning Pattern: Volatile copy of the NVDLP used to enable and deliver the Data Learning Pattern to the outputs. The VDLP may be changed by the host during system operation.	Takes the value of NVDLR during POR or Reset	R/W	Volatile

## 7. PROTECTION MODE

The device supports hardware and software write-protection mechanisms.

### 7.1 HARDWARE WRITE PROTECTION

The Write Protection (WP#) pin provides a hardware write protection method for BP3, BP2, BP1, BP0, SRWD, and QE in the Status Register. Refer to the section 6.1 STATUS REGISTER.

Write inhibit voltage ( $V_{WI}$ ) is specified in the section 9.8 POWER-UP AND POWER-DOWN. All write sequence will be ignored when  $V_{CC}$  drops to  $V_{WI}$ .

**Table 7.1 Hardware Write Protection on Status Register**

SRWD	WP#	Status Register
0	Low	Writable
1	Low	Protected
0	High	Writable
1	High	Writable

**Note:** Before the execution of any program, erase or Write Status/Function Register instruction, the Write Enable Latch (WEL) bit must be enabled by executing a Write Enable (WREN) instruction. If the WEL bit is not enabled, the program, erase or write register instruction will be ignored.

### 7.2 SOFTWARE WRITE PROTECTION

The device also provides two kinds of software write protection feature. One is Block Protection by Block Protection bits (BP3, BP2, BP1, BP0) and another is Advanced Sector/Block Protection (ASP). When Block Protection is enabled (i.e., any BP3-0 are set to "1"), Advanced Sector/Block Protection (ASP) can still be used to protect sectors/blocks not protected by the Block Protection scheme. In the case that both ASP and Block Protection are used on the same sector/block the logical OR of ASP and Block Protection related to the sector/block is used.

**Warning:** ASP and Block Protection should not be used concurrently. Use one or the other, but not both.

#### 7.2.1 BLOCK PROTECTION BITS

The device provides a software write protection feature. The Block Protection bits (BP3, BP2, BP1, BP0) allow part or the whole memory area to be write-protected. For details, see 6.1 Status Register.

### 7.2.2 ADVANCED SECTOR/BLOCK PROTECTION (ASP)

There are two ways to implement software Advanced Sector/Block Protection on this device: Password Protection method or Persistent Protection methods. Through these two protection methods, user can disable or enable the programming or erasing operation to any or all blocks including 32 top 4K sectors or 32 bottom 4K sectors when block size is 64KB. When block size is uniform 256KB, there will be no 32 4KB parameter blocks.

The Figure 7.1 shows an overview of these methods

Every main flash array block/top sector/bottom sector has a non-volatile (PPB) and a volatile (DYB) protection bit associated with it. When either bit is 0, the sector is protected from program and erase operations.

The PPB bits are protected from program and erase when the PPB Lock bit is “0”. The PPB bits are erased so that all main flash array sectors are unprotected when shipped from factory.

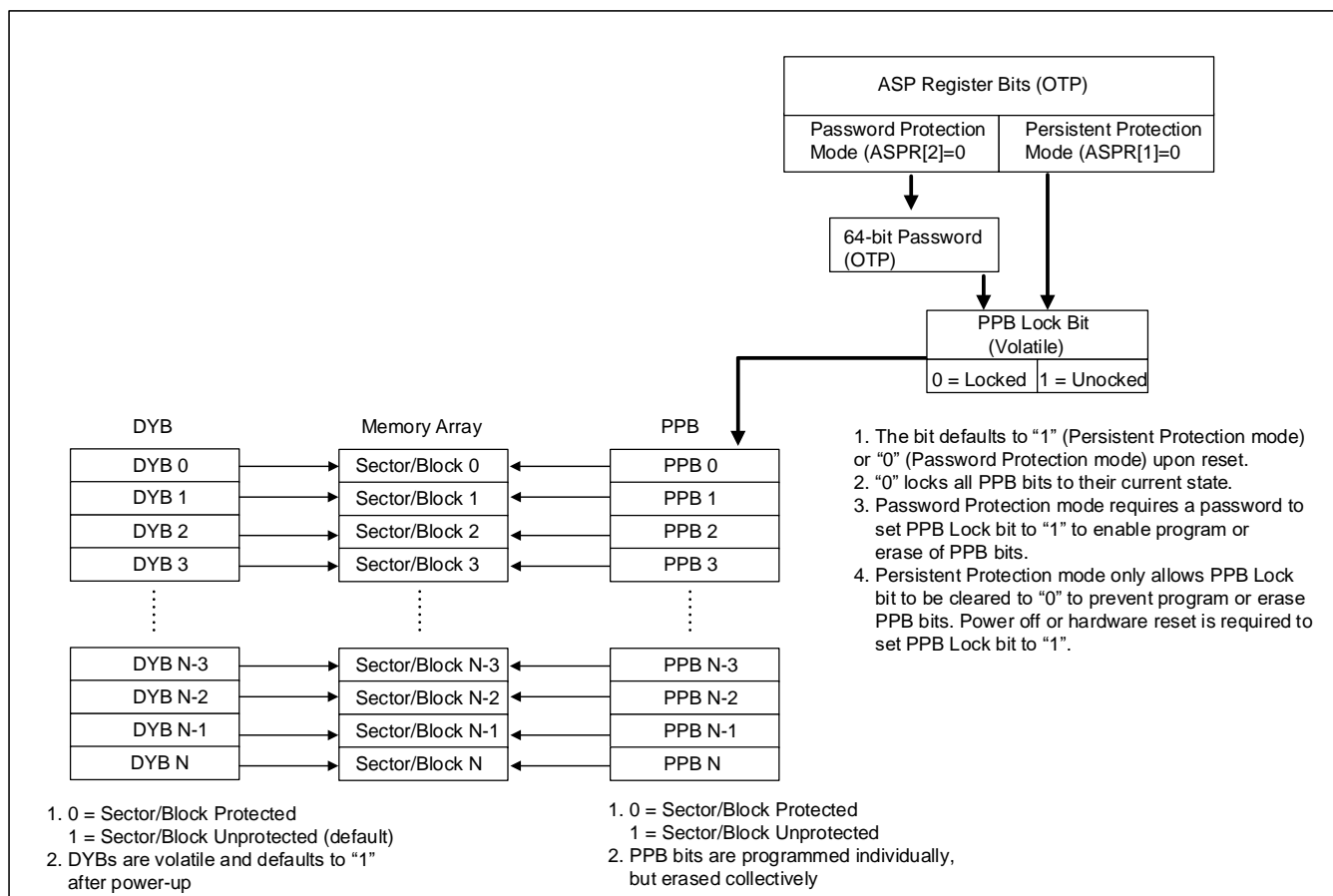
There are two methods for managing the state of the PPB Lock bit, Persistent Protection and Password Protection.

The Persistent Protection Mode sets the PPB Lock bit to “1” during power up or Hardware Reset so that the PPB bits are unprotected. There is a WRPLB command to clear the PPB Lock bit to “0” to protect the PPB bits. There is no command in the Persistent Protection method to set the PPB Lock bit therefore the PPB Lock bit will remain at “0” until the next power up or Hardware Reset. The Persistent Protection method allows boot code the option of changing sector protection by programming or erasing the PPB, then protecting the PPB from further change for the remainder of normal system operation by clearing the PPB Lock bit. This is sometimes called Boot-code controlled sector protection.

The Password Protection Mode requires use of a password to control PPB protection. In the Password Protection Mode, the PPB Lock bit is cleared to “0” during power up or Hardware Reset to protect the PPB bits. A 64-bit password may be permanently programmed and hidden for the Password Protection Mode. The UNPWD command can be used to provide a password for comparison with the hidden password. If the password matches the PPB Lock bit is set to “1” to unprotect the PPB. The WRPLB command can be used to clear the PPB Lock bit to “0”. After clearing the PPB Lock bit to “0”, the UNPWD command can be used again to unprotect the PPB.

The selection of the PPB Lock bit management method is made by programming OTP bits in the ASP Register so as to permanently select the method used.



**Figure 7.1 Advanced Sector/Block Protection**

**Notes:**

- 256M (64KB Block):**  $N = 541 = 32 \text{ (32xTop 4KB sectors or 32xBottom 4KB sectors)} + 510 \text{ (510x64KB blocks)} - 1$
- 256M (256KB Block):**  $N = 127 = 0 + 128 \text{ (128x256KB blocks)} - 1$

**Table 7.2 PPB/DYB and Sector/Block mapping when Block Size is 64KB (TBPARM = 1)**

Memory Density	PPB Group	DYB Group	Block No. (64Kbyte)	Sector No.	Sector Size (Kbyte)	Address Range
<b>256Mb</b>	PPB 0	DYB 0	Block 0	Sector 0	4	000000h - 000FFFh
	:	:		:	:	:
	:	:		:	:	:
	PPB 15	DYB 15	Block 1	Sector 15	4	00F000h - 00FFFFh
	PPB 16	DYB 16		Sector 16	4	010000h - 010FFFh
	:	:		:	:	:
	:	:	Block 2	Sector 31	4	01F000h - 01FFFFh
	PPB 31	DYB 31		Sector 32	4	020000h - 020FFFh
	PPB 32	DYB 32		:	:	:
	:	:	Block 254	:	:	:
	:	:		Sector 4064	4	FE0000h - FE0FFFh
	:	:		:	:	:
	PPB 284	DYB 284	Block 255	Sector 4079	4	FEF000h - FFFFFFFh
	PPB 285	DYB 285		Sector 4080	4	FF0000h - FF0FFFh
	:	:		:	:	:
	:	:	Block 510	Sector 4095	4	FFF000h - FFFFFFFh
	:	:		:	:	:
	:	:		:	:	:
	PPB 540	DYB 540	Block 511	Sector 8160	4	1FE0000h - 1FE0FFFh
	:	:		:	:	:
	:	:		:	:	:
	:	:	Block 511	Sector 8175	4	1FEF000h - 1FEFFFFh
	PPB 541	DYB 541		Sector 8176	4	1FF0000h - 1FF0FFFh
	:	:		:	:	:
	:	:		:	:	:
	:	:		Sector 8191	4	1FFF000h - 1FFFFFFh

**Table 7.3 PPB/DYB and Sector/Block mapping when Block Size is 64KB (TBPARM = 0)**

Memory Density	PPB Group	DYB Group	Block No. (64Kbyte)	Sector No.	Sector Size (Kbyte)	Address Range
<b>256Mb</b>	PPB 0	DYB 0	Block 0	Sector 0	4	000000h - 000FFFh
				:	:	:
				:	:	:
	PPB 1	DYB 1	Block 1	Sector 15	4	00F000h - 00FFFFh
				Sector 16	4	010000h - 010FFFh
				:	:	:
				:	:	:
				Sector 31	4	01F000h - 01FFFFh
	PPB 2	DYB 2	Block 2	Sector 32	4	020000h - 020FFFh
				:	:	:
				:	:	:
				Sector 47	4	02F000h - 02FFFFh
	:	:	:	:	:	:
	PPB 254	DYB 254	Block 254	Sector 4064	4	FE0000h - FE0FFFh
				:	:	:
				:	:	:
				Sector 4079	4	FEF000h - FFFFFFFh
	PPB 255	DYB 255	Block 255	Sector 4080	4	FF0000h - FF0FFFh
				:	:	:
				:	:	:
				Sector 4095	4	FFF000h - FFFFFFFh
	:	:	:	:	:	:
	PPB 510	DYB 510	Block 510	Sector 8160	4	1FE0000h - 1FE0FFFh
	:	:		:	:	:
	:	:		:	:	:
	PPB 525	DYB 525	Block 511	Sector 8175	4	1FEF000h - 1FEFFFFh
	PPB 525	DYB 526		Sector 8176	4	1FF0000h - 1FF0FFFh
	:	:		:	:	:
	:	:		:	:	:
	PPB 541	DYB 541		Sector 8191	4	1FFF000h - 1FFFFFFh

**Table 7.4 PPB/DYB and Sector/Block mapping when Block Size is 256KB (TBPARM bit is Reserved)**

Memory Density	PPB Group	DYB Group	Block No. (256Kbyte)	Sector No.	Sector Size (Kbyte)	Address Range
256Mb	PPB 0	DYB 0	Block 0	Sector 0	4	000000h - 000FFFh
				:	:	:
				:	:	:
				Sector 15	4	00F000h - 00FFFFh
				Sector 16	4	010000h - 010FFFh
				:	:	:
				:	:	:
				Sector 31	4	01F000h - 01FFFFh
				Sector 32	4	020000h - 020FFFh
				:	:	:
				:	:	:
				Sector 47	4	02F000h - 02FFFFh
				Sector 48	4	030000h - 030FFFh
				:	:	:
				:	:	:
				Sector 63	4	03F000h - 03FFFFh
	:	:	:	:	:	:
	PPB 63	DYB 63	Block 63	Sector 4032	4	FC0000h - FC0FFFh
				:	:	:
				:	:	:
				Sector 4047	4	FCF000h - FCFFFFh
				:	:	:
				Sector 4080	4	FF0000h - FF0FFFh
				:	:	:
				:	:	:
				Sector 4095	4	FFF000h - FFFFFFh
	:	:	:	:	:	:
	PPB 127	DYB 127	Block 127	Sector 8128	4	1FC0000h - 1FC0FFFh
				:	:	:
				:	:	:
				Sector 8143	4	1FCF000h - 1FCFFFFh
				Sector 8144	4	1FD0000h - 1FD0FFFh
				:	:	:
				:	:	:
				Sector 8159	4	1FDF000h - 1FDFFFFh
				Sector 8160	4	1FE0000h - 1FE0FFFh
				:	:	:
				:	:	:
				Sector 8175	4	1FEF000h - 1FEFFFFh
				Sector 8176	4	1FF0000h - 1FF0FFFh
				:	:	:
				:	:	:
				Sector 8191	4	1FFF000h - 1FFFFFFh

### **Persistent Protection Bits (PPBs)**

The Persistent Protection Bits (PPBs) are unique for each sector/block and non-volatile (refer to Figure 7.1, Table 7.2, and Table 7.3). It is programmed individually but must be erased as a group, similar to the way individual words may be programmed in the main array but an entire sector/block must be erased at the same time. The PPBs have the same endurance as the Flash memory. Preprogramming and verification prior to erasure are handled by the device, and therefore do not require system monitoring. Programming a PPB bit requires the typical page programming time. Erasing all the PPBs requires typical sector erase time. During PPB bit programming and PPB bit erasing, status is available by reading the Status Register or Extended Read Register. Reading of a PPB bit requires the initial access time of the device.

#### **Notes:**

1. Each PPB is individually programmed to “0” and all are erased to “1” in parallel.
2. The PPB Lock bit must be cleared first before changing the status of a PPB.
3. While programming PPB, array data cannot be read from any sectors/blocks.
4. When reading the PPB of the desired sector/block the address should be location zero within the sector/block. The high order address bits not used must be zero.
5. There are no means for individually erasing a specific PPB and no specific sector/block address is required for this operation.
6. The state of the PPB for a given sector/block can be verified by using a PPB Read command.
7. When the parts are first shipped, the PPBs are cleared (erased to “1”).

### **Dynamic Protection Bits (DYBs)**

Dynamic Protection Bits (DYBs) are volatile and unique for each sector/block and can be individually modified. DYBs only control the protection for unprotected sectors/blocks that have their PPBs cleared (erased to “1”). By issuing the Write DYB command, the DYBs are cleared to “0” or set to “1”, thus placing each sector/block in the protected or unprotected state respectively. This feature allows software to easily protect sectors/blocks against inadvertent changes, yet does not prevent the easy removal of protection when changes are needed. The DYBs can be set or cleared as often as needed as they are volatile bits.

### **Persistent Protection Bit (PPB) Lock Bit**

The PPB Lock bit is a volatile bit for protecting all PPB bits. When cleared to “0”, it locks all PPBs and when set to “1”, it allows the PPBs to be changed. . If the PPB Lock bit is “0”, the PPB Program or Erase command does not execute and fails without programming or erasing the PPB.

In Persistent Protection mode, the PPB Lock bit is set to “1” during power up or Hardware Reset. When cleared to “0”, no software command sequence can set the PPB Lock bit to “1”, only another Hardware Reset or power-up can set the PPB Lock bit.

In the Password Protection mode, the PPB Lock bit is cleared to “0” during power up or a Hardware Reset during power up or a Hardware Reset. The PPB Lock bit can only be set to “1” by the Password Unlock command.

The PPB Lock bit must be cleared to “0” only after all PPBs are configured to the desired settings.

### **Sector/Block Protection States Summary**

Each sector in specific blocks and each of all other blocks except for the specific blocks can be in one of the following protection states:

- Unlocked – The sector/block is unprotected and protection can be changed by a simple command. The protection state defaults to unprotected after a power cycle, software reset, or hardware reset.
- Dynamically Locked – A sector/block is protected and protection can be changed by a simple command. The protection state is not saved across a power cycle.
- Persistently Locked – A sector/block is protected and protection can only be changed if the PPB Lock bit is set to “1”. The protection state is non-volatile and saved across a power cycle or reset. Changing the protection state requires programming and or erase of the PPB bits.

Table 7.4 contains all possible combinations of the DYB, PPB, and PPB Lock bit relating to the status of the sector/block. In summary, if the PPB Lock bit is locked (cleared to “0”), no changes to the PPBs are allowed. The PPB Lock bit can only be unlocked (set to “1”) through a Hardware Reset or power cycle.

**Table 7.5 Sector/Block Protection States**

		Protection Bit values		Assigned Sector/Block State
		PPB	DYB	
“0” = Locked or Protected “1” = Unlocked or Unprotected		1	1	Unprotected
		1	0	Protected
		0	1	Protected
		0	0	Protected
PPB Lock Bit	1	Changeable	Changeable	
	0	NOT changeable	Changeable	

### **Persistent Protection Mode**

The Persistent Protection Mode sets the PPB Lock bit to “1” during power up or Hardware Reset so that the PPB bits are unprotected by a device Hardware Reset. Software Reset does not affect the PPB Lock bit. The WRPLB command can clear the PPB Lock bit to “0” to protect the PPB. There is no command to set the PPB Lock bit therefore the PPB Lock bit will remain at “0” until the next power up or Hardware Reset.

### **Password Protection Mode**

The Password Protection Mode allows an even higher level of security than the Persistent Protection Mode by requiring a 64-bit password for unlocking the device PPB Lock bit. In addition to this password requirement, after power up or Hardware Reset, the PPB Lock bit is cleared to “0” to maintain the password mode of operation. Successful execution of the Unlock Password command by entering the entire password sets the PPB Lock bit to “1”, allowing for sector/block PPBs modifications.

#### **Notes:**

1. The password is all “1”s when shipped from Factory. It is located in its own memory space and is accessible through the use of the Program Password and Read Password commands.
2. Once the Password is programmed and verified, the Password Protection Mode Lock Bit (ASPR[2]=0) in ASP Register must be programmed in order to prevent reading or modifying the password. After the Password Protection Mode Lock Bit is programmed, all further Program and Read commands to the password region are disabled and these commands are ignored so that there is no means to verify what the password is. Password verification is only allowed before selecting the Password Protection Mode.
3. The Program Password Command is only capable of programming “0”s. Programming a “1” after a cell is programmed as a “0” results in the cell left as a “0” with no programming error.
4. All 64-bit password combinations are valid as a password.
5. The Protection Mode Lock Bits in ASP Register are not erasable because they are OTP.
6. The exact password must be entered in order for the unlocking function to occur. If the password provided by Unlock Password command does not match the hidden internal password, the unlock operation fails in the same manner as a programming operation on a protected sector/block. The P\_ERR and PROT\_E are set to 1 and the PPB Lock bit remains cleared to 0. In this case it is a failure to change the state of the PPB Lock bit because it is still protected by the lack of a valid password.
7. The Unlock Password command cannot be accepted any faster than once every  $100\mu\text{s} \pm 20\mu\text{s}$ . This makes it take an unreasonably long time (58 million years) for a hacker to run through all the 64-bit combinations in an attempt to correctly match a password. The Read Status Register command or the Read Extended Read Register may be used to read the WIP bit to determine when the device has completed the Unlock Password command or is ready to accept a new password command. When a valid password is provided the Unlock Password command does not insert the  $100\mu\text{s}$  delay before returning the WIP bit to zero.
8. If the password is lost after selecting the Password Protection Mode, there is no way to set the PPB Lock bit.

### 7.3 RPMC OPERATION

The device supports RPMC (Replay Protection Monotonic Counter) function as an additional security, which could detect Hardware attack.

The device is equipped with four 32-bit Monotonic Counters and can be addressed by the 8-bit Counter Address. These monotonic counters are used by the host controllers to ensure the physical authenticity of the flash devices.

RPMC operation is based on the HMAC-SHA-256 cryptographic algorithm. HMAC-SHA-256 is a type of keyed hash algorithm that is constructed from the SHA-256 hash function and used as a Hash-based Message Authentication Code (HMAC). The HMAC process mixes a secret key with the message data, hashes the result with the hash function, mixes that hash value with the secret key again, and then applies the hash function a second time. The output of hash function is 256 bits in length.

An HMAC can be used to determine whether a message sent over an insecure channel has been tempered with, provided that the sender and receiver share a secret key. The sender computes the hash value for the original data and sends both the original data and hash value as a single message. The receiver recalculates the hash value on the received message and checks that computed HMAC matches the transmitted HMAC.

Any change to the data or hash value results in a mismatch, because knowledge of the secret key is required to change the message and reproduce the correct hash value. Therefore, if the original and computed hash values match, the message is authenticated.

A Reset (Hardware reset, Software Reset, or Power On Reset) operation will cause the part to get fully reset. All volatile memory based resources will be cleared. **Update HMAC Key ree register command must be reissued after a reset.**

While in a deep power down state, OP1, OP2 commands are ignored until the device comes out of deep power down mode.

The device is an MCP (RPMC + SPI NOR Flash), so the device can support concurrent operation of OP1 and subsequent flash transactions.



## 8. DEVICE OPERATION

### 8.1 COMMAND OVERVIEW

The device utilizes an 8-bit instruction register. Refer to Table 8.4. Instruction Set for details on instructions and instruction codes. All instructions, addresses, and data are shifted in with the most significant bit (MSB) first on Serial Data Input (SI) or Serial Data IOs (IO0, IO1, IO2, IO3). The input data on SI or IOs is latched on the rising edge of Serial Clock (SCK) for normal mode and both of rising and falling edges for DTR mode after Chip Enable (CE#) is driven low ( $V_{IL}$ ). Every instruction sequence starts with a one-byte instruction code and is followed by address bytes and/or dummy cycles (configurable) and/or data bytes, depending on the type of instruction. CE# must be driven high ( $V_{IH}$ ) after the last bit of the instruction sequence has been shifted in to end the operation.

Commands are structured as follows:

- Each command begins with a byte (eight bits) instruction.
- The instruction may be stand alone or may be followed by address bits to select a location within one of several address spaces in the device. The address may be either a 24-bit or 32-bit byte boundary address.
- The SPI interface with Multiple IO provides the option for each transfer of address and data information to be done one, two, or four bits in parallel. This enables a tradeoff between the number of signal connections (IO bus width) and the speed of information transfer. If the host system can support a two or four bit wide IO bus the memory performance can be increased by using the instructions that provide parallel two bit (dual) or parallel four bit (quad) transfers.
- The width of all transfers following the instruction are determined by the instruction sent.
- All single bit or parallel bit groups are transferred in most to least significant bit order.
- Some instructions send Mode Bits following the address to indicate that the next command will be of the same type with an implied, rather than an explicit, instruction. The next command thus does not provide an instruction byte, only a new address and mode bits. This reduces the time needed to send each command when the same command type is repeated in a sequence of commands.
- The address or Mode Bits may be followed by Dummy Cycles before read data is returned to the host.
- Dummy Cycles may be zero to several SCK cycles. In fact, Mode Bits will be counted as a part of Dummy Cycles.
- All instruction, address, Mode, and data information is transferred in byte granularity. Addresses are shifted into the device with the Most Significant Byte first. All data is transferred with the lowest address byte sent first. Following bytes of data are sent in lowest to highest byte address order i.e. the byte address increments.
- All attempts to read the flash memory array during a program, erase, or a write cycle (embedded operations) are ignored. The embedded operation will continue to execute without any affect. A very limited set of commands are accepted during an embedded operation. These are discussed in the individual command descriptions. While a program, erase, or write operation is in progress, it is recommended to check that the Write In Progress (WIP) bit is "0" before issuing most commands to the device, to ensure the new command can be accepted.
- Depending on the command, the time for execution varies. A command to read status information from an executing command is available to determine when the command completes execution and whether the command was successful.
- Following are some general signal relationship descriptions to keep in mind.
  - The host always controls the Chip Enable (CE#), Serial Clock (SCK), and Serial Input (SI) - SI for single bit wide transfers. The memory drives Serial Output (SO) for single bit read transfers. The host and memory alternately drive the IO0-IO3 signals during Dual and Quad transfers.
  - All commands begin with the host selecting the memory by driving CE# low before the first rising edge of SCK. CE# is kept low throughout a command and when CE# is returned high the command ends.

Generally, CE# remains low for 8-bit transfer multiples to transfer byte granularity information. All commands will not be accepted if CE# is returned high not at an 8-bit boundary.

## 8.2 COMMAND SET SUMMARY

### Extended Addressing

To accommodate addressing above 128Mb (24-bit) for read, erase, write and program operations, there are three options:

1. New 4-byte (32-bit) address instructions.

New 4-byte instructions are valid in both 4-byte addressing mode and 3-byte addressing mode. See Table 8.2.

2. 4-byte addressing mode with legacy 3-byte address instructions.

There are 2 ways to enter into 4-byte addressing mode.

- The device can be set to 4-byte addressing mode by setting EXTADD bit = “1” in the Bank Address Register (BAR[7])
- If EXTADD bit is set to “0” (3-byte addressing mode), also the device can be switched to 4-byte addressing mode by issuing Enter 4-byte Address Mode Command See Table 8.3.
- BA24 bit is ignored in 4-byte addressing mode

3. 3-byte addressing mode with the 3-byte address instructions:

The device can be set to 3-byte addressing mode by setting EXTADD bit = “0” in the Bank Address Register (BAR[7])

- BA24 selects individual bank (128Mb) in 3-byte addressing mode.
- 3-byte address selects an address within the bank (128Mb)
- The host system writes the Bank Address Register to access beyond the first 128Mbit of memory.
- Bank Address Register bits are volatile.
  - On power up, the default is Bank0 (the lowest address 16 Mbytes).
- For Read, the device will continuously transfer out data until the end of the array.
  - There is no bank to bank delay.
  - The Bank Address Register is not updated.
  - The Bank Address Register value is used only for the initial address of an access.

**Table 8.1 Bank Address Map**

Bank Address Register Bit 0	Bank	Memory Array Address Range	
0	0	00000000h	00FFFFFFh
1	1	01000000h	01FFFFFFh

**Table 8.2 New Instruction Set with 4-byte address**

Instruction Name	Operation	Code	Address Mode
4NORD	4-byte Address Normal Read Mode	13h	4-byte Address
4FRD	4-byte Address Fast Read Mode	0Ch	4-byte Address
4FRDIO	4-byte Address Fast Read Dual I/O	BCh	4-byte Address
4FRDO	4-byte Address Fast Read Dual Output	3Ch	4-byte Address
4FRQIO	4-byte Address Fast Read Quad I/O	ECh	4-byte Address
4FRQO	4-byte Address Fast Read Quad Output	6Ch	4-byte Address
4FRDTR	4-byte Address Fast Read DTR Mode	0Eh	4-byte Address
4FRDDTR	4-byte Address Fast Read Dual I/O DTR	BEh	4-byte Address
4FRQDTR	4-byte Address Fast Read Quad I/O DTR	EEh	4-byte Address
4PP	4-byte Address Serial Input Page Program	12h	4-byte Address
4PPQ	4-byte Address Quad Input Page Program	34h/3Eh	4-byte Address
4SER	4-byte Address Sector Erase	21h	4-byte Address
4BER32 (32KB)	4-byte Address Block Erase 32KB	5Ch	4-byte Address
4BER64/256 (64/256KB)	4-byte Address Block Erase 64KB	DCh	4-byte Address
4SECUNLOCK	4-byte Address Sector Unlock	25h	4-byte Address
4RDDYB	4-byte Address Read DYB	E0	4-byte Address
4WRDYB	4-byte Address Write DYB	E1	4-byte Address
4RDPPB	4-byte Address Read PPB	E2	4-byte Address
4PGPPB	4-byte Address Program PPB	E3	4-byte Address

**Table 8.3 Instruction Set with 3-byte or 4-byte address according to EXTADD Bit setting**

Instruction Name	Operation	Code	Address Mode	
			EXTADD (BAR[7] = 1)	EXTADD (BAR[7]) = 0
NORD	Normal Read Mode	03h	4-byte Address	3-byte Address
FRD	Fast Read Mode	0Bh	4-byte Address	3-byte Address
FRDIO	Fast Read Dual I/O	BBh	4-byte Address	3-byte Address
FRDO	Fast Read Dual Output	3Bh	4-byte Address	3-byte Address
FRQIO	Fast Read Quad I/O	EBh	4-byte Address	3-byte Address
FRQO	Fast Read Quad Output	6Bh	4-byte Address	3-byte Address
FRDTR	Fast Read DTR Mode	0Dh	4-byte Address	3-byte Address
FRDDTR	Fast Read Dual I/O DTR	BDh	4-byte Address	3-byte Address
FRQDTR	Fast Read Quad I/O DTR	EDh	4-byte Address	3-byte Address
PP	Serial Input Page Program	02h	4-byte Address	3-byte Address
PPQ	Quad Input Page Program	32h/38h	4-byte Address	3-byte Address
SER	Sector Erase	D7h/20h	4-byte Address	3-byte Address
BER32 (32KB)	Block Erase 32KB	52h	4-byte Address	3-byte Address
BER64 (64KB/256KB)	Block Erase 64KB/256KB	D8h	4-byte Address	3-byte Address
SECUNLOCK	Sector Unlock	26h	4-byte Address	3-byte Address
RDDYB	Read DYB	FA	4-byte Address	3-byte Address
WRDYB	Write DYB	FB	4-byte Address	3-byte Address
RDPPB	Read PPB	FC	4-byte Address	3-byte Address
PGPPB	Program PPB	FD	4-byte Address	3-byte Address

**Table 8.4 NOR Flash Instruction Set**

Instruction Name	Operation	Mode	Byte0	Byte1	Byte2	Byte3	Byte4	Byte5	Byte6
NORD	Normal Read Mode (3-byte Address)	SPI	03h	A <23:16>	A <15:8>	A <7:0>	Data out		
NORD	Normal Read Mode (4-byte Address)	SPI	03h	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Data out	
4NORD	4-byte Address Normal Read Mode	SPI	13h	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Data out	
FRD	Fast Read Mode (3-byte Address)	SPI QPI	0Bh	A <23:16>	A <15:8>	A <7:0>	Dummy <sup>(1)</sup> Byte	Data out	
FRD	Fast Read Mode (4-byte Address)	SPI QPI	0Bh	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Dummy <sup>(1)</sup> Byte	Data out
4FRD	4-byte Address Fast Read Mode	SPI QPI	0Ch	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Dummy <sup>(1)</sup> Byte	Data out
FRDIO	Fast Read Dual I/O (3-byte Address)	SPI	BBh	A <23:16> Dual	A <15:8> Dual	A <7:0> Dual	AXh <sup>(1),(2)</sup> Dual	Dual Data out	
FRDIO	Fast Read Dual I/O (4-byte Address)	SPI	BBh	A <31::24>	A <23:16> Dual	A <15:8> Dual	A <7:0> Dual	AXh <sup>(1),(2)</sup> Dual	Dual Data out
4FRDIO	4-byte Address Fast Read Dual I/O	SPI	BCh	A <31::24>	A <23:16> Dual	A <15:8> Dual	A <7:0> Dual	AXh <sup>(1),(2)</sup> Dual	Dual Data out
FRDO	Fast Read Dual Output (3-byte Address)	SPI	3Bh	A <23:16>	A <15:8>	A <7:0>	Dummy <sup>(1)</sup> Byte	Dual Data out	
FRDO	Fast Read Dual Output (4-byte Address)	SPI	3Bh	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Dummy <sup>(1)</sup> Byte	Dual Data out
4FRDO	4-byte Address Fast Read Dual Output	SPI	3Ch	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Dummy <sup>(1)</sup> Byte	Dual Data out
FRQIO	Fast Read Quad I/O (3-byte Address)	SPI QPI	EBh	A <23:16> Quad	A <15:8> Quad	A <7:0> Quad	AXh <sup>(1),(2)</sup> Quad	Quad Data out	
FRQIO	Fast Read Quad I/O (4-byte Address)	SPI QPI	EBh	A <31::24> Quad	A <23:16> Quad	A <15:8> Quad	A <7:0> Quad	AXh <sup>(1),(2)</sup> Quad	Quad Data out
4FRQIO	4-byte Address Fast Read Quad I/O	SPI QPI	ECh	A <31::24> Quad	A <23:16> Quad	A <15:8> Quad	A <7:0> Quad	AXh <sup>(1),(2)</sup> Quad	Quad Data out
FRQO	Fast Read Quad Output (3-byte Address)	SPI	6Bh	A <23:16>	A <15:8>	A <7:0>	Dummy <sup>(1)</sup> Byte	Quad Data out	
FRQO	Fast Read Quad Output (4-byte Address)	SPI	6Bh	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Dummy <sup>(1)</sup> Byte	Quad Data out
4FRQO	4-byte Address Fast Read Quad Output	SPI	6Ch	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Dummy <sup>(1)</sup> Byte	Quad Data out

Instruction Name	Operation	Mode	Byte0	Byte1	Byte2	Byte3	Byte4	Byte5	Byte6
FRDTR	Fast Read DTR Mode (3-byte Address)	SPI QPI	0Dh	A <23:16>	A <15:8>	A <7:0>	Dummy <sup>(1)</sup> Byte	Dual Data out	
FRDTR	Fast Read DTR Mode (4-byte Address)	SPI QPI	0Dh	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Dummy <sup>(1)</sup> Byte	Dual Data out
4FRDTR	4-byte Address Fast Read DTR Mode	SPI QPI	0Eh	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Dummy <sup>(1)</sup> Byte	Dual Data out
FRDDTR	Fast Read Dual I/O DTR (3-byte Address)	SPI	BDh	A <23:16> Dual	A <15:8> Dual	A <7:0> Dual	AXh <sup>(1), (2)</sup> Dual	Dual Data out	
FRDDTR	Fast Read Dual I/O DTR (4-byte Address)	SPI	BDh	A <31::24>	A <23:16> Dual	A <15:8> Dual	A <7:0> Dual	AXh <sup>(1), (2)</sup> Dual	Dual Data out
4FRDDTR	4-byte Address Fast Read Dual I/O DTR	SPI	BEh	A <31::24>	A <23:16> Dual	A <15:8> Dual	A <7:0> Dual	AXh <sup>(1), (2)</sup> Dual	Dual Data out
FRQDTR	Fast Read Quad I/O DTR (3-byte Address)	SPI QPI	EDh	A <23:16>	A <15:8>	A <7:0>	AXh <sup>(1), (2)</sup> Quad	Quad Data out	
FRQDTR	Fast Read Quad I/O DTR (4-byte Address)	SPI QPI	EDh	A <31::24>	A <23:16>	A <15:8>	A <7:0>	AXh <sup>(1), (2)</sup> Quad	Quad Data out
4FRQDTR	4-byte Address Fast Read Quad I/O DTR	SPI QPI	EEh	A <31::24>	A <23:16>	A <15:8>	A <7:0>	AXh <sup>(1), (2)</sup> Quad	Quad Data out
PP	Input Page Program (3-byte Address)	SPI QPI	02h	A <23:16>	A <15:8>	A <7:0>	PD (256byte)		
PP	Input Page Program (4-byte Address)	SPI QPI	02h	A <31::24>	A <23:16>	A <15:8>	A <7:0>	PD (256byte)	
4PP	4-byte Address Input Page Program	SPI QPI	12h	A <31::24>	A <23:16>	A <15:8>	A <7:0>	PD (256byte)	
PPQ	Quad Input Page Program (3-byte Address)	SPI	32h/38h	A <23:16>	A <15:8>	A <7:0>	Quad PD (256byte)		
PPQ	Quad Input Page Program (4-byte Address)	SPI	32h/38h	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Quad PD (256byte)	
4PPQ	4-byte Address Quad Input Page Program	SPI	34h/3Eh	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Quad PD (256byte)	
SER	Sector Erase (3-byte Address)	SPI QPI	D7h/20h	A <23:16>	A <15:8>	A <7:0>			
SER	Sector Erase (4-byte Address)	SPI QPI	D7h/20h	A <31::24>	A <23:16>	A <15:8>	A <7:0>		
4SER	4-byte Address Sector Erase	SPI QPI	21h	A <31::24>	A <23:16>	A <15:8>	A <7:0>		

Instruction Name	Operation	Mode	Byte0	Byte1	Byte2	Byte3	Byte4	Byte5	Byte6
BER32 (32KB)	Block Erase 32Kbyte (3-byte Address)	SPI QPI	52h	A <23:16>	A <15:8>	A <7:0>			
BER32 (32KB)	Block Erase 32Kbyte (4-byte Address)	SPI QPI	52h	A <31::24>	A <23:16>	A <15:8>	A <7:0>		
4BER32 (32KB)	4-byte Address Block Erase 32Kbyte	SPI QPI	5Ch	A <31::24>	A <23:16>	A <15:8>	A <7:0>		
BER64 (64KB)	Block Erase 64Kbyte (3-byte Address)	SPI QPI	D8h	A <23:16>	A <15:8>	A <7:0>			
BER64 (64KB)	Block Erase 64Kbyte (4-byte Address)	SPI QPI	D8h	A <31::24>	A <23:16>	A <15:8>	A <7:0>		
4BER64 (64KB)	4-byte Address Block Erase 64Kbyte	SPI QPI	DCh	A <31::24>	A <23:16>	A <15:8>	A <7:0>		
CER	Chip Erase	SPI QPI	C7h/60h						
WREN	Write Enable	SPI QPI	06h						
WRDI	Write Disable	SPI QPI	04h						
RDSR	Read Status Register	SPI QPI	05h	Data out					
WRSR	Write Status Register	SPI QPI	01h	Data in					
RDFR	Read Function Register	SPI QPI	48h	Data out					
WRFR	Write Function Register	SPI QPI	42h	Data in					
QPIEN	Enter QPI mode	SPI	35h						
QPIDI	Exit QPI mode	QPI	F5h						
PERSUS	Suspend during program/erase	SPI QPI	75h/B0h						
PERRSM	Resume program/erase	SPI QPI	7Ah/30h						
DP	Deep Power Down	SPI QPI	B9h						
RDID, RDPD	Read ID / Release Power Down	SPI QPI	ABh	XXh <sup>(3)</sup>	XXh <sup>(3)</sup>	XXh <sup>(3)</sup>	ID7-ID0		
SRPNV	Set Read Parameters (Non-Volatile)	SPI QPI	65h	Data in					
SRPV	Set Read Parameters (Volatile)	SPI QPI	C0h/63h	Data in					
SERPNV	Set Extended Read Parameters (Non-Volatile)	SPI QPI	85h	Data in					

Instruction Name	Operation	Mode	Byte0	Byte1	Byte2	Byte3	Byte4	Byte5	Byte6
SERPv	Set Extended Read Parameters (Volatile)	SPI QPI	83h	Data in					
RDRP	Read Read Parameters (Volatile)	SPI QPI	61h	Data out					
RDERP	Read Extended Read Parameters (Volatile)	SPI QPI	81h	Data out					
CLERP	Clear Extended Read Register	SPI QPI	82h						
RDJDID	Read JEDEC ID Command	SPI QPI	9Fh	MF7-MF0	ID15-ID8	ID7-ID0			
RDMDID	Read Manufacturer & Device ID	SPI QPI	90h	XXh <sup>(3)</sup>	XXh <sup>(3)</sup>	00h	MF7-MF0	ID7-ID0	
						01h	ID7-ID0	MF7-MF0	
RDJDIDQ	Read JEDEC ID QPI mode	QPI	AFh	MF7-MF0	ID15-ID8	ID7-ID0			
RDUID	Read Unique ID	SPI QPI	4Bh	A <sup>(4)</sup> <23:16>	A <sup>(4)</sup> <15:8>	A <sup>(4)</sup> <7:0>	Dummy Byte	Data out	
RDSFDP	SFDP Read	SPI QPI	5Ah	A <23:16>	A <15:8>	A <7:0>	Dummy Byte	Data out	
NOP	No Operation	SPI QPI	00h						
RSTEN	Software Reset Enable	SPI QPI	66h						
RST	Software Reset	SPI QPI	99h						
IRER	Erase Information Row	SPI QPI	64h	A <23:16>	A <15:8>	A <7:0>			
IRP	Program Information Row	SPI QPI	62h	A <23:16>	A <15:8>	A <7:0>	PD (256byte)		
IRRD	Read Information Row	SPI QPI	68h	A <23:16>	A <15:8>	A <7:0>	Dummy Byte	Data out	
SECUNLOCK	Sector Unlock (3-byte Address)	SPI QPI	26h	A <23:16>	A <15:8>	A <7:0>			
SECUNLOCK	Sector Unlock (4-byte Address)	SPI QPI	26h	A <31::24>	A <23:16>	A <15:8>	A <7:0>		
4SECUNLOCK	4-byte Address Sector Unlock	SPI QPI	25h	A <31::24>	A <23:16>	A <15:8>	A <7:0>		
SECLOCK	Sector Lock	SPI QPI	24h						
RDABR	Read AutoBoot Register	SPI	14h						
WRABR	Write AutoBoot Register	SPI QPI	15h	Data in 1	Data in 2	Data in 3	Data in 4		



Instruction Name	Operation	Mode	Byte0	Byte1	Byte2	Byte3	Byte4	Byte5	Byte6
RDBR	Read Bank Address Register (Volatile)	SPI QPI	16h/C8h	Data out					
WRBRV	Write Bank Address Register (Volatile)	SPI QPI	17h/C5h	Data in					
WRBRNV	Write Bank Address Register (Non-Volatile)	SPI QPI	18h	Data in					
EN4B	Enter 4-byte Address Mode	SPI QPI	B7h						
EX4B	Exit 4-byte Address Mode	SPI QPI	29h						
RDDYB	Read DYB (3-byte Address)	SPI	FAh	A <23:16>	A <15:8>	A <7:0>	Data out		
RDDYB	Read DYB (4-byte Address)	SPI	FAh	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Data out	
4RDDYB	4-byte Address Read DYB	SPI	E0h	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Data out	
WRDYB	Write DYB (3-byte Address)	SPI QPI	FBh	A <23:16>	A <15:8>	A <7:0>	Data in		
WRDYB	Write DYB (4-byte Address)	SPI QPI	FBh	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Data in	
4WRDYB	4-byte Address Write DYB	SPI QPI	E1h	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Data in	
RDPPB	Read PPB (3-byte Address)	SPI	FCh	A <23:16>	A <15:8>	A <7:0>	Data out		
RDPPB	Read PPB (4-byte Address)	SPI	FCh	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Data out	
4RDPPB	4-byte Address Read PPB	SPI	E2h	A <31::24>	A <23:16>	A <15:8>	A <7:0>	Data out	
PGPPB	Program PPB (Individually) (3-byte Address)	SPI QPI	FDh	A <23:16>	A <15:8>	A <7:0>			
PGPPB	Program PPB (Individually) (4-byte Address)	SPI QPI	FDh	A <31::24>	A <23:16>	A <15:8>	A <7:0>		
4PGPPB	4-byte Address Program PPB (Individually)	SPI QPI	E3h	A <31::24>	A <23:16>	A <15:8>	A <7:0>		
ERPPB	Erase PPB (as a group)	SPI QPI	E4h						
RDASP	Read ASP	SPI	2Bh	Data out (2 byte)					
PGASP	Program ASP	SPI QPI	2Fh	PD (2 byte)					
RDPLB	Read PPB Lock Bit	SPI	A7h	Data out					
WRPLB	Write PPB Lock Bit	SPI QPI	A6h						
SFRZ	Set FREEZE bit	SPI QPI	91h						

Instruction Name	Operation	Mode	Byte0	Byte1	Byte2	Byte3	Byte4	Byte5	Byte6
RDPWD	Read Password	SPI	E7h	Data out (8 byte)					
PGPWD	Program Password	SPI QPI	E8h	PD (8 byte)					
UNPWD	Unlock Password	SPI QPI	E9h	Data in (8 byte)					
GBLK	Set all DYB bits (Gang Sector/Block Lock)	SPI QPI	7Eh						
GBUN	Clear all DYB bits (Gang Sector/Block Unlock)	SPI QPI	98h						
RDDLDP	Read Data Learning Pattern	SPI QPI	41h	Data out					
PNVDLR	Program DLP Register (Non-Volatile)	SPI QPI	43h	Data in					
WRVDLR	Write DLP Register (Volatile)	SPI QPI	4Ah	Data in					

**Notes:**

1. The number of dummy cycles depends on the value setting in the Table 6.11 Read Dummy Cycles.
2. AXh has to be counted as a part of dummy cycles. X means “don’t care”.
3. XX means “don’t care”.
4. A<23:9> are “don’t care” and A<8:4> are always “0”.

**Table 8.5 RPMC Command**

OP1 Input Command (OP1) <sup>(1)</sup>						
Operation	Command (BYTE 0)	Cmd Type (Byte 1)	Counter Address (Byte 2)	Reserved <sup>(2)</sup> (Byte 3)	Data	Signature
Write Root Key Register	9Bh	00h	[7:0]	[7:0]	Byte 4 – 35 RootKey[255:0]	Byte 36 – 63 TruncatedSign[223:0]
Update HMAC Key Register	9Bh	01h	[7:0]	[7:0]	Byte 4 – 7 KeyData [31:0]	Byte 8 – 39 Signature [255:0]
Increment Monotonic Counter	9Bh	02h	[7:0]	[7:0]	Byte 4 – 7 CounterData [31:0]	Byte 8 – 39 Signature [255:0]
Request Monotonic Counter	9Bh	03h	[7:0]	[7:0]	Byte 4 – 15 Tag [95:0]	Byte 16 – 47 Signature [255:0]
Reserved	9Bh	04h ~ FFh	Reserved			
OP2 Output Command (OP2) <sup>(1)</sup>						
Operation	Command (BYTE 0)	Dummy (Byte 1)	Status	Data		Signature
Read RPMC Status/Data	96h	8 clock cycles	Byte 2 RPMC Status [7:0]	Byte 3-14 Tag [ 95:0]	Byte 15-18 Counter Data [31:0]	Byte 19 – 50 Signature [255:0]

**Notes:**

1. All RPMC operations are in SPI mode (1 input & 1 output). Each input/output byte requires 8 clocks.
2. The reserved field (Byte 3) for RPMC OP1 operation is “Don’t Care”
3. The host controller may terminate the Read RPMC Status/Data operation (96h) at any time without going through the entire data output (Byte 2 ~Byte 50) sequence
4. When BUSY=1, the device will output the RPMC\_Status [7:0] value only after 8 dummy cycles, repeatedly until CE# terminates the operation. The device will NOT output Tag, CounterData & Signature fields when BUSY=1. Once BUSY becomes 0, another OP2 command must be issued to read out the correct Tag, CounterData & Signature fields.

### 8.3 NORMAL READ OPERATION (NORD, 03h or 4NORD, 13h)

The Normal Read (NORD) instruction is used to read memory contents at a maximum frequency of 80MHz.

- 03h (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- 03h (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- 13h is followed by a 4-byte address (A31-A0)

The Normal Read instruction code is transmitted via the SI line, followed by three (A23 - A0) or four (A31 - A0) address bytes of the first memory location to be read as above. A total of 24 or 32 address bits are shifted in, but only  $A_{VMSB}$  (Valid Most Significant Bit) -  $A_0$  are decoded. The remaining bits ( $A_{31} - A_{VMSB+1}$ ) are ignored. The first byte addressed can be at any memory location. Upon completion, any data on the SI will be ignored. Refer to Table 8.5 for the related Address Key.

The first byte data (D7 - D0) is shifted out on the SO line, MSB first. A single byte of data, or up to the whole memory array, can be read out in one Normal Read instruction. The address is automatically incremented by one after each byte of data is shifted out. The read operation can be terminated at any time by driving CE# high (VIH) after the data comes out. When the highest address of the device is reached, the address counter will roll over to the 000000h address, allowing the entire memory to be read in one continuous Read instruction.

If the Normal Read instruction is issued while an Erase, Program or Write operation is in process (WIP=1) the instruction is ignored and will not have any effects on the current operation.

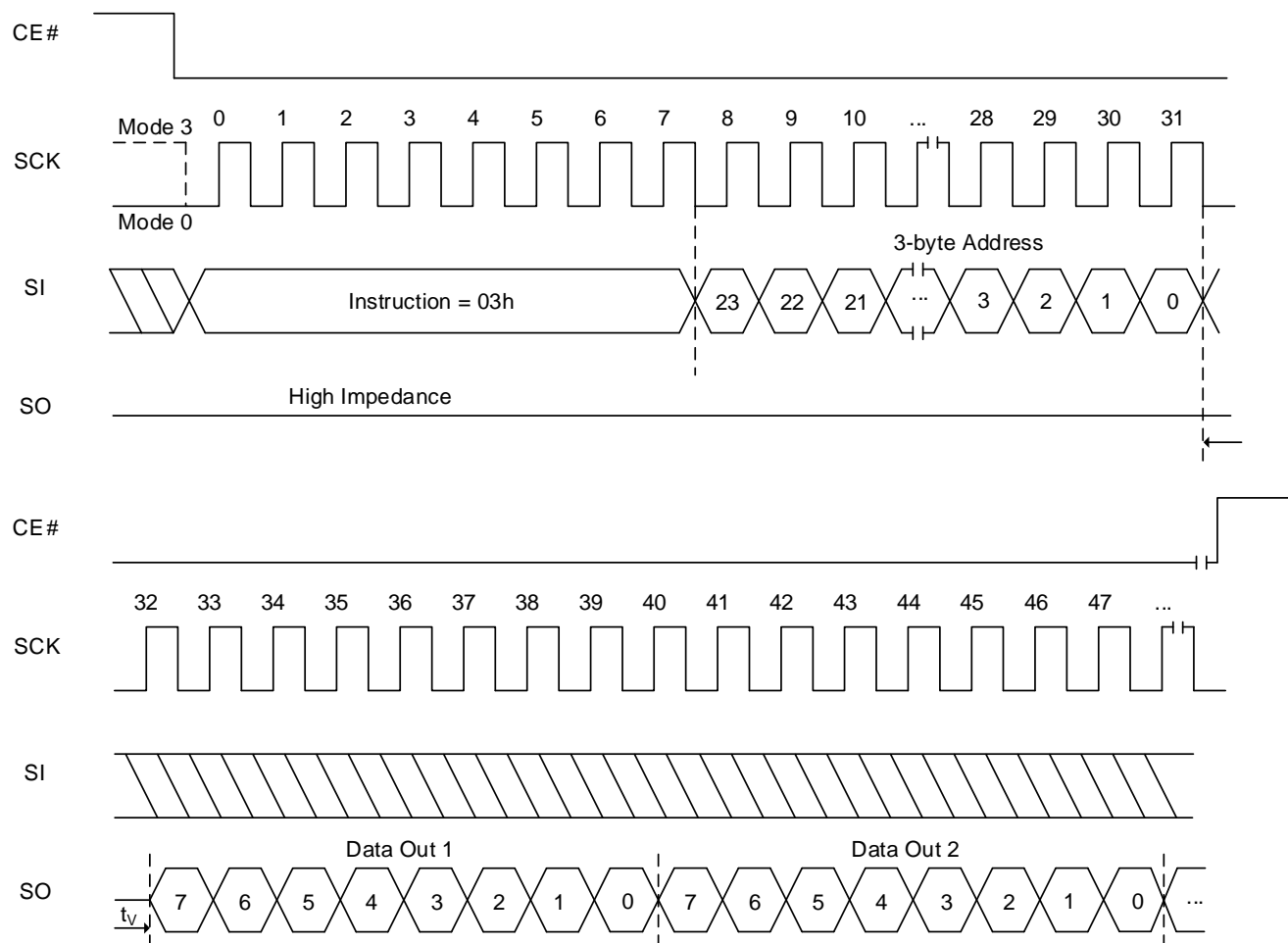
**Table 8.6 Address Key**

Mode	Valid Address	256Mb
3 byte address	$A_{VMSB}^{(2)} - A_0$	A23-A0
4 byte address		A24-A0 ( $A_{31}-A_{25}=X$ ) <sup>(2)</sup>

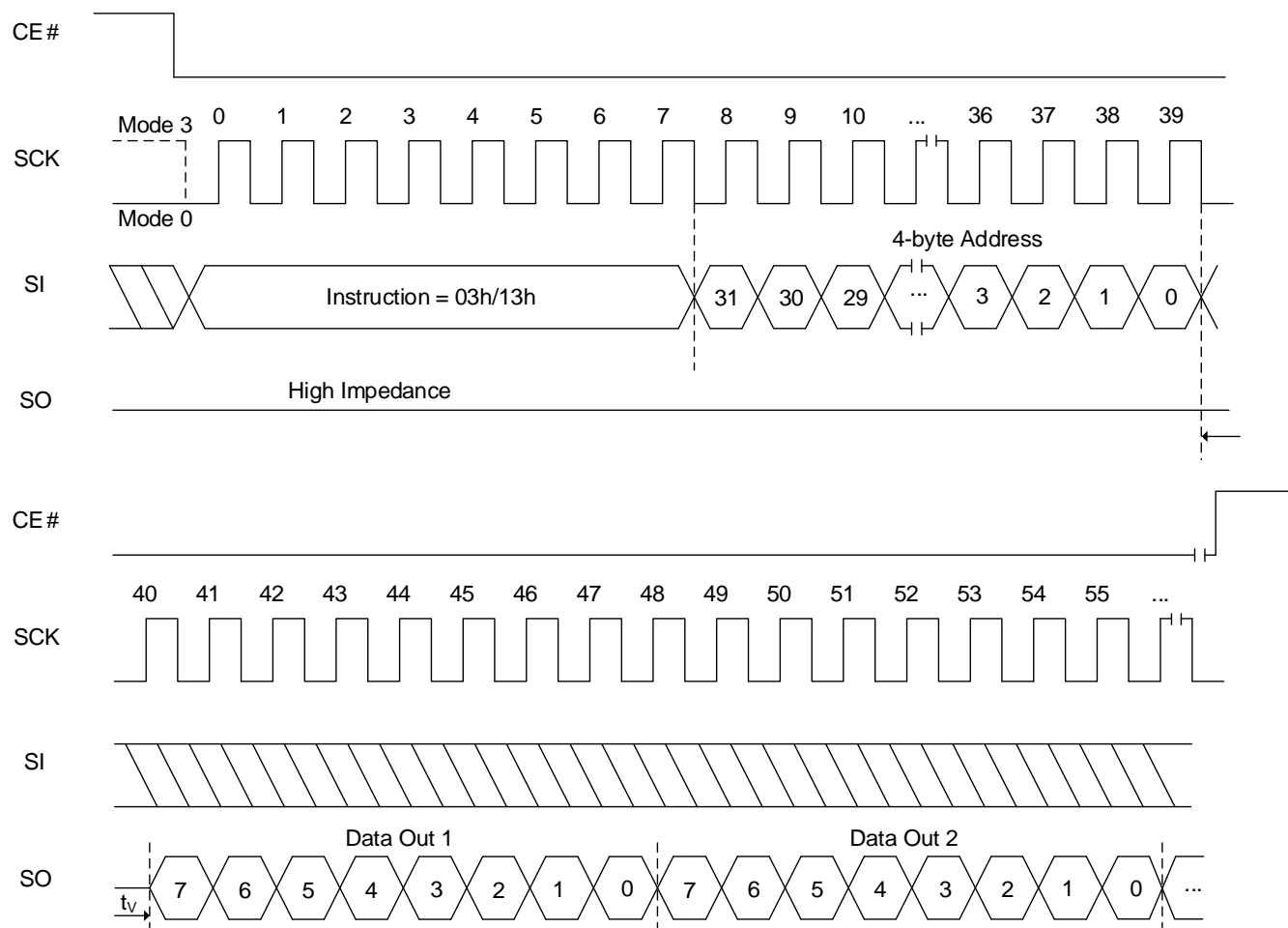
**Notes:**

1. **X=Don't Care**
2.  **$A_{VMSB}$  is a Valid MSB. In 4 byte address for 256Mb, A31 is an MSB, and A24 is a Valid MSB.**

**Figure 8.1 Normal Read Sequence (03h [EXTADD=0], 3-byte address)**



**Figure 8.2 Normal Read Sequence (03h [EXTADD=1] or 13h, 4-byte address)**



#### **8.4 FAST READ OPERATION (FRD, 0Bh or 4FRD, 0Ch)**

The Fast Read (FRD, 4FRD) instruction is used to read memory data at up to a 166MHz clock.

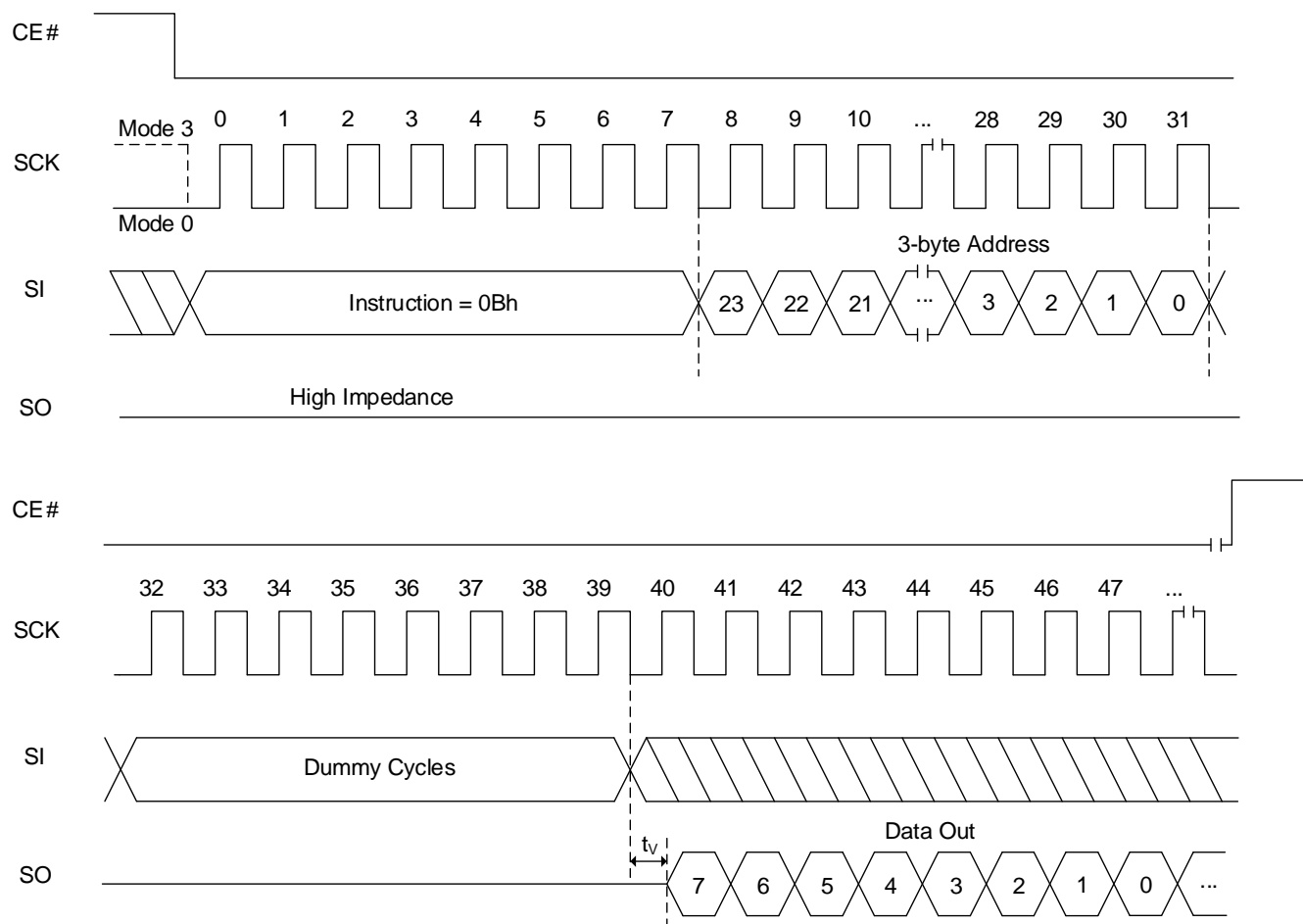
- 0Bh (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- 0Bh (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- 0Ch is followed by a 4-byte address (A31-A0)

The Fast Read instruction code is followed by three or four address bytes as above and dummy cycles (configurable, default is 8 clocks), transmitted via the SI line, with each bit latched-in during the rising edge of SCK. Then the first data byte from the address is shifted out on the SO line, with each bit shifted out at a maximum frequency  $f_{CT}$ , during the falling edge of SCK.

The first byte addressed can be at any memory location. The address is automatically incremented by one after each byte of data is shifted out. When the highest address is reached, the address counter will roll over to the 000000h address, allowing the entire memory to be read with a single Fast Read instruction. The Fast Read instruction is terminated by driving CE# high (VIH).

If the Fast Read instruction is issued while an Erase, Program or Write cycle is in process (WIP=1) the instruction is ignored without affecting the current cycle.

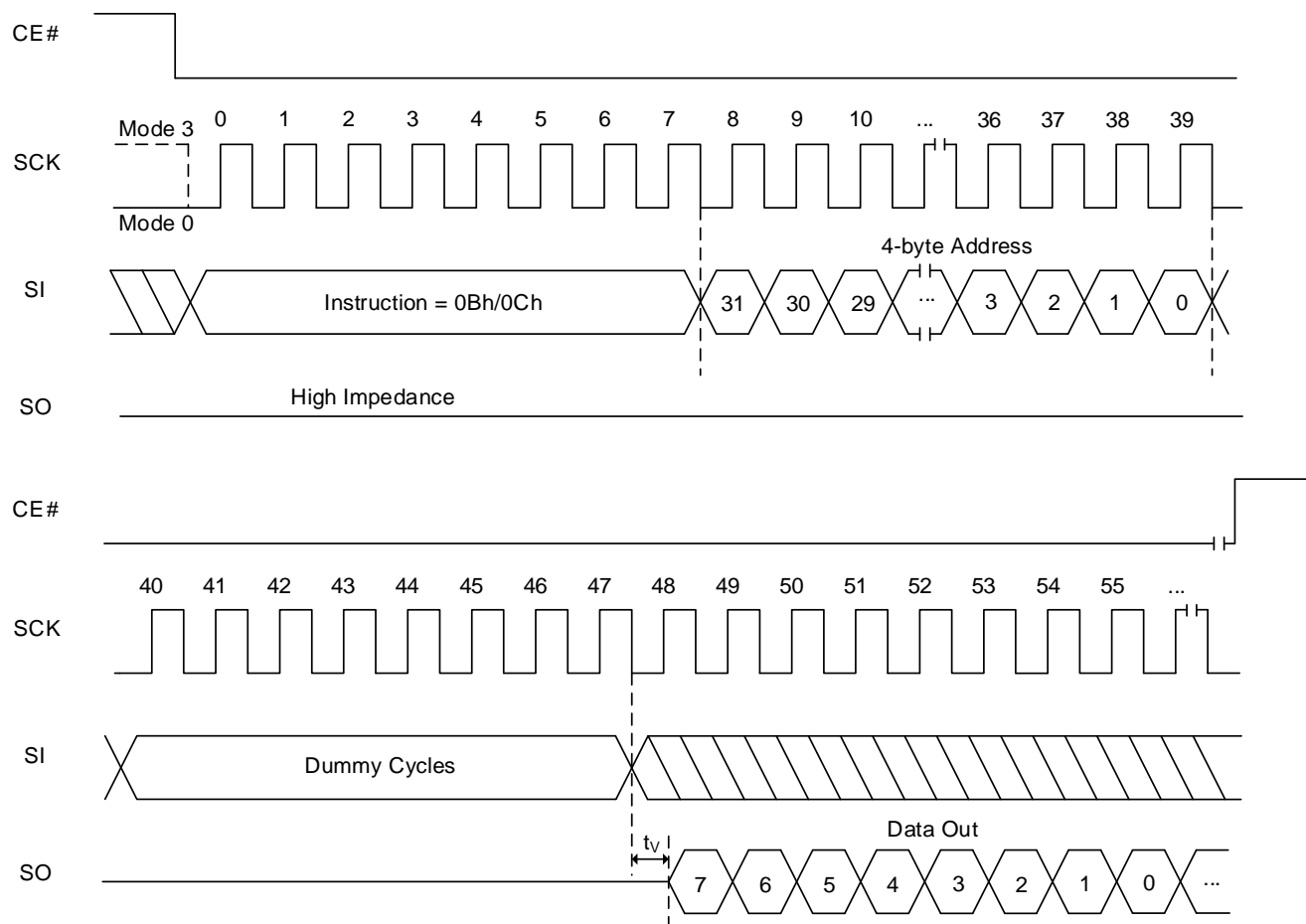
**Figure 8.3 Fast Read Sequence (0Bh [EXTADD=0], 3-byte address)**



**Note:** Dummy cycles depends on Read Parameter setting. Detailed information in Table 6.11 Read Dummy Cycles.



**Figure 8.4 Fast Read Sequence (0Bh [EXTADD=1] or 0Ch, 4-byte address)**



**Note:** Dummy cycles depends on Read Parameter setting. Detailed information in Table 6.11 Read Dummy Cycles.

### FAST READ OPERATION IN QPI MODE (FRD, 0Bh or 4FRD, 0Ch)

The Fast Read (FRD) instruction in QPI mode is used to read memory data at up to a 166MHz clock.

- 0Bh (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- 0Bh (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- 0Ch is followed by a 4-byte address (A31-A0)

The Fast Read instruction code (2 clocks) is followed by three (6 clocks) or four (8 clocks) address bytes as above and 6 dummy cycles (configurable, default is 6 clocks), transmitted via the IO3, IO2, IO1 and IO0 lines, with each bit latched-in during the rising edge of SCK. Then the first data byte addressed is shifted out on the IO3, IO2, IO1 and IO0 lines, with each bit shifted out at a maximum frequency  $f_{CT}$ , during the falling edge of SCK.

The first byte addressed can be at any memory location. The address is automatically incremented by one after each byte of data is shifted out. When the highest address is reached, the address counter will roll over to the 000000h address, allowing the entire memory to be read with a single Fast Read QPI instruction. The Fast Read QPI instruction is terminated by driving CE# high (VIH).

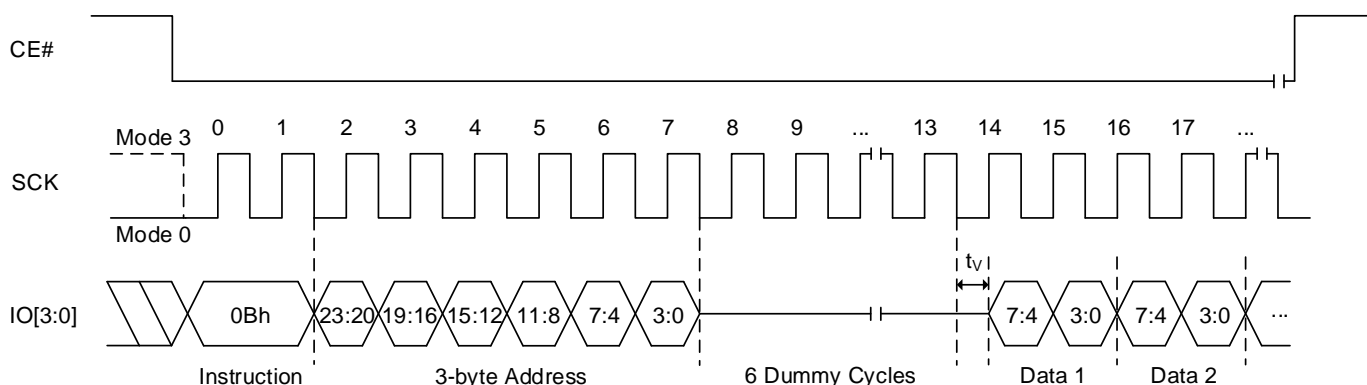
If the Fast Read instruction in QPI mode is issued while an Erase, Program or Write cycle is in process (WIP=1) the instruction is ignored without affecting the current cycle.

The Fast Read sequence in QPI mode is also applied to the commands in the following table 8.7. However, only 3-byte address mode QPI sequence is applied for RDUID, and IRRD commands.

**Table 8.7 Instructions that Fast Read sequence in QPI Mode is applied to**

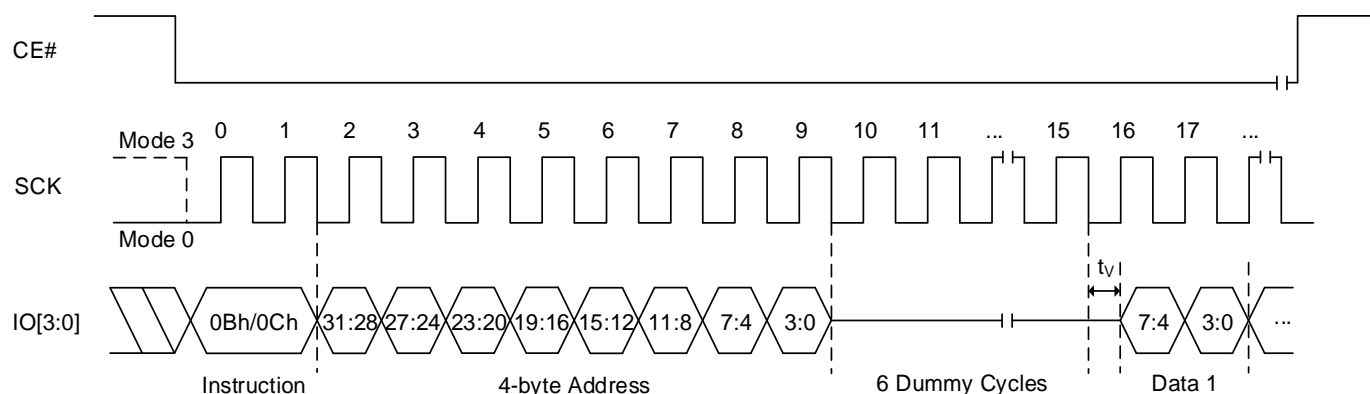
Instruction Name	Operation	Hex Code
FRQIO	Fast Read Quad I/O	EBh
RDUID	Read Unique ID	4Bh
IRRD	Read Information Row	68h

**Figure 8.5 Fast Read Sequence In QPI Mode (0Bh [EXTADD=0], 3-byte address)**



**Note: Number of dummy cycles depends on Read Parameter setting. Detailed information in Table 6.11 Read Dummy Cycles.**

Figure 8.6 Fast Read Sequence In QPI Mode (0Bh [EXTADD=1] or 0Ch, 4-byte address)



**Note:** Number of dummy cycles depends on Read Parameter setting. Detailed information in Table 6.11 Read Dummy Cycles.

## 8.5 HOLD OPERATION

HOLD# is used in conjunction with CE# to select the device. When the device is selected and a serial sequence is underway, HOLD# can be used to pause the serial communication with the master device without resetting the serial sequence. To pause, HOLD# is brought low while the SCK signal is low. To resume serial communication, HOLD# is brought high while the SCK signal is low (SCK may still toggle during HOLD). Inputs to SI will be ignored while SO is in the high impedance state, during HOLD.

**Note:** HOLD is not supported with QE=1.

Timing graph can be referenced in AC Parameters Figure 9.4.

**Call Factory for HOLD# function.**

## 8.6 FAST READ DUAL I/O OPERATION (FRDIO, BBh or 4FRDIO, BCh)

The Fast Read Dual I/O (FRDIO, 4FRDIO) instruction allows the address bits to be input two bits at a time. This may allow for code to be executed directly from the SPI in some applications.

- BBh (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- BBh (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- BCh is followed by a 4-byte address (A31-A0)

The FRDIO/4FRDIO instruction code is followed by three or four address bytes as above and dummy cycles (configurable, default is 4 clocks), transmitted via the IO1 and IO0 lines, with each pair of bits latched-in during the rising edge of SCK. The address MSB is input on IO1, the next bit on IO0, and this shift pattern continues to alternate between the two lines. Depending on the usage of AX read operation mode, a mode byte may be located after address input.

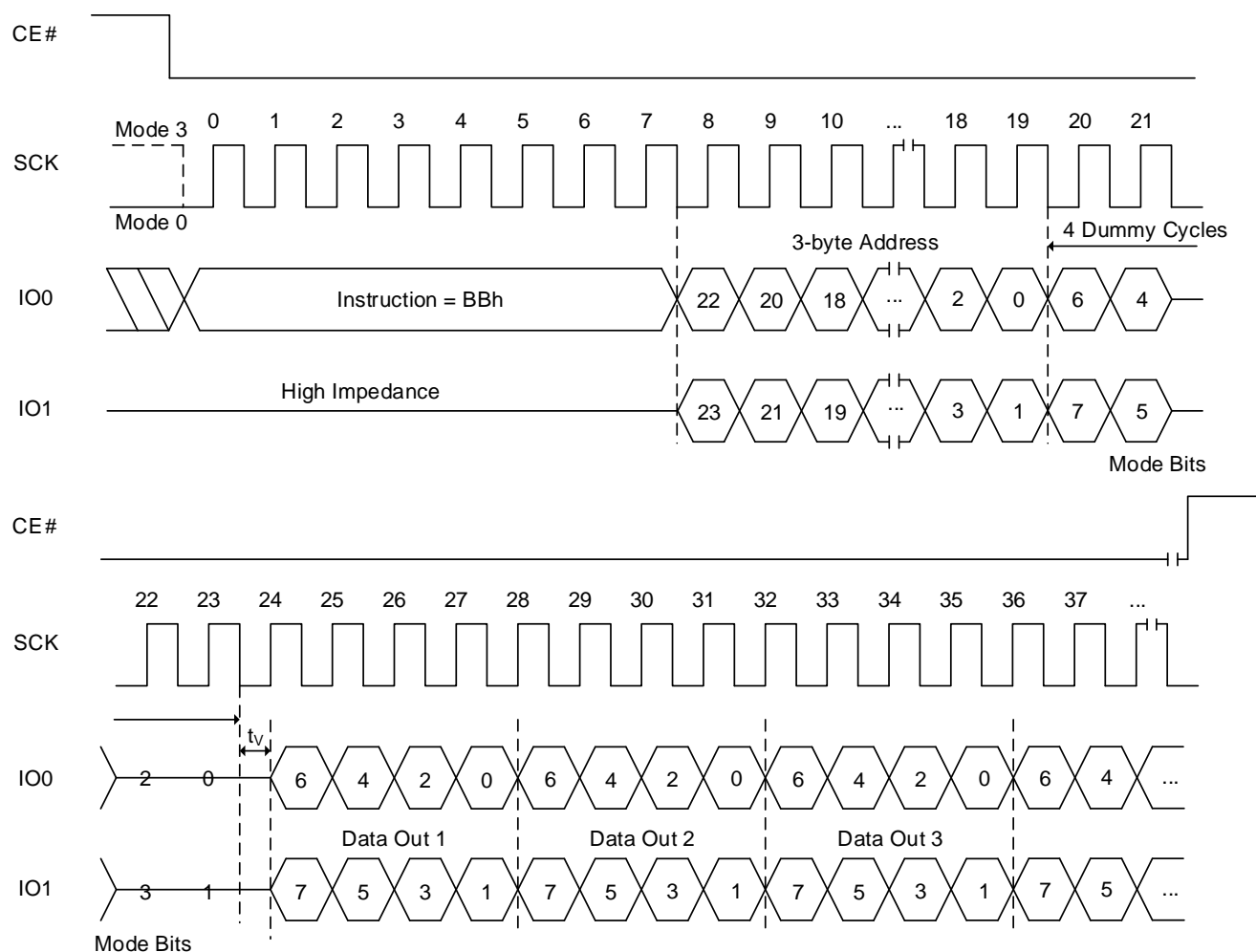
The first data byte addressed is shifted out on the IO1 and IO0 lines, with each pair of bits shifted out at a maximum frequency  $f_{CT}$ , during the falling edge of SCK. The first bit (MSB) is output on IO1, while simultaneously the second bit is output on IO0. Figures 8.7 and 8.8 illustrates the timing sequence.

The first byte addressed can be at any memory location. The address is automatically incremented by one after each byte of data is shifted out. When the highest address is reached, the address counter will roll over to the 000000h address, allowing the entire memory to be read with a single FRDIO/4FRDIO instruction. The FRDIO/4FRDIO instruction is terminated by driving CE# high ( $V_{IH}$ ).

The device supports the AX read operation by applying mode bits during dummy period. Mode bits consist of 8 bits, such as M7 to M0. Four cycles after address input are reserved for Mode bits in FRDIO/4FRDIO execution. M7 to M4 are important for enabling this mode. M3 to M0 become don't care for future use. When M[7:4]=1010(Ah), it enables the AX read operation and subsequent FRDIO/4FRDIO execution skips command code. It saves cycles as described in Figures 8.9 and 8.10. When the code is different from AXh (where X is don't care), the device exits the AX read operation. After finishing the read operation, device becomes ready to receive a new command. SPI or QPI mode configuration retains the prior setting. Mode bit must be applied during dummy cycles. Number of dummy cycles in Table 6.11 includes number of mode bit cycles. If dummy cycles are configured as 4 cycles, data output will start right after mode bit is applied.

If the FRDIO/4FRDIO instruction is issued while an Erase, Program or Write cycle is in process (WIP=1) the instruction is ignored and will not affect the current cycle.

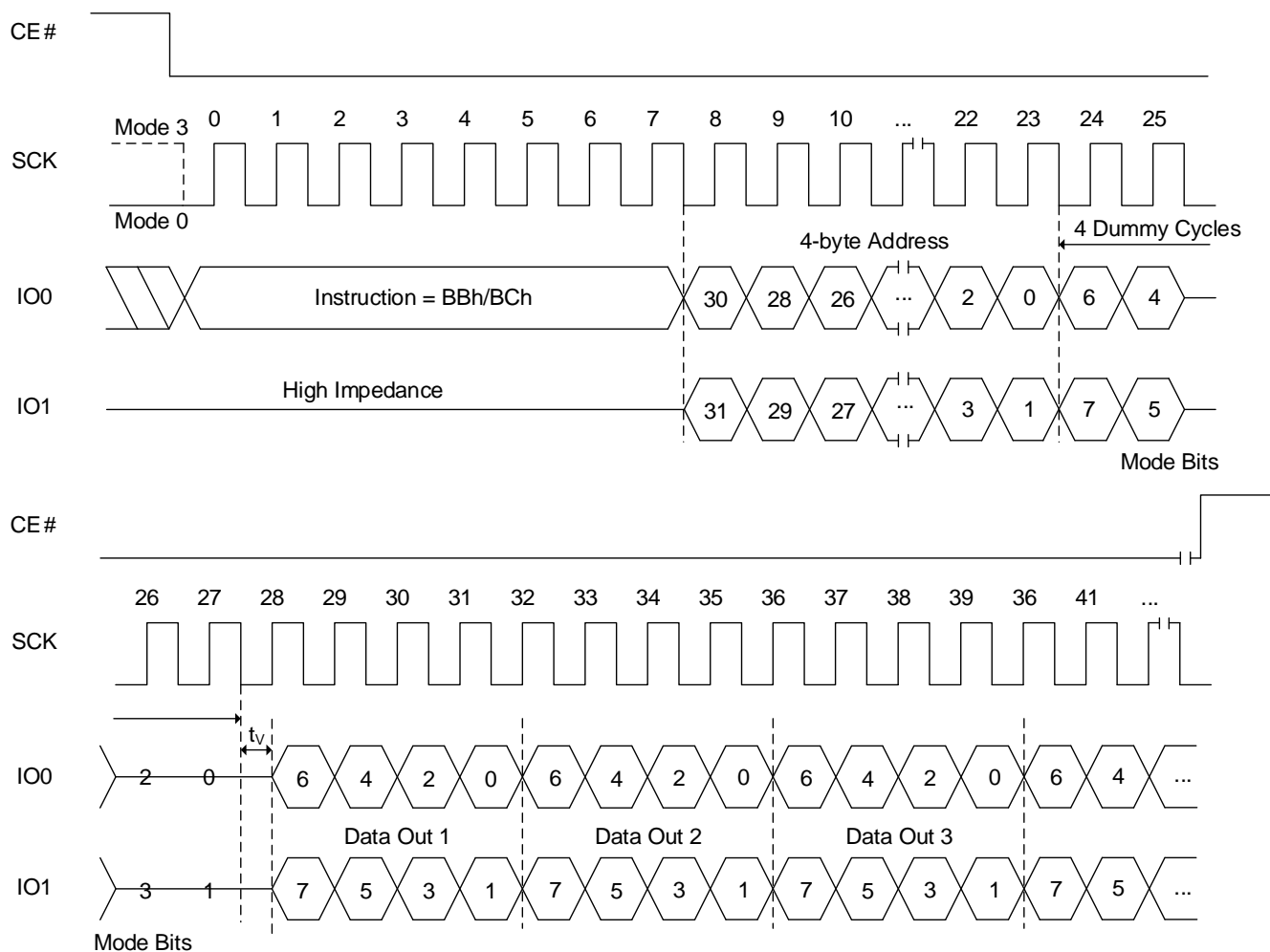
Figure 8.7 Fast Read Dual I/O Sequence (BBh [EXTADD=0], 3-byte address)



**Notes:**

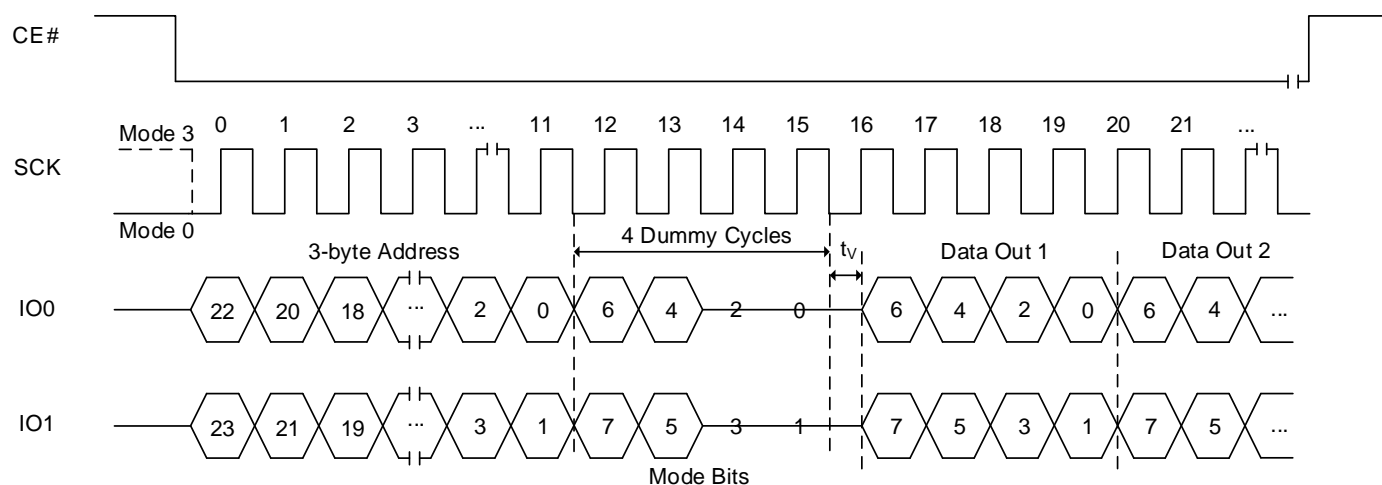
1. If the mode bits=AXh (where X is don't care), it can execute the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

**Figure 8.8 Fast Read Dual I/O Sequence (BBh [EXTADD=1] or BCh, 4-byte address)**

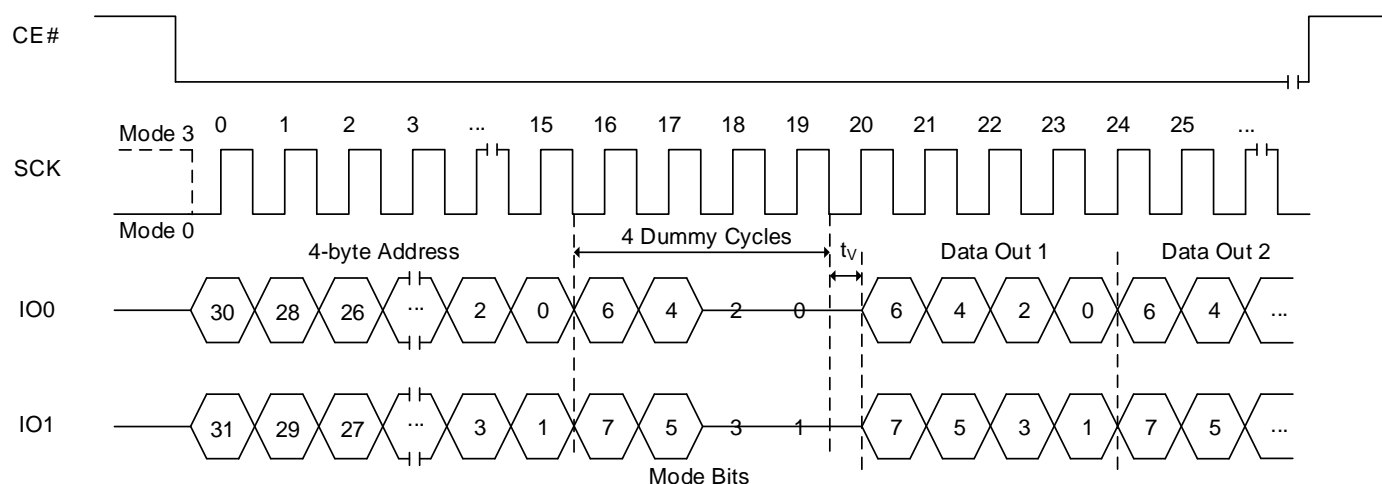


**Notes:**

1. If the mode bits=AXh (where X is don't care), it can execute the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

**Figure 8.9 Fast Read Dual I/O AX Read Sequence (BBh [EXTADD=0], 3-byte address)**

**Notes:**

1. If the mode bits=AXh (where X is don't care), it will keep executing the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

**Figure 8.10 Fast Read Dual I/O AX Read Sequence (BBh [EXTADD=1] or BCh, 4-byte address)**

**Notes:**

1. If the mode bits=AXh (where X is don't care), it will keep executing the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

### 8.7 FAST READ DUAL OUTPUT OPERATION (FRDO, 3Bh or 4FRDO, 3Ch)

The FRDO/4FRDO instruction is used to read memory data on two output pins each at up to a 166MHz clock.

- 3Bh (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- 3Bh (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- 3Ch is followed by a 4-byte address (A31-A0)

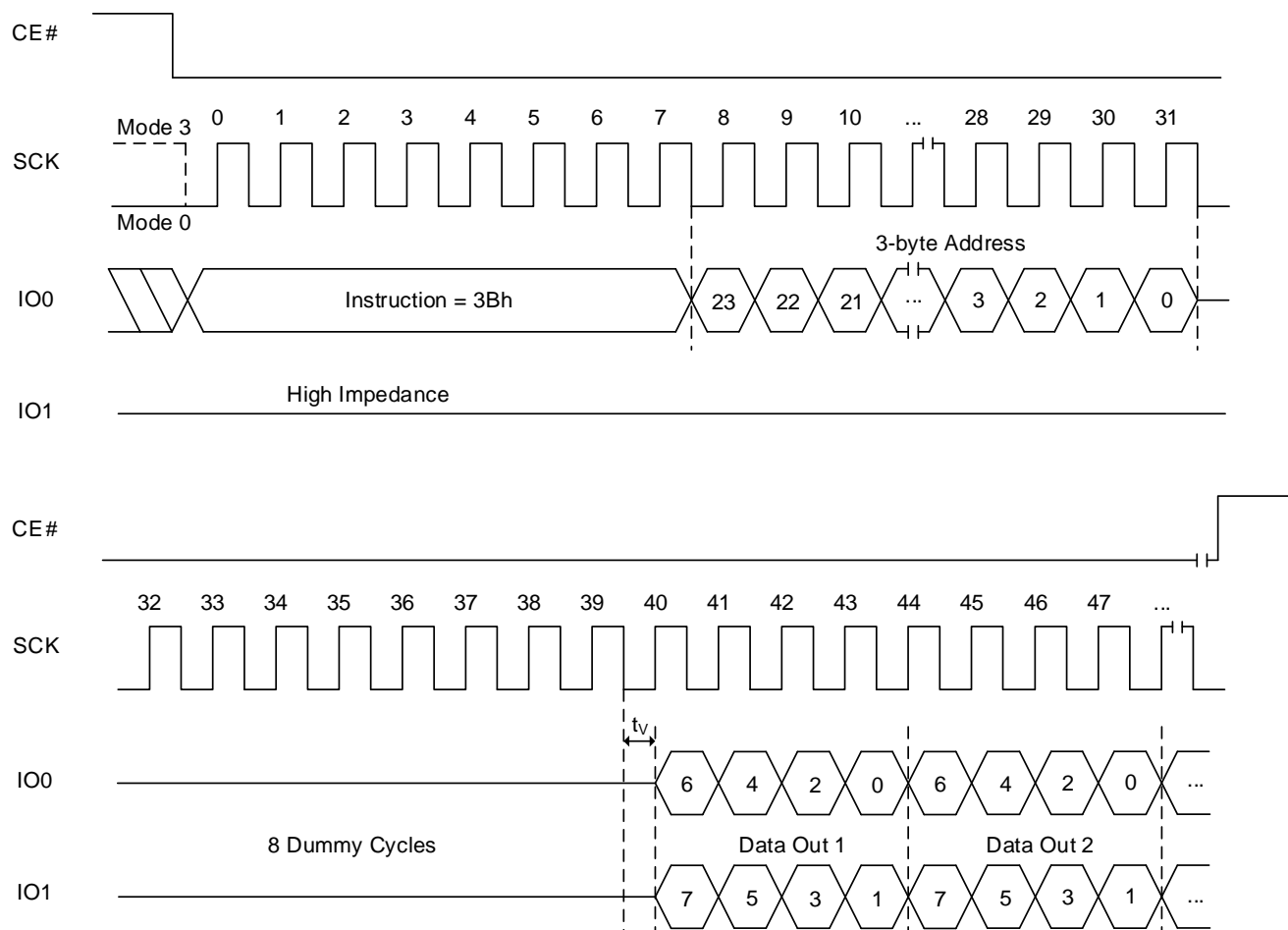
The FRDO/4FRDO instruction code is followed by three or four address bytes as above and dummy cycles (configurable, default is 8 clocks), transmitted via the IO0 line, with each bit latched-in during the rising edge of SCK. Then the first data byte addressed is shifted out on the IO1 and IO0 lines, with each pair of bits shifted out at a maximum frequency  $f_{CT}$ , during the falling edge of SCK. The first bit (MSB) is output on IO1. Simultaneously, the second bit is output on IO0.

The first byte addressed can be at any memory location. The address is automatically incremented by one after each byte of data is shifted out. When the highest address is reached, the address counter will roll over to the 000000h address, allowing the entire memory to be read with a single FRDO/4FRDO instruction. The instruction FRDO/4FRDO is terminated by driving CE# high (VIH).

If the FRDO/4FRDO instruction is issued while an Erase, Program or Write cycle is in process (WIP=1) the instruction is ignored and will not have any effects on the current cycle.

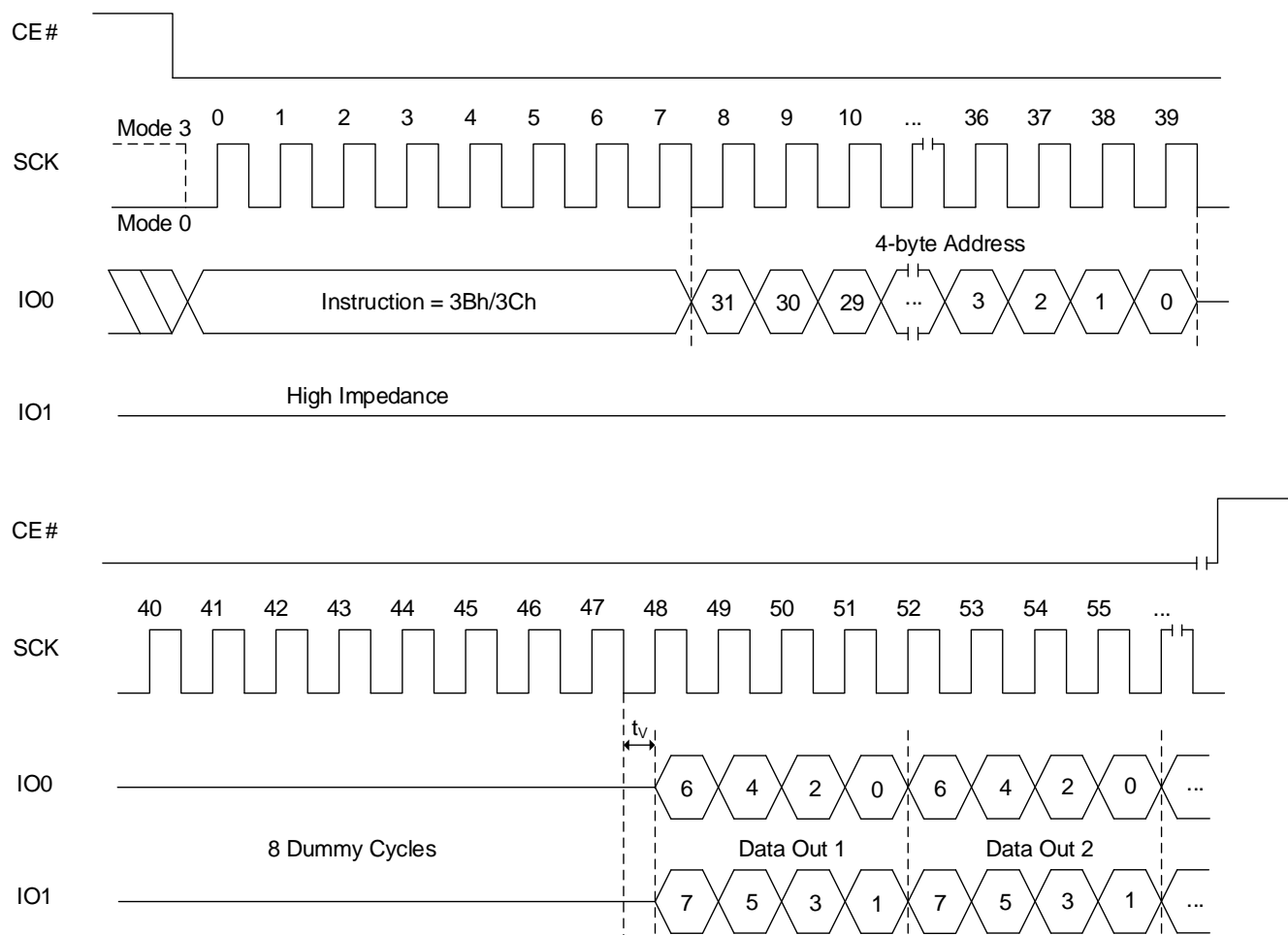


**Figure 8.11 Fast Read Dual Output Sequence (3Bh [EXTADD=0], 3-byte address)**



**Note:** Dummy cycles depends on Read Parameter setting. Detailed information in Table 6.11 Read Dummy Cycles.

**Figure 8.12 Fast Read Dual Output Sequence (3Bh [EXTADD=1] or 3Ch, 4-byte address)**



**Note:** Dummy cycles depends on Read Parameter setting. Detailed information in Table 6.11 Read Dummy Cycles.

### 8.8 FAST READ QUAD OUTPUT OPERATION (FRQO, 6Bh or 4FRQO 6Ch)

The FRQO/4FRQO instruction is used to read memory data on four output pins each at up to a 166 MHz clock.

A Quad Enable (QE) bit of Status Register must be set to "1" before sending the Fast Read Quad Output instruction.

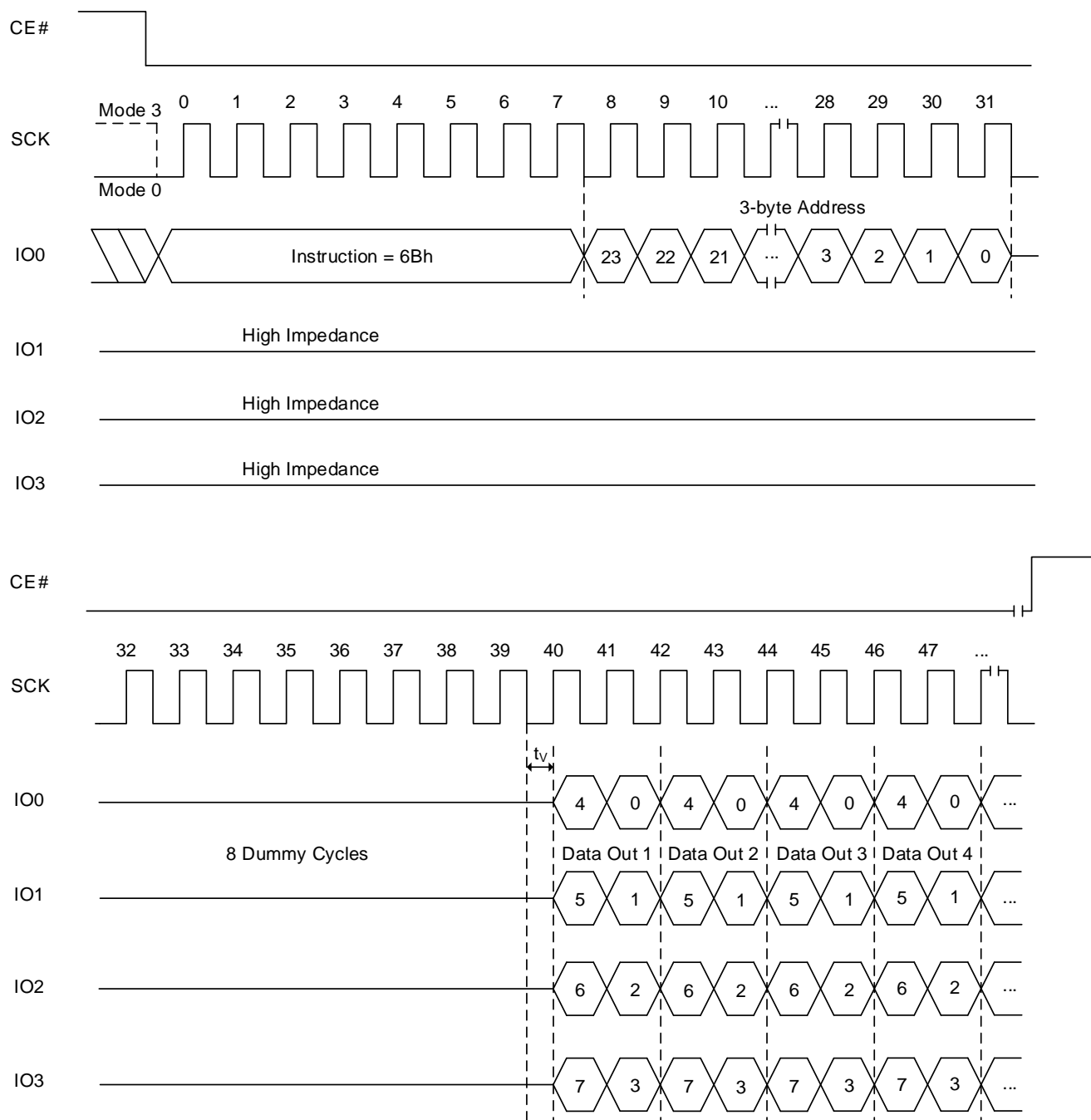
- 6Bh (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- 6Bh (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- 6Ch is followed by a 4-byte address (A31-A0)

The FRQO/4FRQO instruction code is followed by three or four address bytes as above and dummy cycles (configurable, default is 8 clocks), transmitted via the IO0 line, with each bit latched-in during the rising edge of SCK. Then the first data byte addressed is shifted out on the IO3, IO2, IO1 and IO0 lines, with each group of four bits shifted out at a maximum frequency  $f_{CT}$ , during the falling edge of SCK. The first bit (MSB) is output on IO3, while simultaneously the second bit is output on IO2, the third bit is output on IO1, etc.

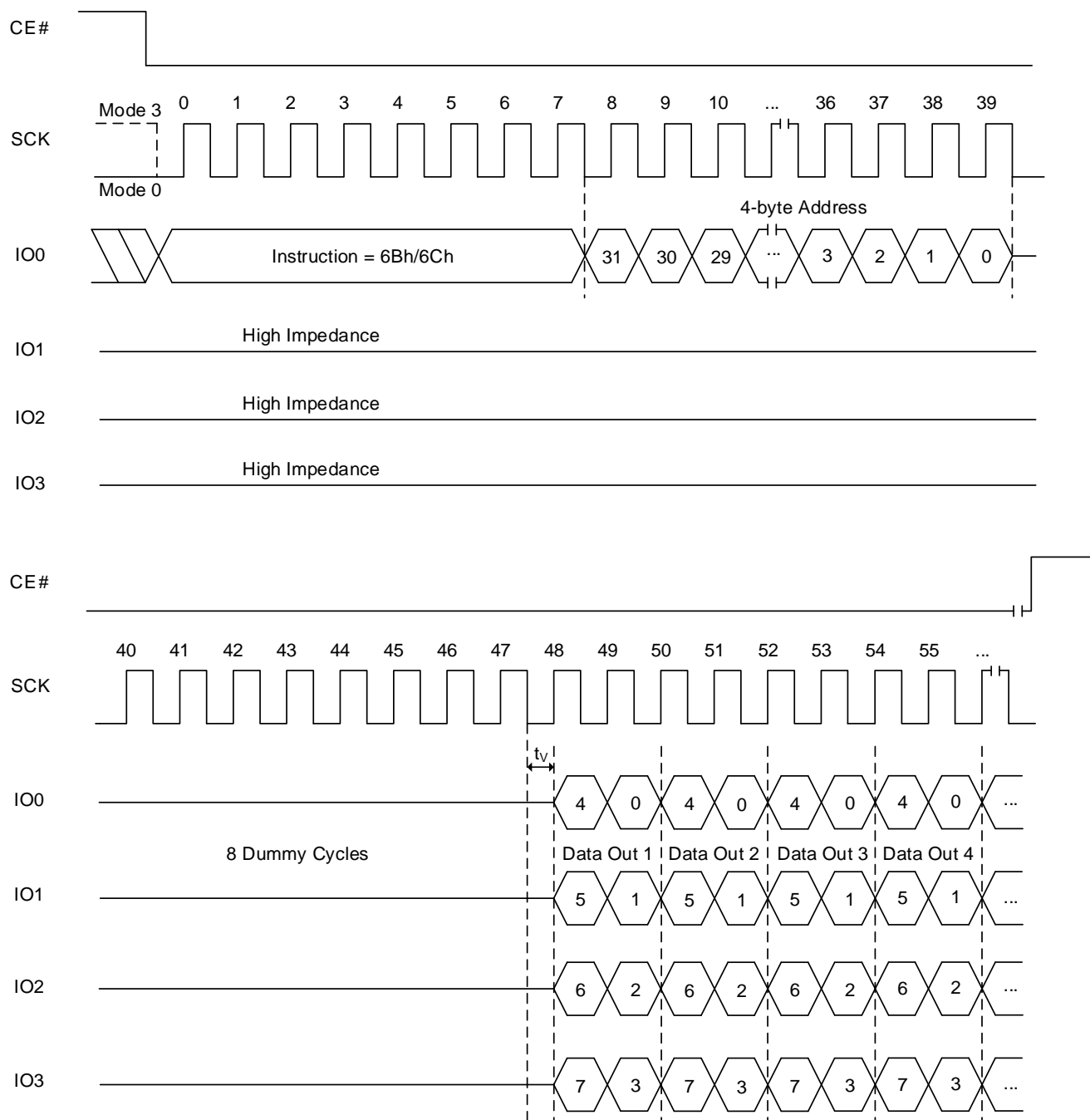
The first byte addressed can be at any memory location. The address is automatically incremented after each byte of data is shifted out. When the highest address is reached, the address counter will roll over to the 000000h address, allowing the entire memory to be read with a single FRQO/4FRQO instruction. FRQO/4FRQO instruction is terminated by driving CE# high (VIH).

If a FRQO/4FRQO instruction is issued while an Erase, Program or Write cycle is in process (WIP=1) the instruction is ignored and will not have any effects on the current cycle.

**Figure 8.13 Fast Read Quad Output Sequence (6Bh [EXTADD=0], 3-byte address)**



**Figure 8.14 Fast Read Quad Output Sequence (6Bh [EXTADD=1] or 6Ch, 4-byte address)**



**Note:** Dummy cycles depends on Read Parameter setting. Detailed information in Table 6.11 Read Dummy Cycles.

### 8.9 FAST READ QUAD I/O OPERATION (FRQIO, EBh or 4FRQIO, ECh)

The FRQIO/4FRQIO instruction allows the address bits to be input four bits at a time. This may allow for code to be executed directly from the SPI in some applications.

A Quad Enable (QE) bit of Status Register must be set to "1" before sending the Fast Read Quad I/O instruction.

- EBh (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- EBh (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- ECh is followed by a 4-byte address (A31-A0)

The FRQIO/4FRQIO instruction code is followed by three or four address bytes as above and dummy cycles (configurable, default is 6 clocks), transmitted via the IO3, IO2, IO1 and IO0 lines, with each group of four bits latched-in during the rising edge of SCK. The address of MSB inputs on IO3, the next bit on IO2, the next bit on IO1, the next bit on IO0, and continue to shift in alternating on the four. Depending on the usage of AX read operation mode, a mode byte may be located after address input.

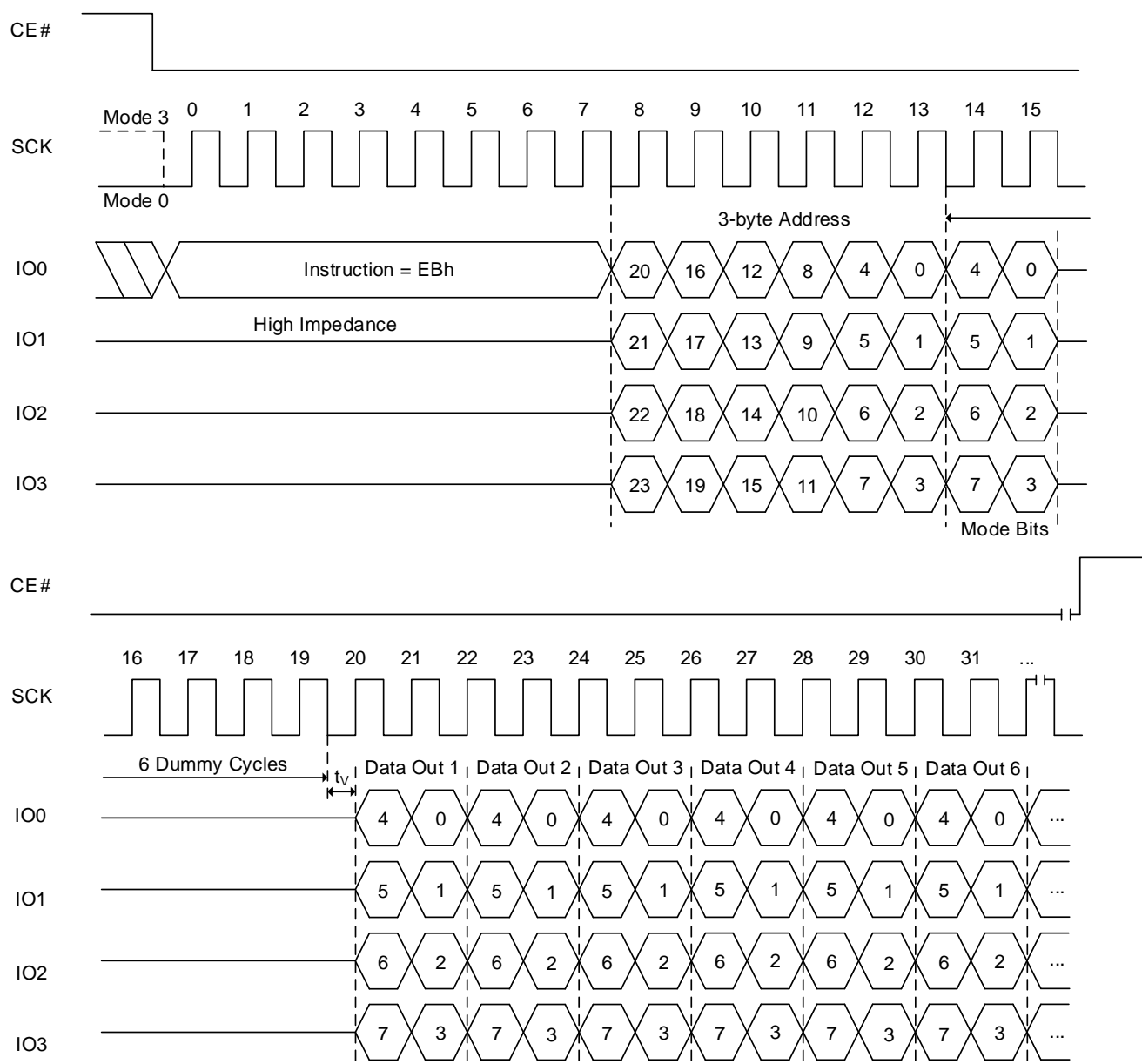
The first data byte addressed is shifted out on the IO3, IO2, IO1 and IO0 lines, with each group of four bits shifted out at a maximum frequency  $f_{CT}$ , during the falling edge of SCK. The first bit (MSB) is output on IO3, while simultaneously the second bit is output on IO2, the third bit is output on IO1, etc. Figures 8.15 and 8.16 illustrates the timing sequence.

The first byte addressed can be at any memory location. The address is automatically incremented after each byte of data is shifted out. When the highest address is reached, the address counter will roll over to the 000000h address, allowing the entire memory to be read with a single FRQIO/4FRQIO instruction. FRQIO/4FRQIO instruction is terminated by driving CE# high ( $V_{IH}$ ).

The device supports the AX read operation by applying mode bits during dummy period. Mode bits consist of 8 bits, such as M7 to M0. Two cycles after address input are reserved for Mode bits in FRQIO/4FRQIO execution. M7 to M4 are important for enabling this mode. M3 to M0 become don't care for future use. When M[7:4]=1010(Ah), it enables the AX read operation and subsequent FRQIO/4FRQIO execution skips command code. It saves cycles as described in Figures 8.17 and 8.18. When the code is different from AXh (where X is don't care), the device exits the AX read operation. After finishing the read operation, device becomes ready to receive a new command. SPI or QPI mode configuration retains the prior setting. Mode bit must be applied during dummy cycles. Number of dummy cycles in Table 6.11 includes number of mode bit cycles. If dummy cycles are configured as 6 cycles, data output will start right after mode bits and 4 additional dummy cycles are applied.

If the FRQIO/4FRQIO instruction is issued while an Erase, Program or Write cycle is in process (WIP=1) the instruction is ignored and will not have any effects on the current cycle.

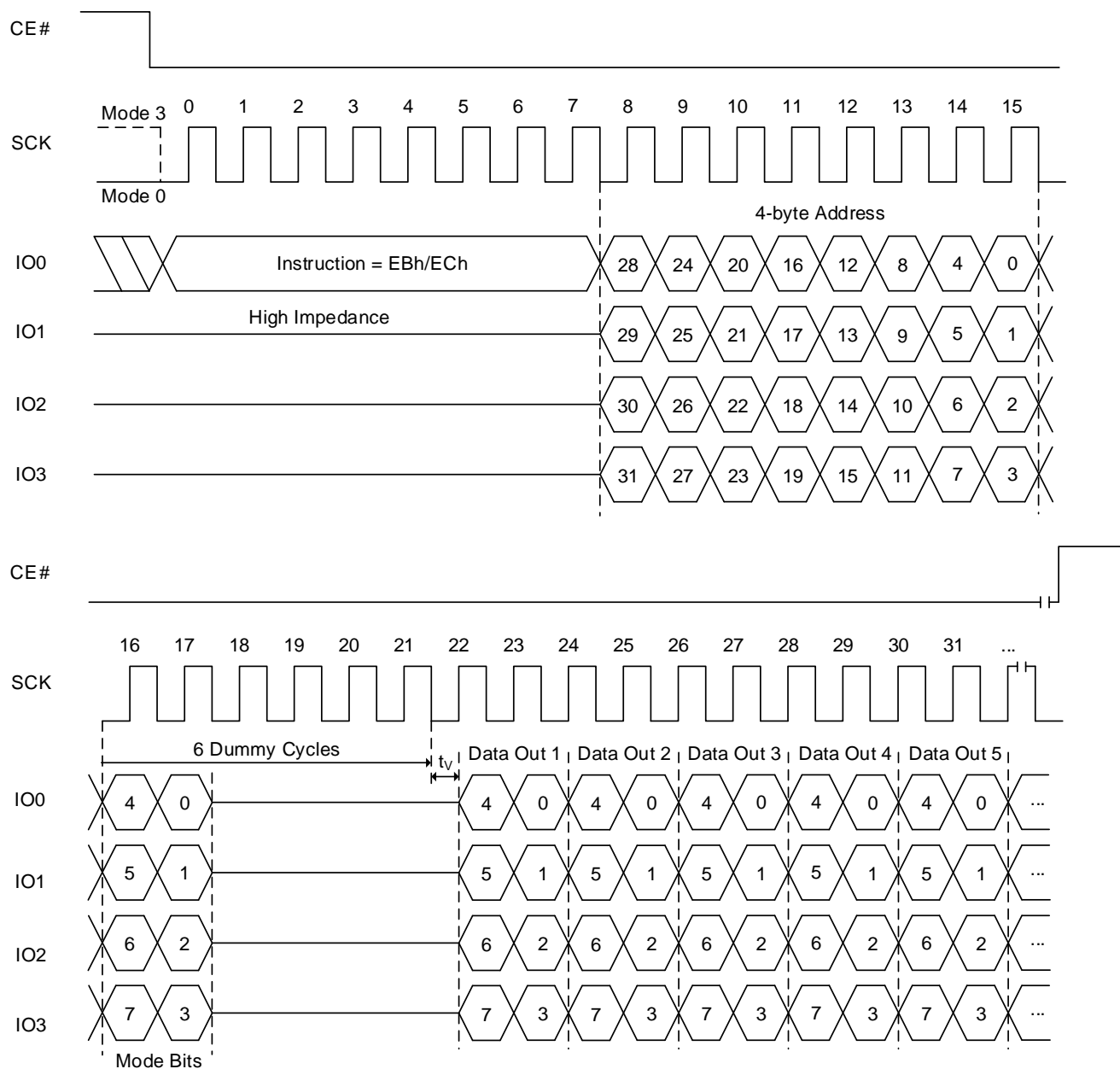
**Figure 8.15 Fast Read Quad I/O Sequence (EBh [EXTADD=0], 3-byte address)**



**Notes:**

1. If the mode bits=AXh (where X is don't care), it can execute the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

**Figure 8.16 Fast Read Quad I/O Sequence (EBh [EXTADD=1] or ECh, 4-byte address)**

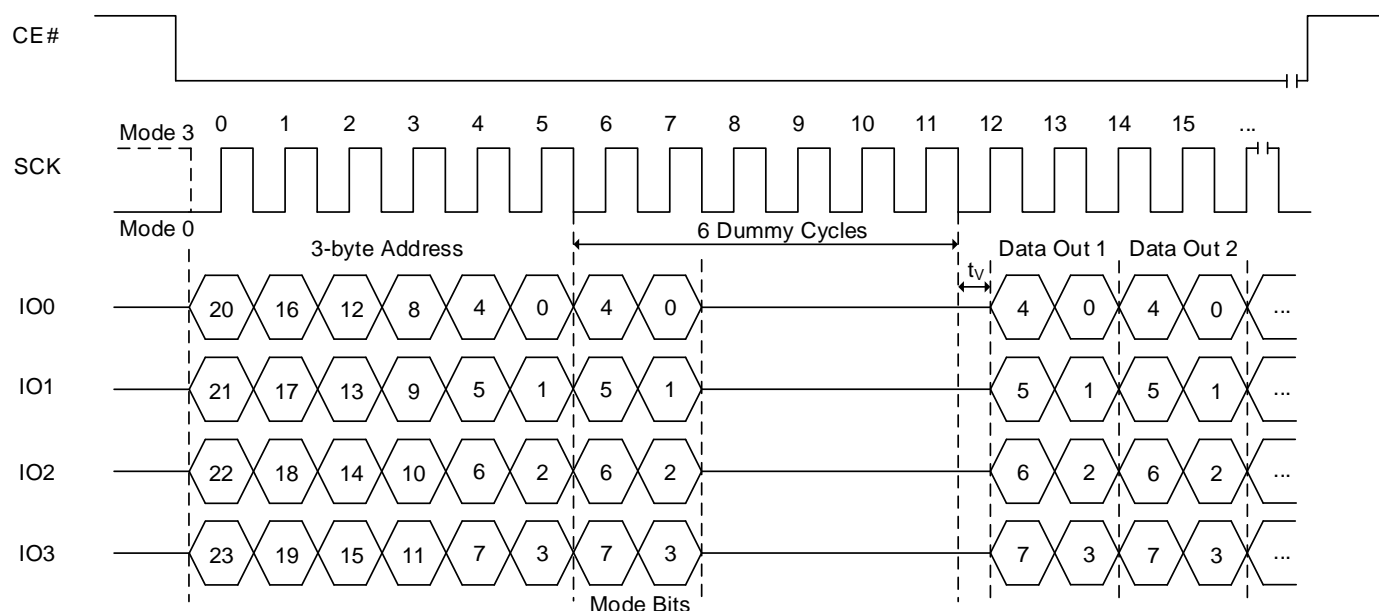


**Notes:**

1. If the mode bits=AXh (where X is don't care), it can execute the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.



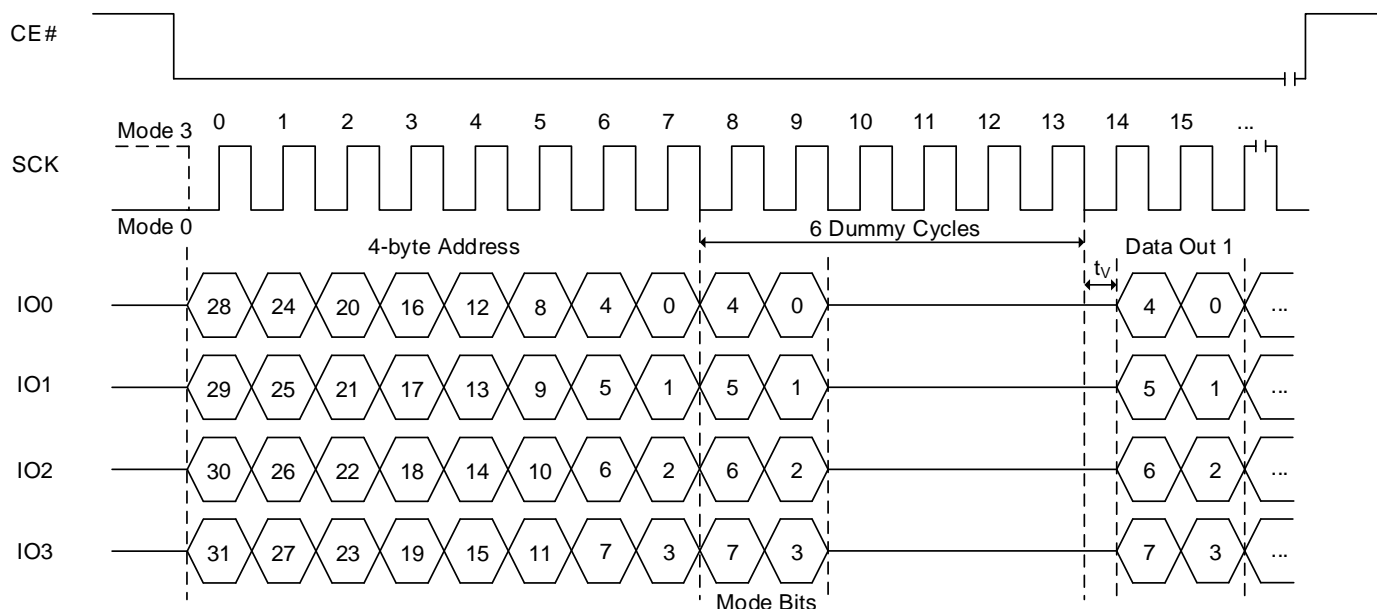
Figure 8.17 Fast Read Quad I/O AX Read Sequence (EBh [EXTADD=0], 3-byte address)



**Notes:**

1. If the mode bits=AXh (where X is don't care), it will keep executing the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

**Figure 8.18 Fast Read Quad I/O AX Read Sequence (EBh [EXTADD=1] or ECh, 4-byte address)**



**Notes:**

1. If the mode bits=AXh (where X is don't care), it will keep executing the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

### FAST READ QUAD I/O OPERATION IN QPI MODE (FRQIO, EBh or 4FRQIO, ECh)

The FRQIO/4FRQIO instruction in QPI mode is used to read memory data at up to a 166MHz clock.

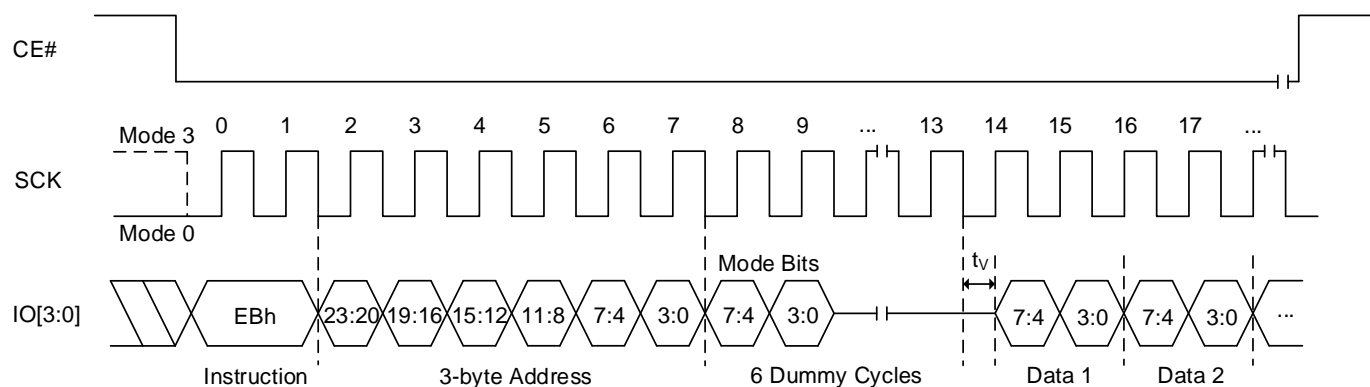
It is not required to set QE bit to “1”.before Fast Read Quad I/O instruction in QPI mode.

The FRQIO/4FRQIO instruction in QPI mode utilizes all four IO lines to input the instruction code so that only two clocks are required, while the FRQIO/4FRQIO instruction requires that the byte-long instruction code is shifted into the device only via IO0 line in eight clocks. As a result, 6 command cycles will be reduced by the FRQIO/4FRQIO QPI instruction. In addition, subsequent address and data out are shifted in/out via all four IO lines like the FRQIO/4FRQIO instruction. In fact, except for the command cycle, the FRQIO/4FRQIO operation in QPI mode is exactly same as the FRQIO/4FRQIO operation in SPI mode.

The device supports the AX read operation by applying mode bits during dummy period. Mode bits consist of 8 bits, such as M7 to M0. Two cycles after address input are reserved for Mode bits in FRQIO/4FRQIO execution. M7 to M4 are important for enabling this mode. M3 to M0 become don't care for future use. When M[7:4]=1010(Ah), it enables the AX read operation and subsequent FRQIO/4FRQIO execution skips command code. It saves cycles as described in Figures 8.17 and 8.18. When the code is different from AXh (where X is don't care), the device exits the AX read operation. After finishing the read operation, device becomes ready to receive a new command. SPI or QPI mode configuration retains the prior setting. Mode bit must be applied during dummy cycles. Number of dummy cycles in Table 6.11 includes number of mode bit cycles. If dummy cycles are configured as 6 cycles, data output will start right after mode bits and 4 additional dummy cycles are applied.

If the FRQIO/4FRQIO instruction in QPI mode is issued while an Erase, Program or Write cycle is in process (WIP=1) the instruction is ignored and will not have any effects on the current cycle.

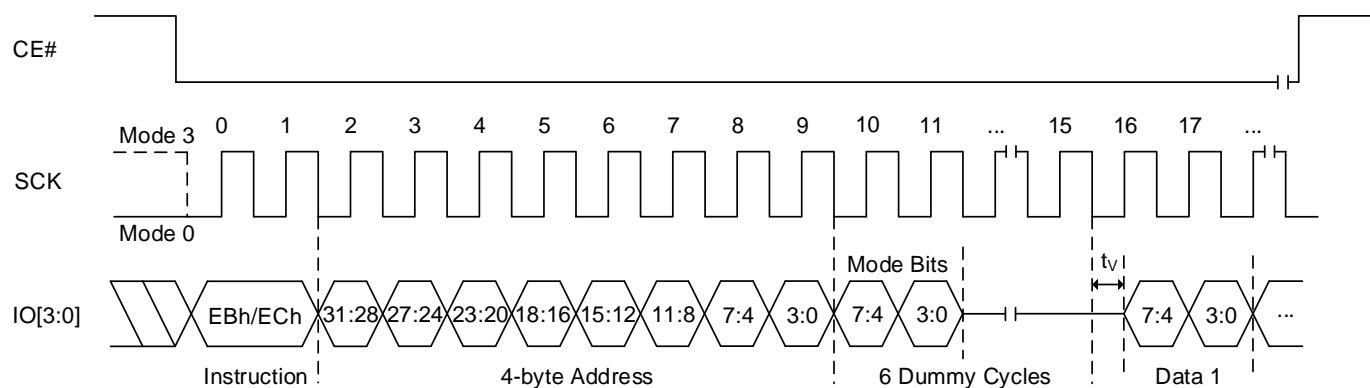
**Figure 8.19 Fast Read Quad I/O Sequence In QPI Mode (EBh [EXTADD=0], 3-byte address)**



**Notes:**

1. If the mode bits=AXh (where X is don't care), it can execute the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on Read Parameter setting. Detailed information in Table 6.11 Read Dummy Cycles.

**Figure 8.20 Fast Read Quad I/O Sequence In QPI Mode (EBh [EXTADD=1] or ECh, 4-byte address)**



**Notes:**

1. If the mode bits=AXh (where X is don't care), it can execute the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on Read Parameter setting. Detailed information in Table 6.11 Read Dummy Cycles.

### **8.10 PAGE PROGRAM OPERATION (PP, 02h or 4PP, 12h)**

The Page Program (PP/4PP) instruction allows up to 256/512 bytes data to be programmed into memory in a single operation. The destination of the memory to be programmed must be outside the protected memory area set by the Block Protection bits (BP3, BP2, BP1, BP0) or ASP. A PP/4PP instruction which attempts to program into a page that is write-protected will be ignored. Before the execution of PP/4PP instruction, the Write Enable Latch (WEL) must be enabled through a Write Enable (WREN) instruction.

- 02h (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- 02h (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- 12h is followed by a 4-byte address (A31-A0)

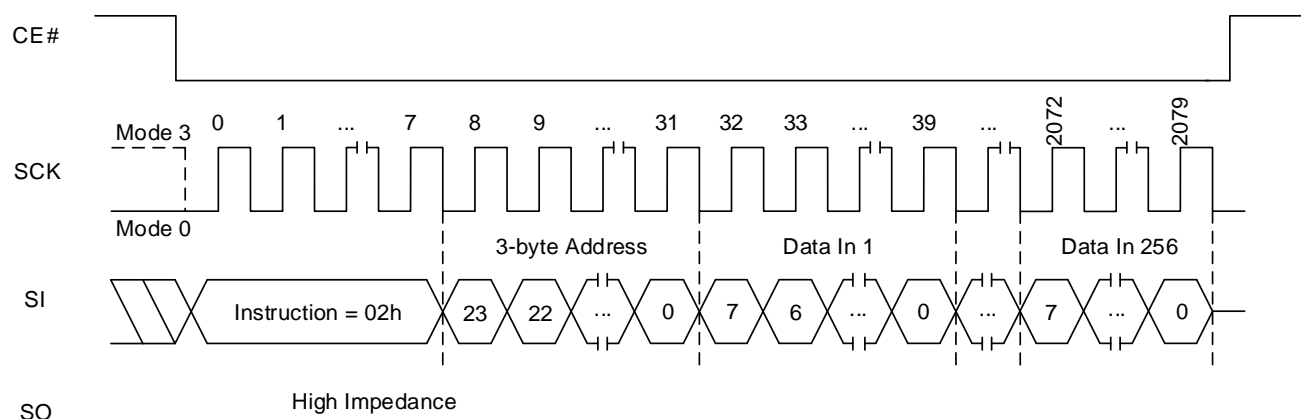
The PP/4PP instruction code, three or four address bytes as above and program data (1 to 256/512 bytes) are input via the SI line. Program operation will start immediately after the CE# is brought high, otherwise the PP/4PP instruction will not be executed. The internal control logic automatically handles the programming voltages and timing. The progress or completion of the program operation can be determined by reading the WIP bit. If the WIP bit is "1", the program operation is still in progress. If WIP bit is "0", the program operation has completed.

If more than 256/512 bytes data are sent to a device, the address counter rolls over within the same page, the previously latched data are discarded, and the last 256/512 bytes are kept to be programmed into the page. The starting byte can be anywhere within the page. When the end of the page is reached, the address will wrap around to the beginning of the same page. If the data to be programmed are less than a full page, the data of all other bytes on the same page will remain unchanged.

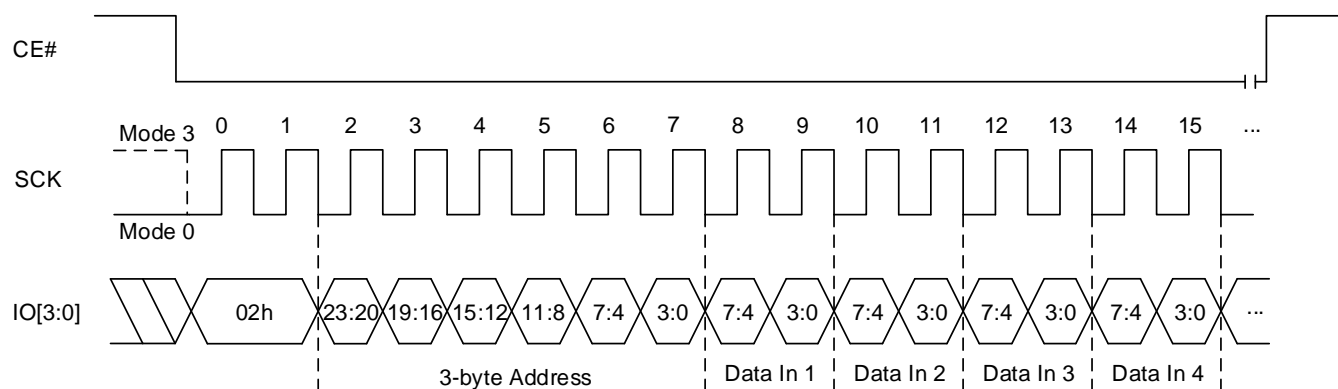
**Notes:**

**A program operation can alter "1"s into "0"s, but an erase operation is required to change "0"s back to "1"s.**

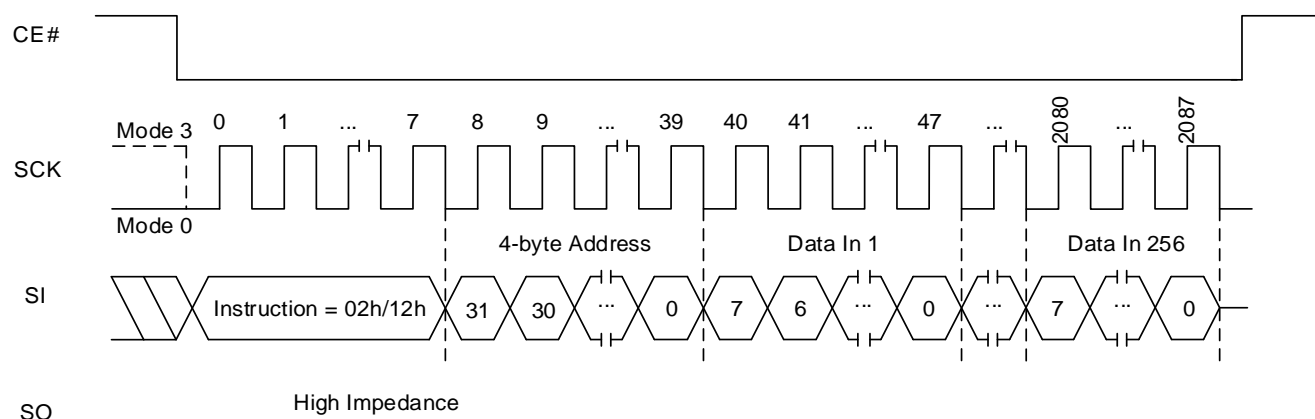
**Figure 8.21 Page Program Sequence In SPI Mode (02h [EXTADD=0], 3-byte address)**



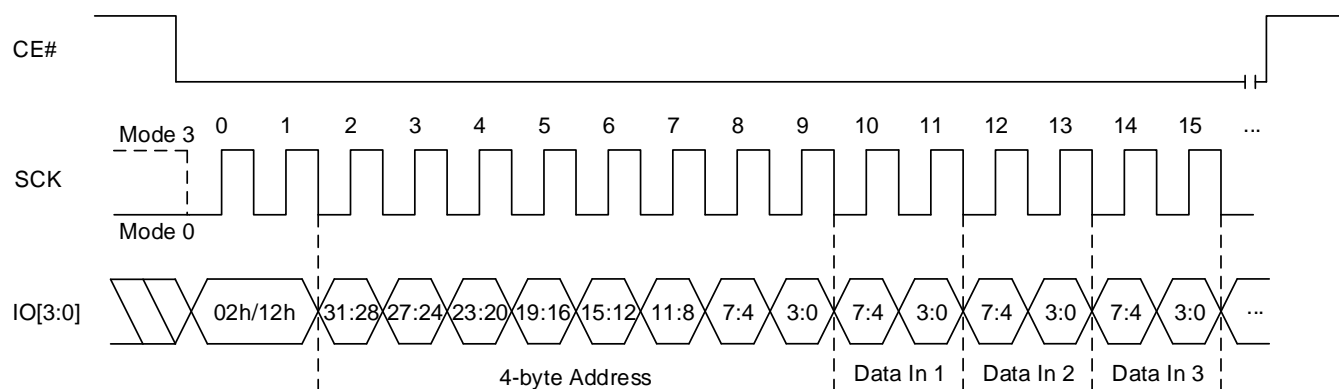
**Figure 8.22 Page Program Sequence In QPI Mode (02h [EXTADD=0], 3-byte address)**



**Figure 8.23 Page Program Sequence In SPI Mode (02h [EXTADD=1] or 12h, 4-byte address)**



**Figure 8.24 Page Program Sequence In QPI Mode (02h [EXTADD=1] or 12h, 4-byte address)**



### 8.11 QUAD INPUT PAGE PROGRAM OPERATION (PPQ, 32h/38h or 4PPQ, 34h/3Eh)

The Quad Input Page Program instruction allows up to 256/512 bytes data to be programmed into memory in a single operation with four pins (IO0, IO1, IO2 and IO3). The destination of the memory to be programmed must be outside the protected memory area set by the Block Protection (BP3, BP2, BP1, BP0) bits or ASP. A Quad Input Page Program instruction which attempts to program into a page that is write-protected will be ignored.

Before the execution of Quad Input Page Program instruction, the QE bit in the Status Register must be set to “1” and the Write Enable Latch (WEL) must be enabled through a Write Enable (WREN) instruction.

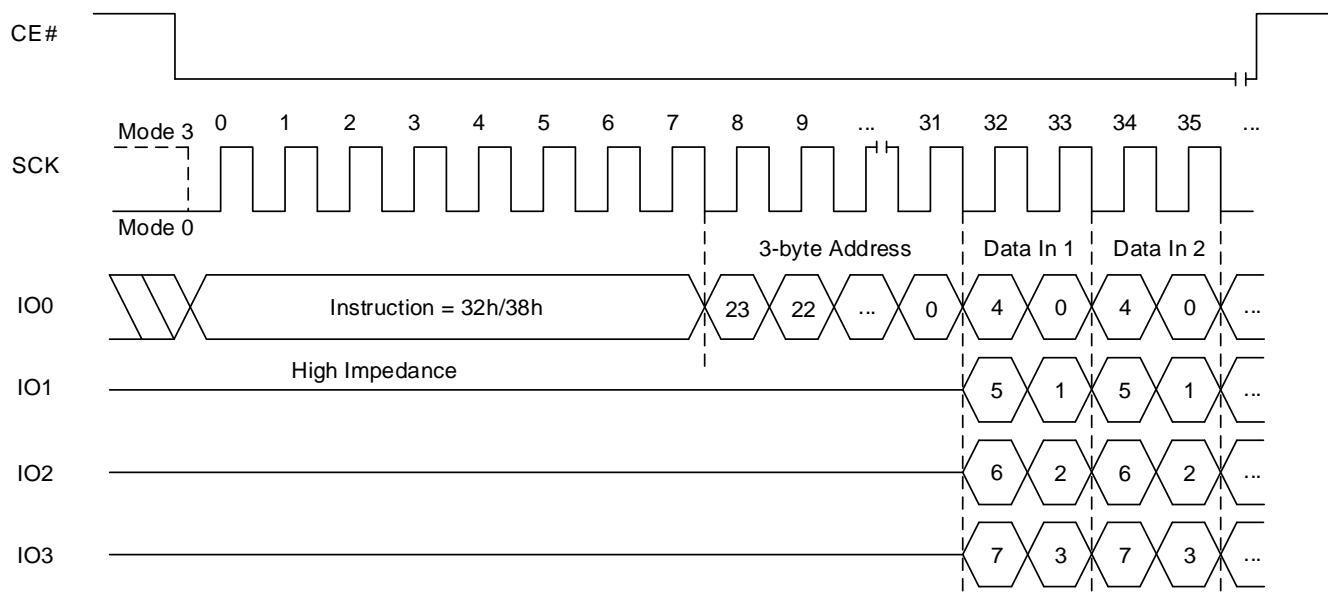
- 32h/38h (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- 32h/38h (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- 34h/3Eh is followed by a 4-byte address (A31-A0)

Program operation will start immediately after the CE# is brought high, otherwise the Quad Input Page Program instruction will not be executed. The internal control logic automatically handles the programming voltages and timing. The progress or completion of the program operation can be determined by reading the WIP bit. If the WIP bit is “1”, the program operation is still in progress. If WIP bit is “0”, the program operation has completed.

If more than 256/512 bytes data are sent to a device, the address counter rolls over within the same page, the previously latched data are discarded, and the last 256/512 bytes data are kept to be programmed into the page. The starting byte can be anywhere within the page. When the end of the page is reached, the address will wrap around to the beginning of the same page. If the data to be programmed are less than a full page, the data of all other bytes on the same page will remain unchanged.

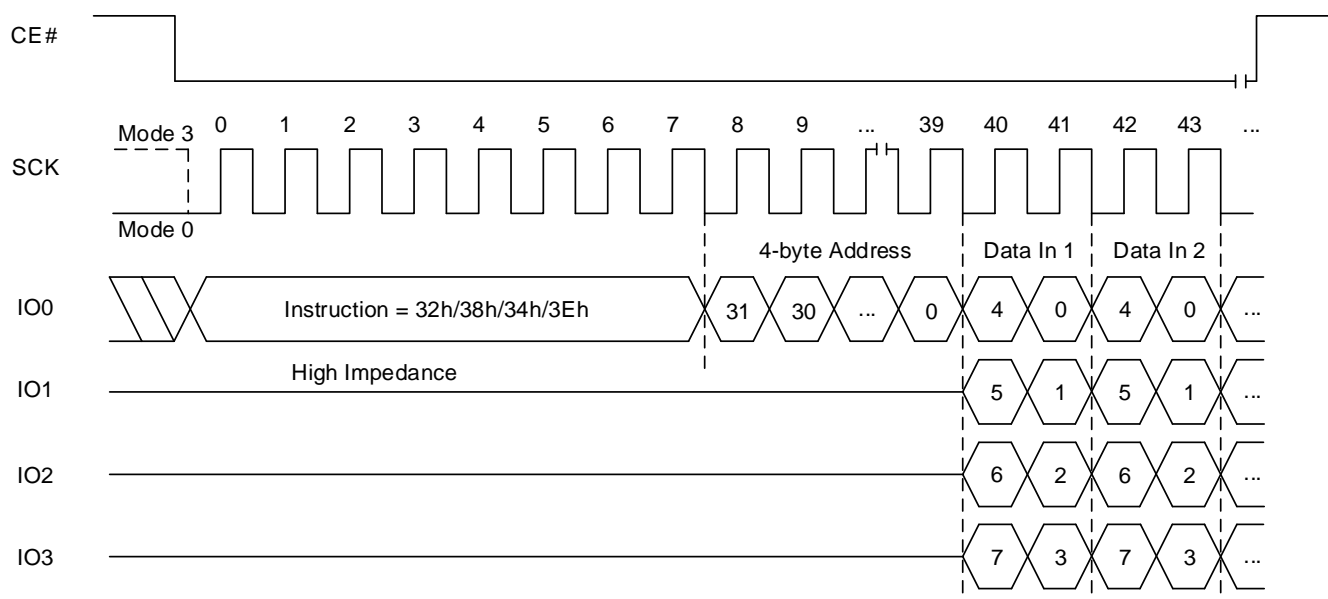
**Note:** A program operation can alter “1”s into “0”s, but an erase operation is required to change “0”s back to “1”s.

**Figure 8.25 Quad Input Page Program operation (32h/38h [EXTADD=0], 3-byte address)**





**Figure 8.26 Quad Input Page Program operation (32h/38h [EXTADD=1] or 34h/3Eh, 4-byte address)**



## 8.12 ERASE OPERATION

The Erase command sets all bits in the addressed sector or block to “1”s.

The memory array of the device is organized into uniform 4 Kbyte sectors or 32/64/256 Kbyte uniform blocks (a block consists of eight/sixteen adjacent sectors respectively).

Before a byte is reprogrammed, the sector or block that contains the byte must be erased (erasing sets bits to “1”). In order to erase the device, there are three erase instructions available: Sector Erase (SER), Block Erase (BER), and Chip Erase (CER). A sector erase operation allows any individual sector to be erased without affecting the data in other sectors. A block erase operation erases any individual block. A chip erase operation erases the whole memory array of a device. A sector erase, block erase, or chip erase operation can be executed prior to any programming operation.

### 8.13 SECTOR ERASE OPERATION (SER, D7h/20h or 4SER, 21h)

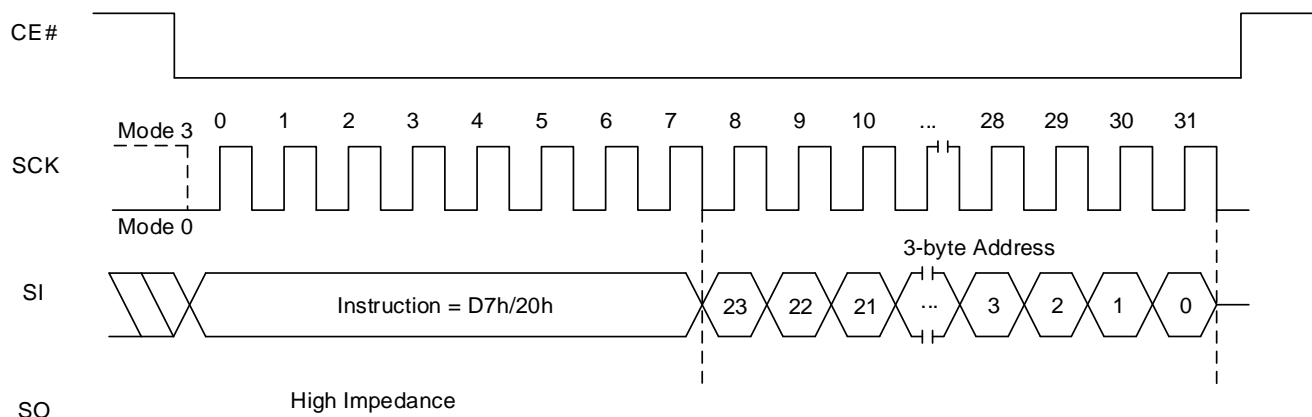
A Sector Erase (SER/4SER) instruction erases a 4 Kbyte sector before the execution of a SER/4SER instruction, the Write Enable Latch (WEL) must be set via a Write Enable (WREN) instruction. The WEL bit is automatically reset after the completion of Sector Erase operation.

- D7h/20h (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- D7h/20h (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- 21h is followed by a 4-byte address (A31-A0)

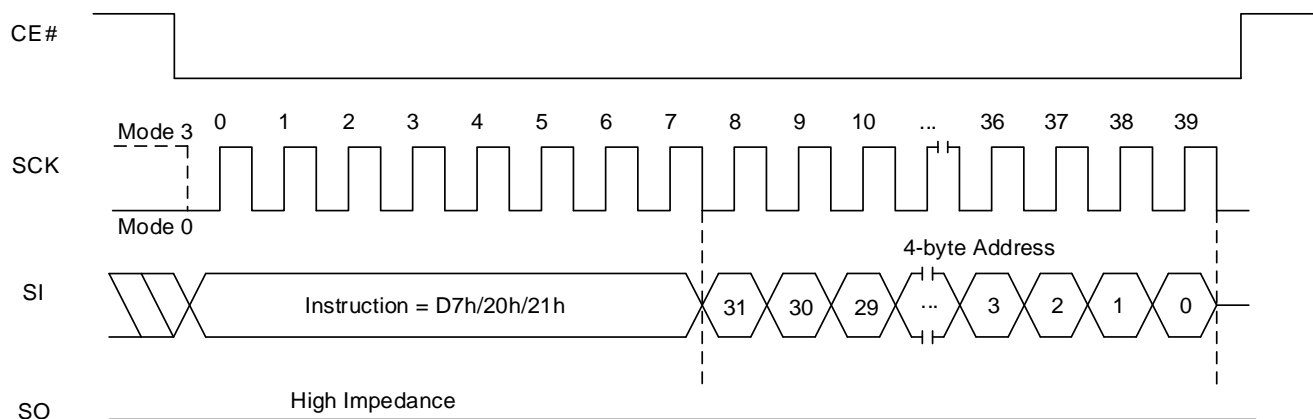
A SER/4SER instruction is entered, after CE# is pulled low to select the device and stays low during the entire instruction sequence. The SER/4SER instruction code, and three or four address bytes as above are input via SI. Erase operation will start immediately after CE# is pulled high. The internal control logic automatically handles the erase voltage and timing.

The progress or completion of the erase operation can be determined by reading the WIP bit. If the WIP bit is "1", the erase operation is still in progress. If the WIP bit is "0", the erase operation has been completed.

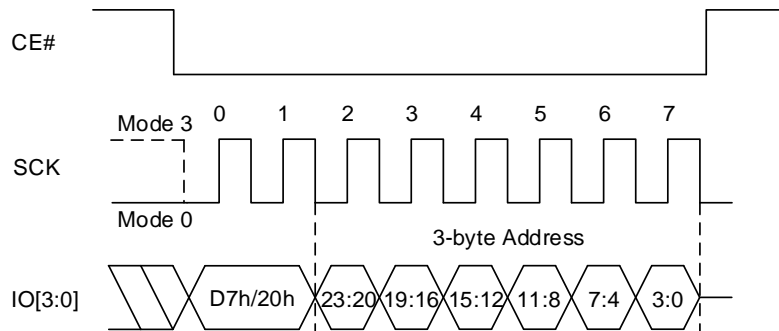
**Figure 8.27 Sector Erase Sequence In SPI Mode (D7h/20h [EXTADD=0], 3-byte address)**



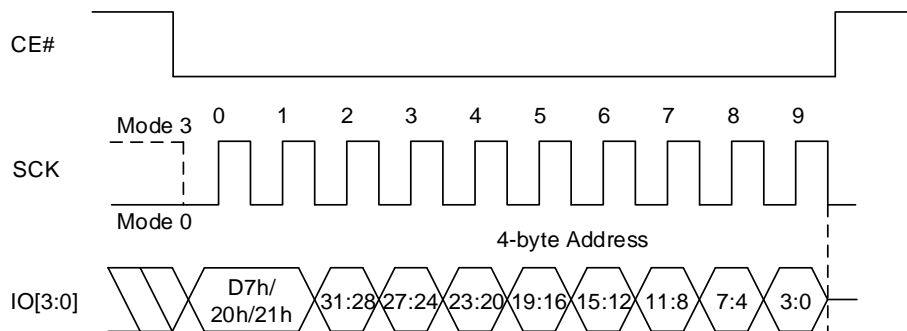
**Figure 8.28 Sector Erase Sequence In SPI Mode (D7h/20h [EXTADD=1] or 21h, 4-byte address)**



**Figure 8.29 Sector Erase Sequence In QPI Mode (D7h/20h [EXTADD=0], 3-byte address)**



**Figure 8.30 Sector Erase Sequence In QPI Mode (D7h/20h [EXTADD=1] or 21h, 4-byte address)**



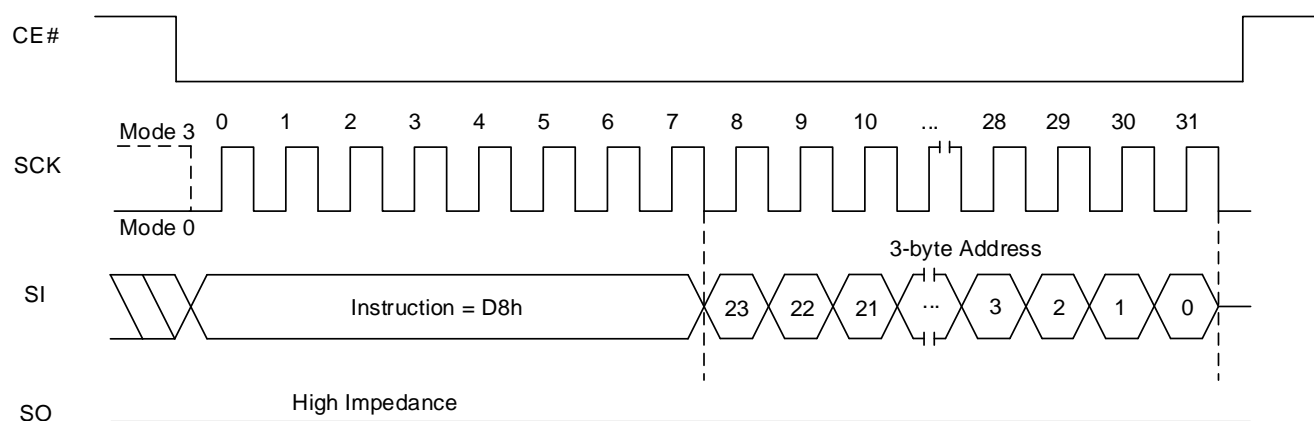
#### 8.14 BLOCK ERASE OPERATION (BER32K:52h or 4BER32K:5Ch, BER64K/256K:D8h or 4BER64K/256K:DCh)

A Block Erase (BER) instruction erases a 32/64 or 32/256 Kbyte block. Before the execution of a BER instruction, the Write Enable Latch (WEL) must be set via a Write Enable (WREN) instruction. The WEL is reset automatically after the completion of a block erase operation.

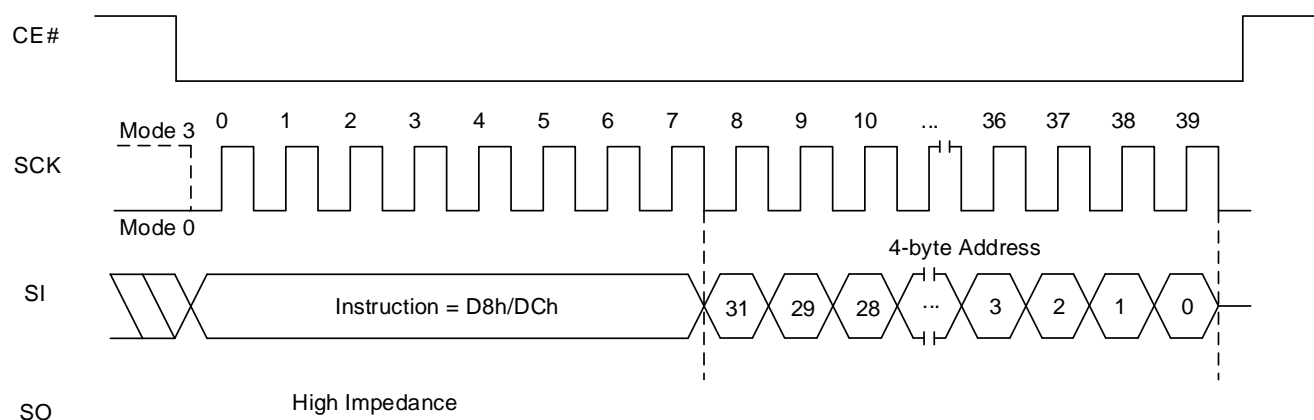
- 52h/D8h (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- 52h/D8h (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- 5Ch/DCh is followed by a 4-byte address (A31-A0)

The BER instruction code and three or four address bytes as above are input via SI. Erase operation will start immediately after the CE# is pulled high, otherwise the BER instruction will not be executed. The internal control logic automatically handles the erase voltage and timing.

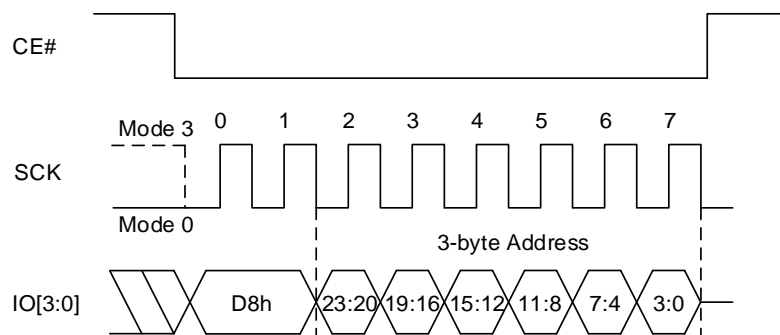
**Figure 8.31 Block Erase (64K/256K) Sequence In SPI Mode (D8h [EXTADD=0], 3-byte address)**



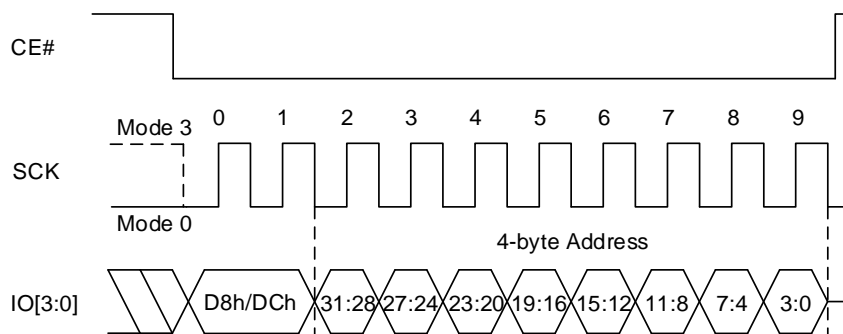
**Figure 8.32 Block Erase (64K/256K) Sequence In SPI Mode (D8h [EXTADD=1] or DCh, 4-byte address)**



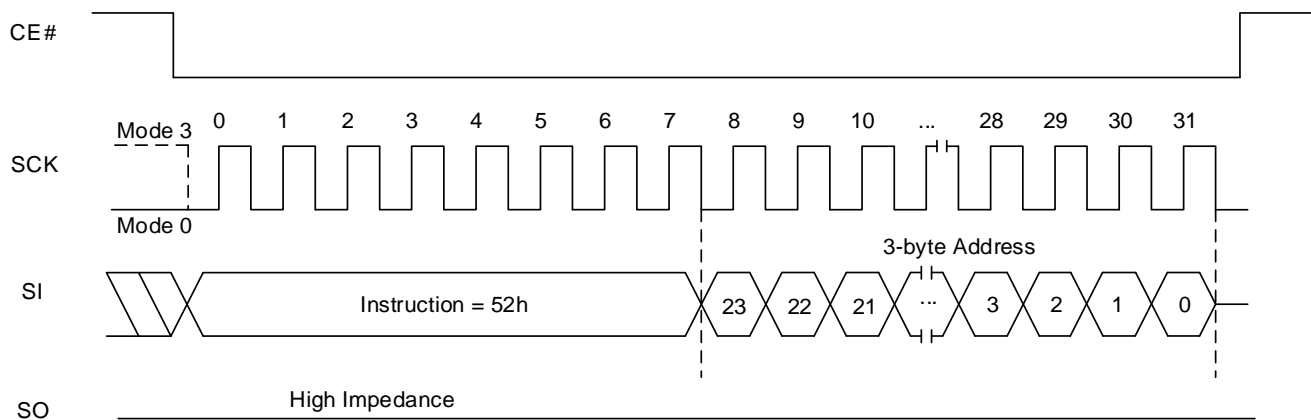
**Figure 8.33 Block Erase (64K/256K) Sequence In QPI Mode (D8h [EXTADD=0], 3-byte address)**



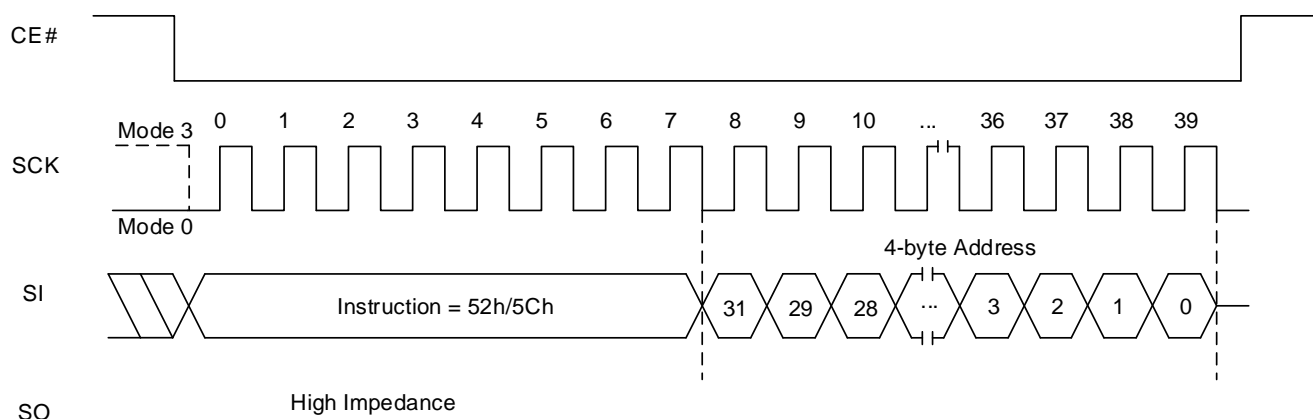
**Figure 8.34 Block Erase (64K/256K) Sequence In QPI Mode (D8h [EXTADD=1] or DCh, 4-byte address)**



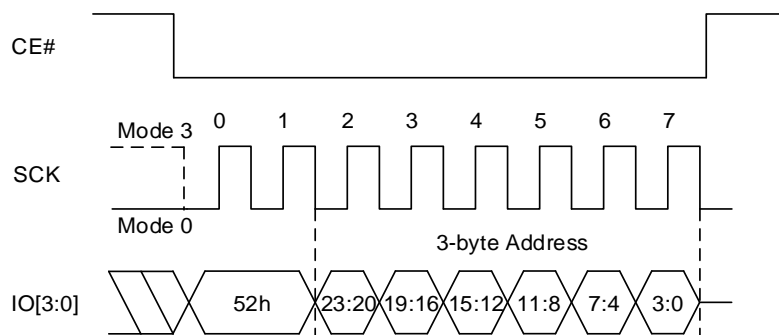
**Figure 8.35 Block Erase (32K) Sequence In SPI Mode (52h [EXTADD=0], 3-byte address)**



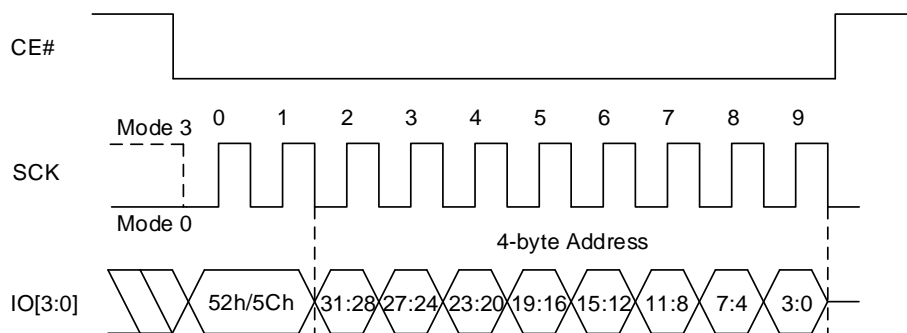
**Figure 8.36 Block Erase (32K) Sequence In SPI Mode (52h [EXTADD=1] or 5Ch, 4-byte address)**



**Figure 8.37 Block Erase (32K) Sequence In QPI Mode (52h [EXTADD=0], 3-byte address)**



**Figure 8.38 Block Erase (32K) Sequence In QPI Mode (52h [EXTADD=1] or 5Ch, 4-byte address)**



### 8.15 CHIP ERASE OPERATION (CER, C7h/60h)

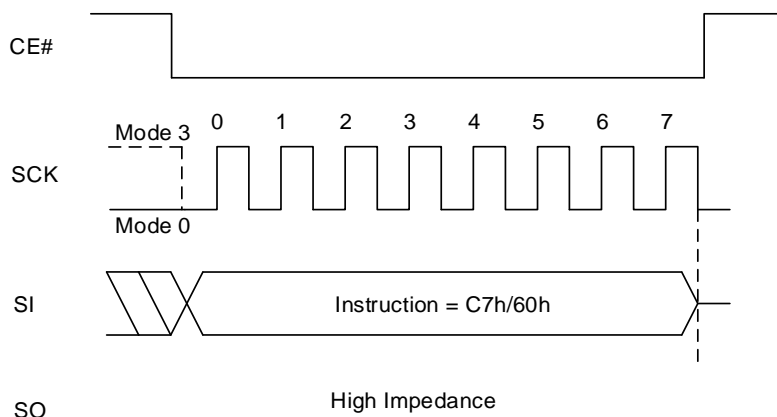
A Chip Erase (CER) instruction erases the entire memory array. Before the execution of CER instruction, the Write Enable Latch (WEL) must be set via a Write Enable (WREN) instruction. The WEL is automatically reset after completion of a chip erase operation.

The CER instruction code is input via the SI. Erase operation will start immediately after CE# is pulled high, otherwise the CER instruction will not be executed. The internal control logic automatically handles the erase voltage and timing.

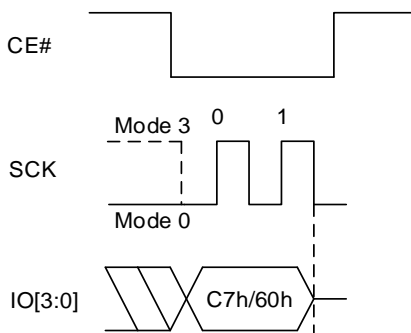
Chip Erase (CER) instruction can be executed only when Block Protection (BP3~BP0) bits are set to 0s. If the BP bits are not 0, the CER command is not executed and E\_ERR and PROT\_E are set.

Chip Erase (CER) instruction will skip sectors/blocks protected by ASP (DYB bits or PPB bits) and will **not** set E\_ERR and PROT\_E if sectors/blocks are protected by ASP only.

**Figure 8.39 Chip Erase Sequence In SPI Mode**



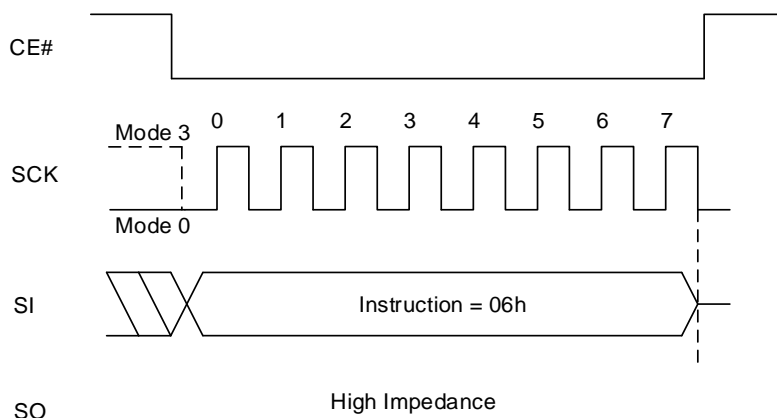
**Figure 8.40 Chip Erase Sequence In QPI Mode**



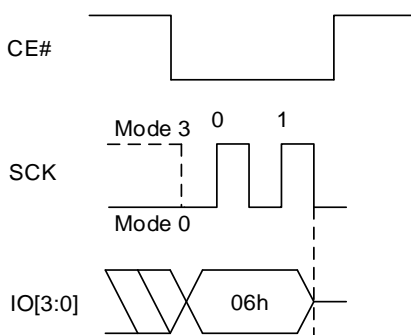
## 8.16 WRITE ENABLE OPERATION (WREN, 06h)

The Write Enable (WREN) instruction is used to set the Write Enable Latch (WEL) bit. The WEL bit is reset to the write-protected state after power-up. The WEL bit must be write enabled before any write operation, including Sector Erase, Block Erase, Chip Erase, Page Program, Program Information Row, Write Status Register, Write Function Register, Set Non-Volatile Read Register, Set Non-Volatile Extended Read Register, Write Autoboot Register operations, Set Volatile Read Register and Set Volatile Extended Read Register. The WEL bit will be reset to the write-protected state automatically upon completion of a write operation. The WREN instruction is required before any above operation is executed.

**Figure 8.41 Write Enable Sequence In SPI Mode**



**Figure 8.42 Write Enable Sequence In QPI Mode**

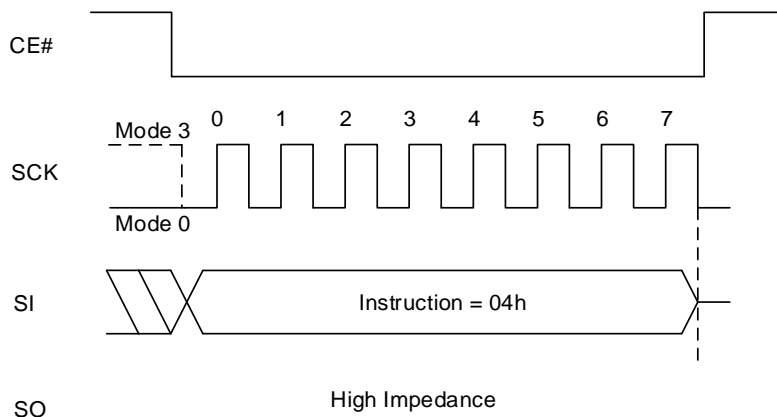




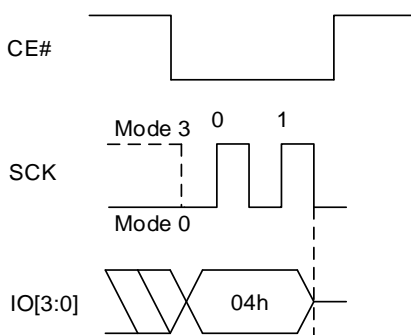
### 8.17 WRITE DISABLE OPERATION (WRDI, 04h)

The Write Disable (WRDI) instruction resets the WEL bit and disables all write instructions. The WRDI instruction is not required after the execution of a write instruction, since the WEL bit is automatically reset.

**Figure 8.43 Write Disable Sequence In SPI Mode**



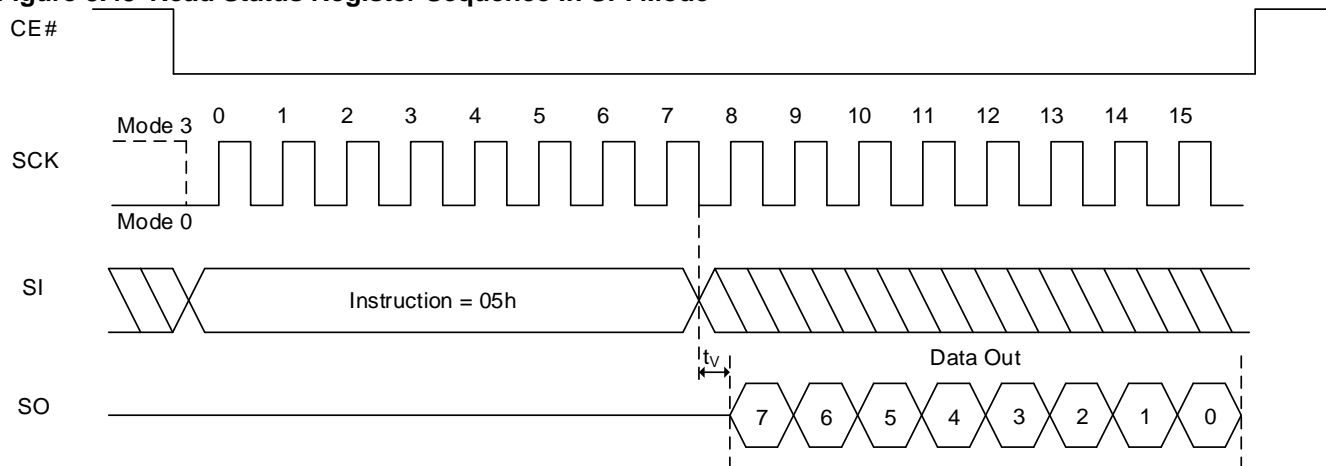
**Figure 8.44 Write Disable Sequence In QPI Mode**



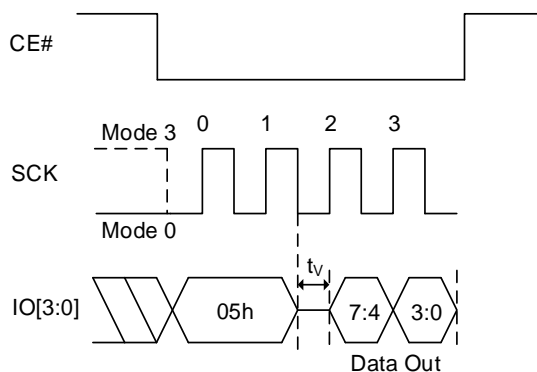
## 8.18 READ STATUS REGISTER OPERATION (RDSR, 05h)

The Read Status Register (RDSR) instruction provides access to the Status Register. During the execution of a program, erase or Write Status Register operation, the RDSR instruction can be used to check the progress or completion of an operation by reading the WIP bit.

**Figure 8.45 Read Status Register Sequence In SPI Mode**



**Figure 8.46 Read Status Register Sequence In QPI Mode**

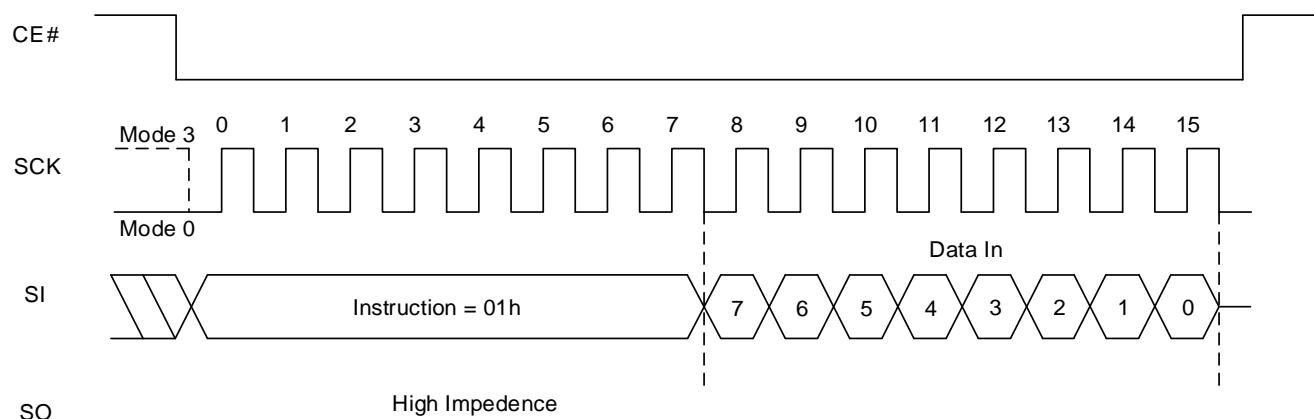


### 8.19 WRITE STATUS REGISTER OPERATION (WRSR, 01h)

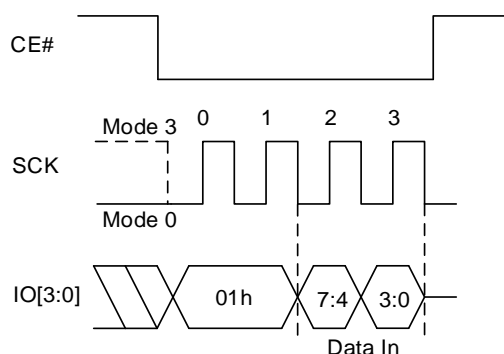
The Write Status Register (WRSR) instruction allows the user to enable or disable the block protection and Status Register write protection features by writing “0”s or “1”s into BP3, BP2, BP1, BP0, and SRWD bits. Also WRSR instruction allows the user to disable or enable quad operation by writing “0” or “1” into QE bit.

To write Status Register bits, a standard Write Enable (06h) instruction must previously have been executed for the device to accept Write Status Register (01h) instruction (Status Register bit WEL must equal 1).

**Figure 8.47 Write Status Register Sequence In SPI Mode**



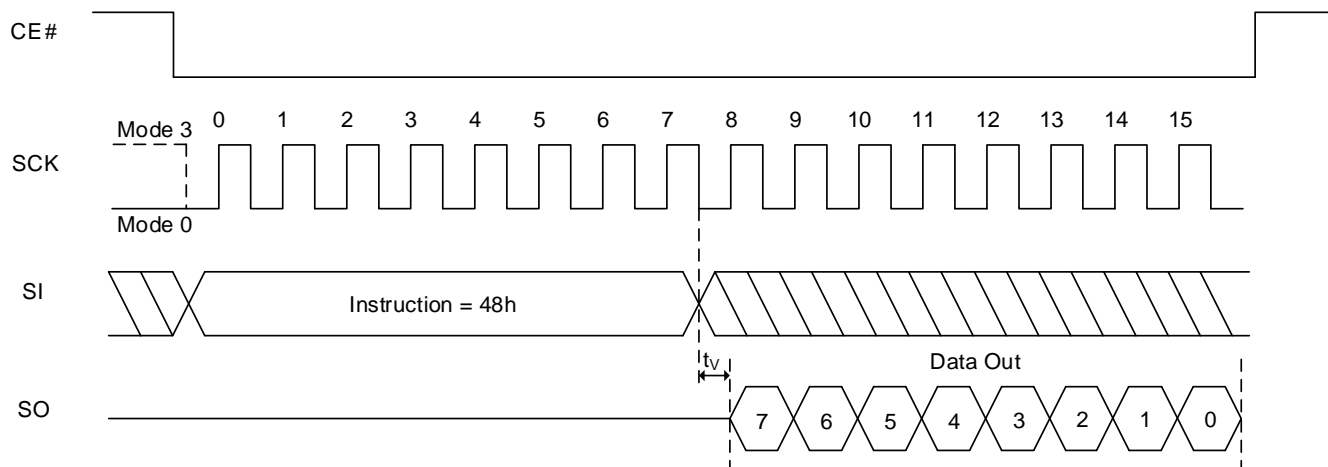
**Figure 8.48 Write Status Register QPI Sequence**



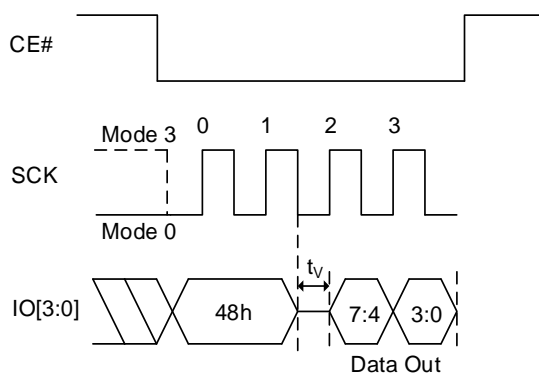
## 8.20 READ FUNCTION REGISTER OPERATION (RDFR, 48h)

The Read Function Register (RDFR) instruction provides access to the Function Register. Refer to Table 6.6 Function Register Bit Definition for more detail.

**Figure 8.49 Read Function Register Sequence In SPI Mode**



**Figure 8.50 Read Function Register QPI Sequence**

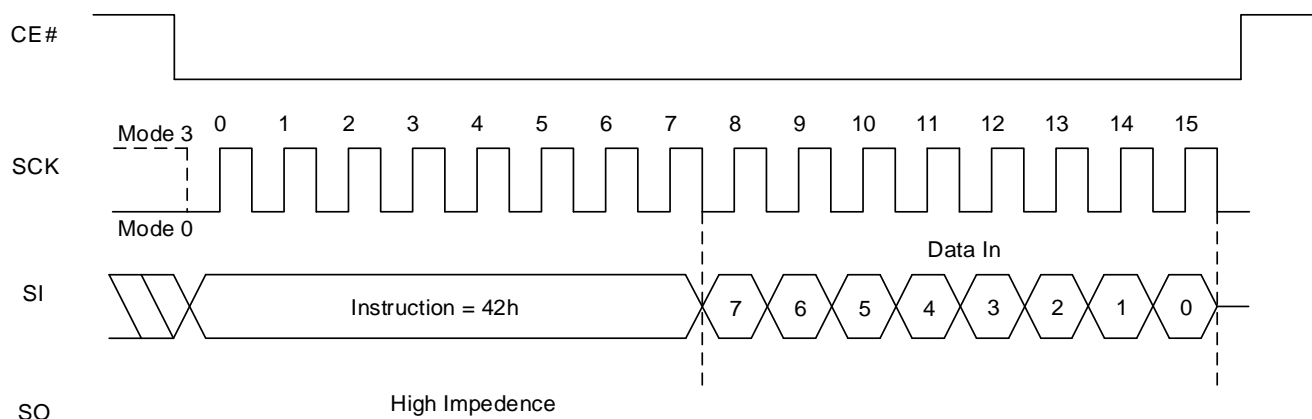


## 8.21 WRITE FUNCTION REGISTER OPERATION (WRFR, 42h)

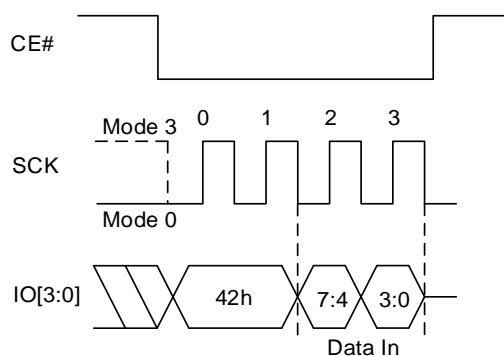
The Write Function Register (WRFR) instruction allows Information Row Lock bits (IRL3~IRL0) can be set to “1” individually by WRFR instruction in order to lock Information Row.

Since IRL bits are OTP, once they are set to “1”, they cannot be set back to “0” again

**Figure 8.51 Write Function Register Sequence In SPI Mode**



**Figure 8.52 Write Function Register QPI Sequence**

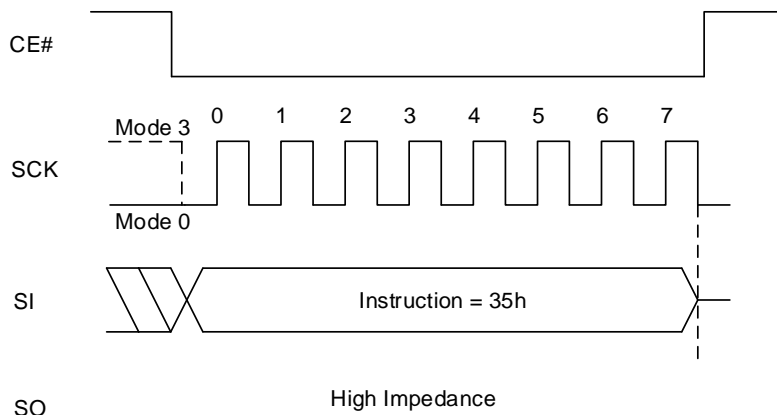


## 8.22 ENTER QUAD PERIPHERAL INTERFACE (QPI) MODE OPERATION (QPIEN,35h; QPIDI,F5h)

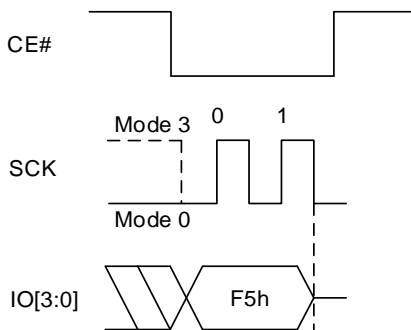
The Enter QPI (QPIEN) instruction, 35h, enables the Flash device for QPI bus operation. Upon completion of the instruction, all instructions thereafter will be 4-bit multiplexed input/output until a power cycle or an Exit QPI instruction is sent to device.

The Exit QPI instruction, F5h, resets the device to 1-bit SPI protocol operation. To execute an Exit QPI operation, the host drives CE# low, sends the Exit QPI command cycle, then drives CE# high. The device just accepts QPI (2 clocks) command cycles.

**Figure 8.53 Enter Quad Peripheral Interface (QPI) Mode Sequence**



**Figure 8.54 Exit Quad Peripheral Interface (QPI) Mode Sequence**



## 8.23 PROGRAM/ERASE SUSPEND & RESUME

The device allows the interruption of Sector Erase, Block Erase, or Page Program operations to conduct other operations. 75h/B0h command for suspend and 7Ah/30h for resume will be used. (SPI/QPI all acceptable) Function Register bit2 (PSUS) and bit3 (ESUS) are used to check whether or not the device is in suspend mode.

Suspend to read ready timing: 100µs

Resume to another suspend timing: 400µs

### PROGRAM/ERASE SUSPEND DURING SECTOR-ERASE OR BLOCK-ERASE (PERSUS 75h/B0h)

The Program/Erase Suspend allows the interruption of Sector Erase and Block Erase operations. After the Program/Erase Suspend, program, read related, resume and reset commands can be accepted. But Write Status Register instruction (01h) and Erase instructions are not allowed during Erase Suspend.

Refer to Table 8.8 for more detail.

To execute the Program/Erase Suspend operation, the host drives CE# low, sends the Program/Erase Suspend command cycle (75h/B0h), then drives CE# high. The Function Register indicates that the Erase has been suspended by changing the ESUS bit from “0” to “1”, but the device will not accept another command until it is ready. To determine when the device will accept a new command, poll the WIP bit or wait the specified time  $t_{SUS}$ . When ESUS bit is issued, the Write Enable Latch (WEL) bit will be reset.

### PROGRAM/ERASE SUSPEND DURING PAGE PROGRAMMING (PERSUS 75h/B0h)

The Program/Erase Suspend allows the interruption of all array Program operations. After the Program/Erase Suspend command, WEL bit will be disabled, therefore only read related, write Volatile register, resume and reset commands can be accepted.

Write Status Register instruction (01h) and Program instructions are not allowed during Program Suspend. Refer to Table 8.8 for more detail.

To execute the Program/Erase Suspend operation, the host drives CE# low, sends the Program/Erase Suspend command cycle (75h/B0h), then drives CE# high. The Function Register indicates that the programming has been suspended by changing the PSUS bit from “0” to “1”, but the device will not accept another command until it is ready. To determine when the device will accept a new command, poll the WIP bit or wait the specified time  $t_{SUS}$ .

### PROGRAM/ERASE RESUME (PERRSM 7Ah/30h)

The Program/Erase Resume restarts the Program or Erase command that was suspended, and changes the suspend status bit in the Function Register (ESUS or PSUS bits) back to “0”. To execute the Program/Erase Resume operation, the host drives CE# low, sends the Program/Erase Resume command cycle (7Ah/30h), then drives CE# high. A cycle is two nibbles long, most significant nibble first. To determine if the internal, self-timed Write operation completed, poll the WIP bit, or wait the specified time  $t_{SE}$ ,  $t_{BE}$  or  $t_{PP}$  for Sector Erase, Block Erase, or Page Programming, respectively. The total write time before suspend and after resume will not exceed the uninterrupted write times  $t_{SE}$ ,  $t_{BE}$  or  $t_{PP}$ .

**Table 8.8 Instructions accepted during Suspend**

Operation Suspended	Instruction Allowed		
	Name	Hex Code	Operation
Program or Erase	NORD	03h	Normal Read Mode
Program or Erase	4NORD	13h	4-byte Address Normal Read Mode
Program or Erase	FRD	0Bh	Fast Read Mode
Program or Erase	4FRD	0Ch	4-byte Address Fast Read Mode
Program or Erase	FRDIO	BBh	Fast Read Dual I/O
Program or Erase	4FRDIO	BCh	4-byte Address Fast Read Dual I/O
Program or Erase	FRDO	3Bh	Fast Read Dual Output
Program or Erase	4FRDO	3Ch	4-byte Address Fast Read Dual Output
Program or Erase	FRQIO	EBh	Fast Read Quad I/O
Program or Erase	4FRQIO	ECh	4-byte Address Fast Read Quad I/O
Program or Erase	FRQO	6Bh	Fast Read Quad Output
Program or Erase	4FRQO	6Ch	4-byte Address Fast Read Quad Output
Program or Erase	FRDTR	0Dh	Fast Read DTR Mode
Program or Erase	4FRDTR	0Eh	4-byte Address Fast Read DTR Mode
Program or Erase	FRDDTR	BDh	Fast Read Dual I/O DTR
Program or Erase	4FRDDTR	BEh	4-byte Address Fast Read Dual I/O DTR
Program or Erase	FRQDTR	EDh	Fast Read Quad I/O DTR
Program or Erase	4FRQDTR	EEh	4-byte Address Fast Read Quad I/O DTR
Program or Erase	WREN	06h	Write Enable
Program or Erase	WRDI	04hh	Write Disable
Erase	PP	02h	Serial Input Page Program
Erase	4PP	12h	4-byte Address Serial Input Page Program
Erase	PPQ	32h/38h	Quad Input Page Program
Erase	4PPQ	34h/3Eh	4-byte Address Quad Input Page Program
Program or Erase	RDSR	05h	Read Status Register
Program or Erase	RDFR	48h	Read Function Register
Program or Erase	CLERP	82h	Clear Extended Read Register

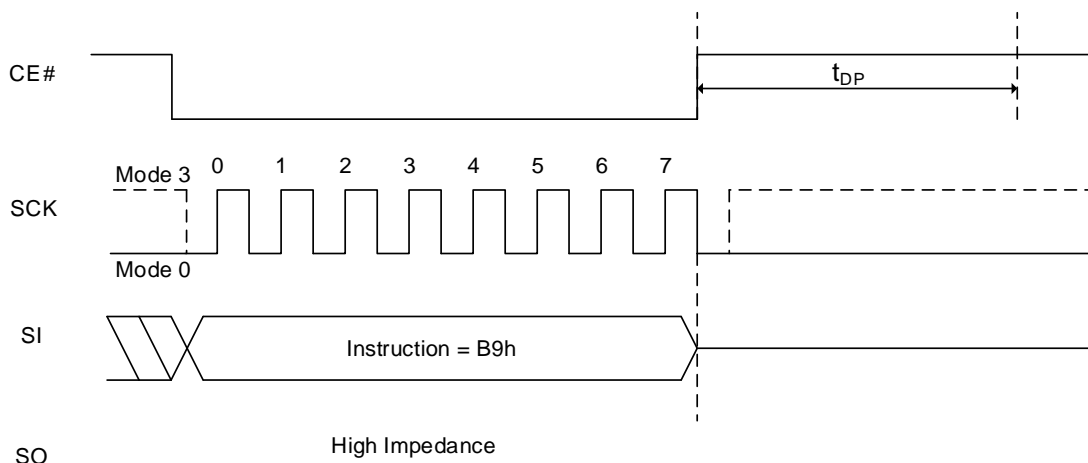


Operation Suspended	Instruction Allowed		
	Name	Hex Code	Operation
Program or Erase	PERRSM	7Ah/30h	Program/Erase Resume
Erase	PERSUS	75h/B0h	Program/Erase Suspend
Program or Erase	RDID	ABh	Read Manufacturer and Product ID
Program or Erase	SRPV	C0/63h	Set Read Parameters (Volatile)
Program or Erase	SERPv	83h	Set Extended Read Parameters (Volatile)
Program or Erase	RDRP	61h	Read Read Parameters (Volatile)
Program or Erase	RDERP	81h	Read Extended Read Parameters (Volatile)
Program or Erase	RDJDID	9Fh	Read Manufacturer and Product ID by JEDEC ID Command
Program or Erase	RDMDID	90h	Read Manufacturer and Device ID
Program or Erase	RDJDIDQ	AFh	Read JEDEC ID QPI mode
Program or Erase	RDUID	4Bh	Read Unique ID Number
Program or Erase	RDSFDP	5Ah	SFDP Read
Program or Erase	WRVDLR	4Ah	Write Volatile DLP Register

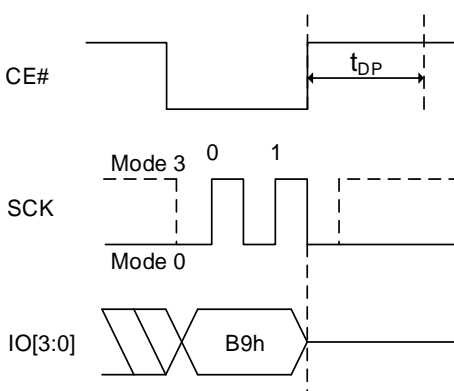
## 8.24 ENTER DEEP POWER DOWN (DP, B9h)

The Deep Power-down (DP) instruction is for setting the device on the minimizing the power consumption (enter into Power-down mode). During this mode, standby current is reduced from  $I_{sb1}$  to  $I_{sb2}$ . While in the Power-down mode, the device is not active and all Write/Program/Erase instructions are ignored. The instruction is initiated by driving the CE# pin low and shifting the instruction code into the device. The CE# pin must be driven high after the instruction has been latched, or Power-down mode will not engage. Once CE# pin driven high, the Power-down mode will be entered within the time duration of  $t_{DP}$ . While in the Power-down mode only the Release from Power-down/RDID instruction, which restores the device to normal operation, will be recognized. All other instructions are ignored, including the Read Status Register instruction which is always available during normal operation. Ignoring all but one instruction makes the Power Down state a useful condition for securing maximum write protection. It is available in both SPI and QPI mode.

**Figure 8.55 Enter Deep Power Down Mode Sequence In SPI Mode**



**Figure 8.56 Enter Deep Power Down Mode Sequence In QPI Mode**

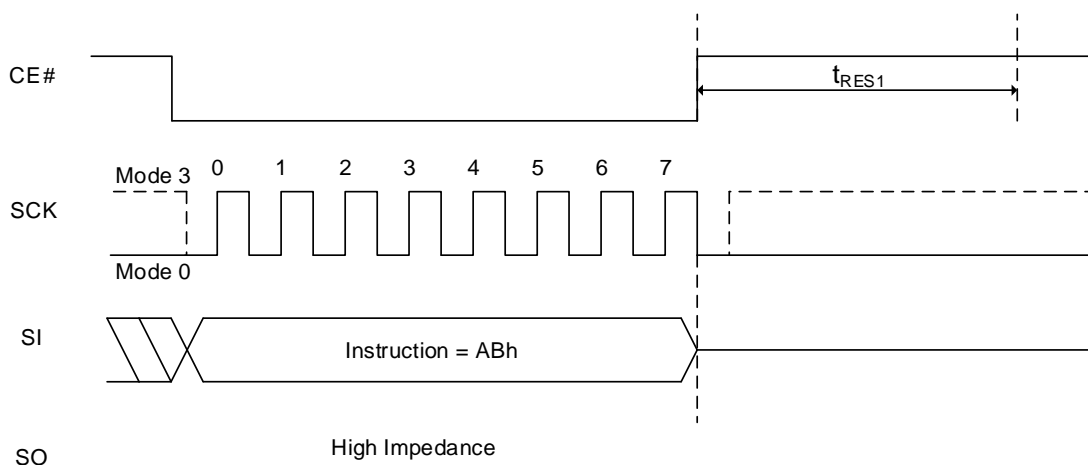


## 8.25 RELEASE DEEP POWER DOWN (RDPD, ABh)

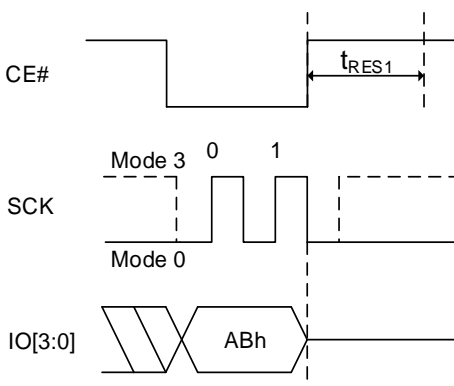
The Release Deep Power-down/Read Device ID instruction is a multi-purpose command. To release the device from the deep power-down mode, the instruction is issued by driving the CE# pin low, shifting the instruction code into the device and driving CE# high.

Releasing the device from Power-down mode will take the time duration of  $t_{RES1}$  before normal operation is restored and other instructions are accepted. The CE# pin must remain high during the  $t_{RES1}$  time duration. If the Release Deep Power-down/RDID instruction is issued while an Erase, Program or Write cycle is in progress (WIP=1) the instruction is ignored and will not have any effects on the current cycle.

**Figure 8.57 Release Power Down Mode Sequence In SPI Mode**



**Figure 8.58 Release Power Down Mode Sequence In QPI Mode**



## 8.26 SET READ PARAMETERS OPERATION (SRPNV: 65h, SRPV: C0h/63h)

### Set Read Parameter Bits

This device supports configurable burst length and dummy cycles in both SPI and QPI mode by setting three bits (P2, P1, P0) and four bits (P6, P5, P4, P3) within the Read Register, respectively. To set those bits the SRPNV and SRPV operation instruction are used. Details regarding burst length and dummy cycles can be found in Table 6.9, Table 6.10, and Table 6.11.

SRPNV is used to set the non-volatile Read Register, while SRPV is used to set the volatile Read Register.

To write non-volatile Read Parameter bits, a standard Write Enable (06h) instruction must previously have been executed for the device to accept SRPNV(65h) instruction (Status Register bit WEL must equal "1").

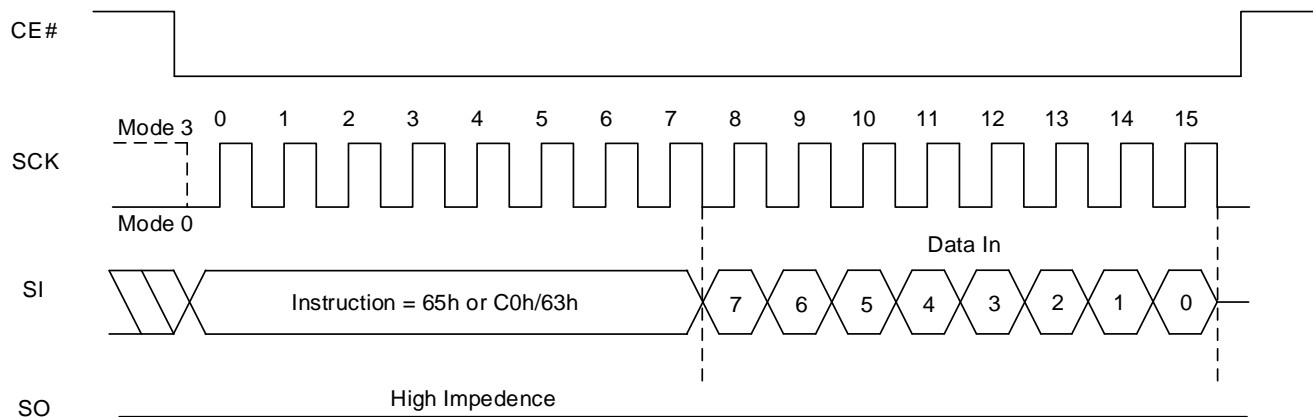
To write volatile Read Parameter bits (SRPV), 63h or C0h command can be used.

When using 63h instruction, a standard Write Enable (06h) instruction must previously have been executed for the device to accept SRPV (63h) instruction (Status Register bit WEL must equal "1").

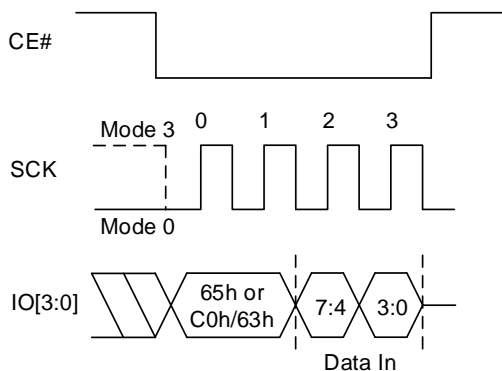
But C0h instruction does not require a standard Write Enable (06h) operation. (Status Register bit WEL remains "0").

**Note: When SRPNV is executed, the volatile Read Register is set as well as the non-volatile Read Register.**

**Figure 8.59 Set Read Parameters Sequence In SPI Mode**



**Figure 8.60 Set Read Parameters Sequence In QPI Mode**



### **Read with “8/16/32/64-Byte Wrap Around”**

The device is capable of burst read with wrap around in both SPI and QPI mode. The size of burst length is configurable by using P0, P1, and P2 bits in Read Register. P2 bit (Wrap enable) enables the burst mode feature. P0 and P1 define the size of burst. Burst lengths of 8, 16, 32, and 64 bytes are supported. By default, address increases by one up through the entire array. By setting the burst length, the data being accessed can be limited to the length of burst boundary within a 256 byte page. The first output will be the data at the initial address which is specified in the instruction. Following data will come out from the next address within the burst boundary. Once the address reaches the end of boundary, it will automatically move to the first address of the boundary. CE# high will terminate the command.

For example, if burst length of 8 and initial address being applied is 0h, following byte output will be from address 00h and continue to 01h,...,07h, 00h, 01h... until CE# terminates the operation. If burst length of 8 and initial address being applied is FEh(254d), following byte output will be from address FEh and continue to FFh, F8h, F9h, FAh, FBh, FCh, FDh, and repeat from FEh until CE# terminates the operation.

The commands, “SRPV (65h) or SRPNV (C0h or 63h)”, are used to configure the burst length. If the following data input is one of “00h”, “01h”, “02h”, and “03h”, the device will be in default operation mode. It will be continuous burst read of the whole array. If the following data input is one of “04h”, “05h”, “06h”, and “07h”, the device will set the burst length as 8, 16, 32 and 64, respectively.

To exit the burst mode, another “C0h or 63h” command is necessary to set P2 to 0. Otherwise, the burst mode will be retained until either power down or reset operation. To change burst length, another “C0h or 63h” command should be executed to set P0 and P1 (Detailed information in Table 6.9 Burst Length Data). All read commands will operate in burst mode once the Read Register is set to enable burst mode.

Refer to Figure 8.59 and Figure 8.60 for instruction sequence.

## 8.27 SET EXTENDED READ PARAMETERS OPERATION (SERPNV: 85h, SERPV: 83h)

### Set Read Operational Driver Strength

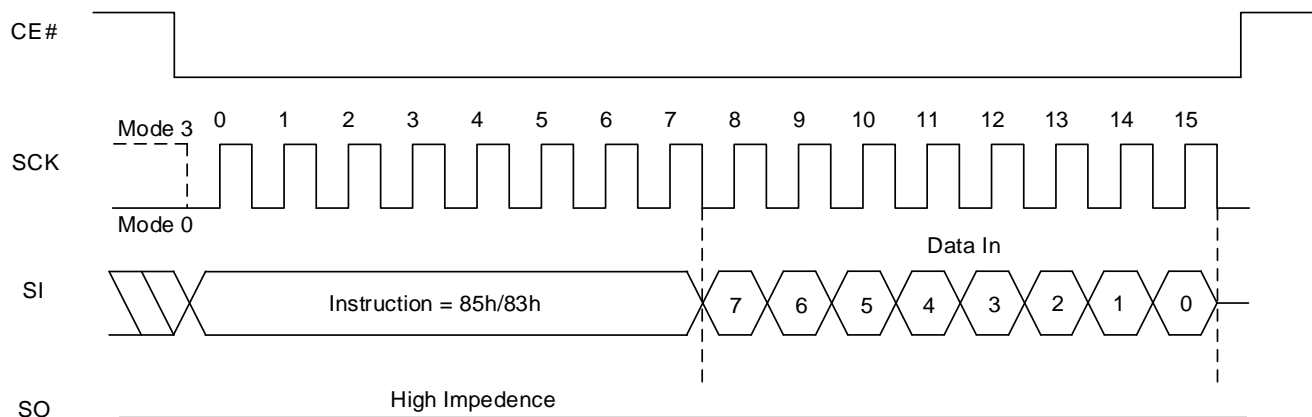
This device supports configurable Operational Driver Strength in both SPI and QPI modes by setting three bits (ODS0, ODS1, ODS2) within the Extended Read Register. To set the ODS bits the SERPNV and SERPV operation instructions are required. The device's driver strength can be reduced as low as 12.50% of full drive strength. Details regarding the driver strength can be found in Table 6.14.

SERPNV is used to set the non-volatile Extended Read register, while SERPV is used to set the volatile Extended Read register.

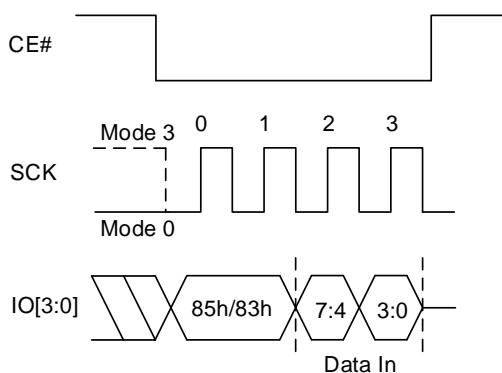
#### Notes:

1. The default driver strength is set to 50%.
2. When SERPNV is executed, the volatile Read Extended Register is set as well as the non-volatile Read Extended Register.

**Figure 8.61 Set Extended Read Parameters Sequence In SPI Mode**



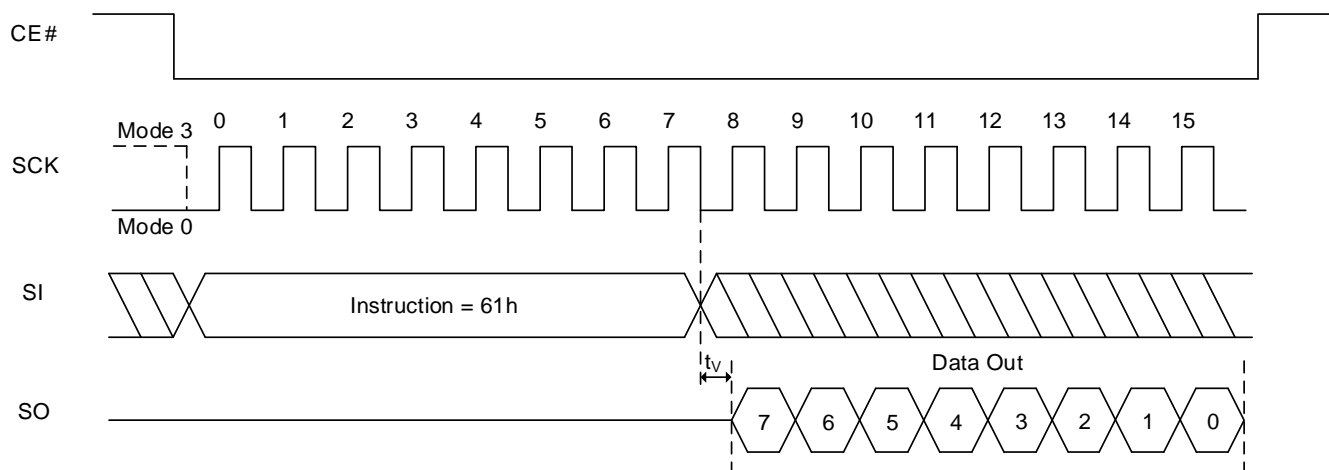
**Figure 8.62 Set Extended Read Parameters Sequence In QPI Mode**



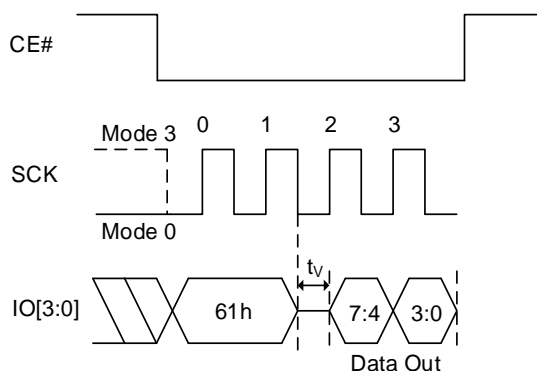
## 8.28 READ READ PARAMETERS OPERATION (RDRP, 61h)

Prior to, or after setting Read Register, the data of the Read Register can be confirmed by the RDRP command. The instruction is only applicable for the volatile Read Register, not for the non-volatile Read Register.

**Figure 8.63 Read Read Parameters Sequence In SPI Mode**



**Figure 8.64 Read Read Parameters Sequence In QPI Mode**

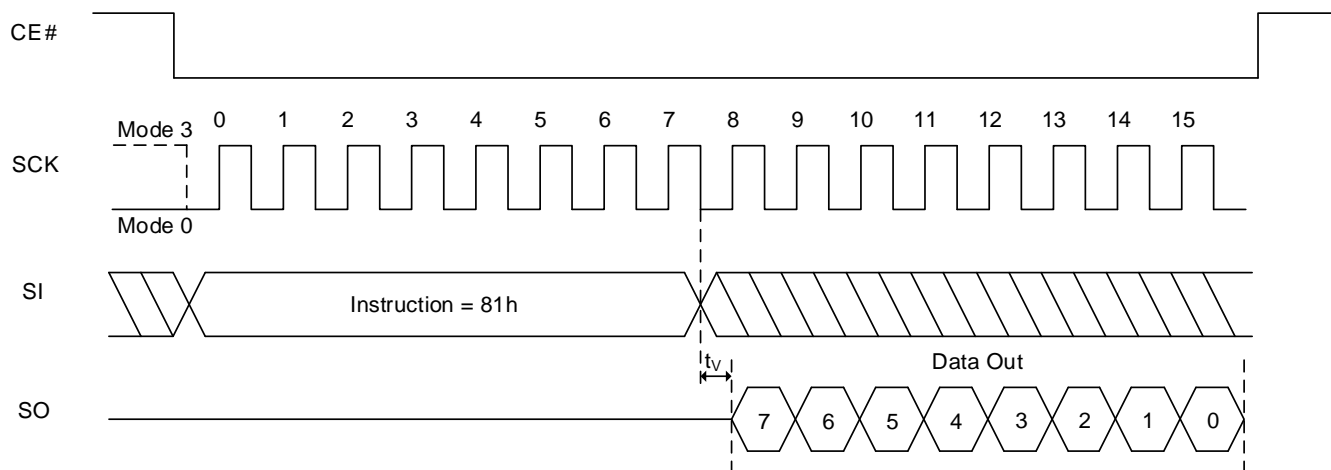


### 8.29 READ EXTENDED READ PARAMETERS OPERATION (RDERP, 81h)

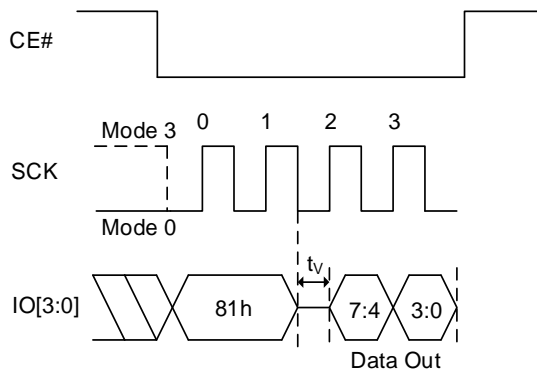
Prior to, or after setting Extended Read Register, the data of the Extended Read Register can be confirmed by the RDERP command. The instruction is only applicable for the volatile Extended Read Register, not for the non-volatile Extended Read Register.

During the execution of a Program, Erase or Write Non-Volatile Register operation, the RDERP instruction will be executed, which can be used to check the progress or completion of an operation by reading the WIP bit.

**Figure 8.65 Read Extended Read Parameters Sequence In SPI Mode**



**Figure 8.66 Read Extended Read Parameters Sequence In QPI Mode**

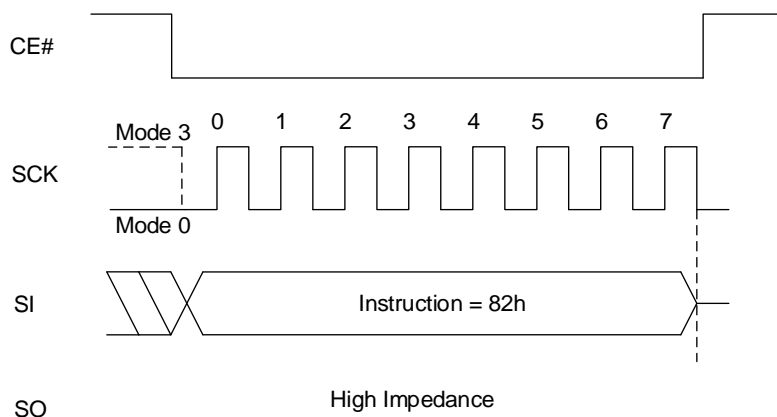




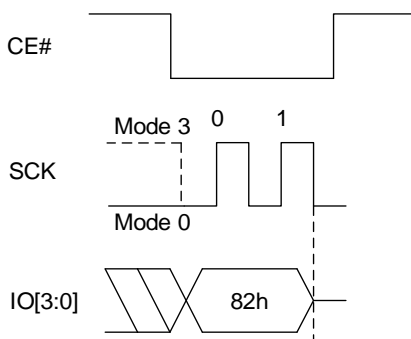
### 8.30 CLEAR EXTENDED READ REGISTER OPERATION (CLERP, 82h)

A Clear Extended Read Register (CLERP) instruction clears PROT\_E, P\_ERR, and E\_ERR error bits in the Extended Read Register to “0” when the error bits are set to “1”. Once the error bits are set to “1”, they remain set to “1” until they are cleared to “0” with a CLERP command.

**Figure 8.67 Clear Extended Read Register Sequence In SPI Mode**



**Figure 8.68 Clear Extended Read Register Sequence In QPI Mode**



### 8.31 READ PRODUCT IDENTIFICATION (RDID, ABh)

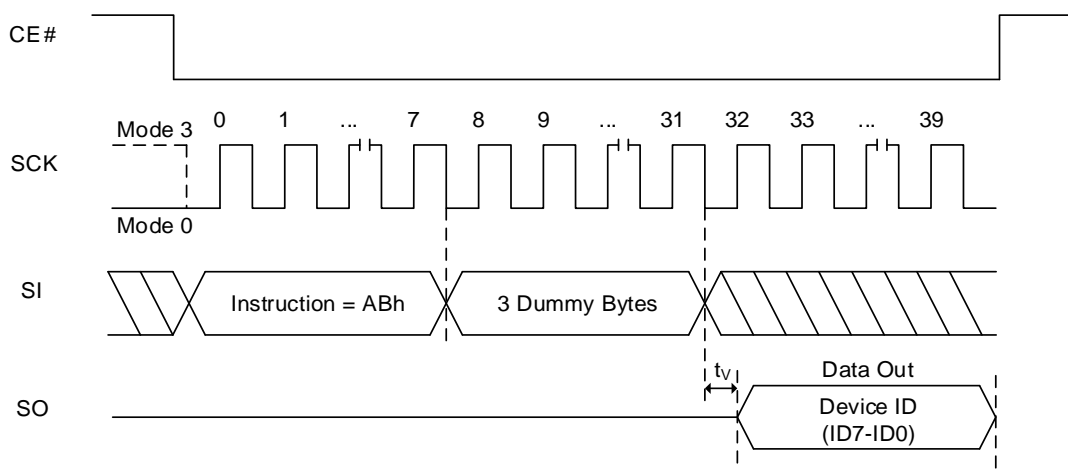
The Release from Power-down/Read Device ID instruction is a multi-purpose instruction. It can support both SPI and QPI modes. The Read Product Identification (RDID) instruction is for reading out the old style of 8-bit Electronic Signature, whose values are shown as the table of Product Identification.

The RDID instruction code is followed by three dummy bytes, each bit being latched-in on SI during the rising SCK edge. Then the Device ID is shifted out on SO with the MSB first, each bit been shifted out during the falling edge of SCK. The RDID instruction is ended by driving CE# high. The Device ID (ID7-ID0) outputs repeatedly if additional clock cycles are continuously sent to SCK while CE# is at low.

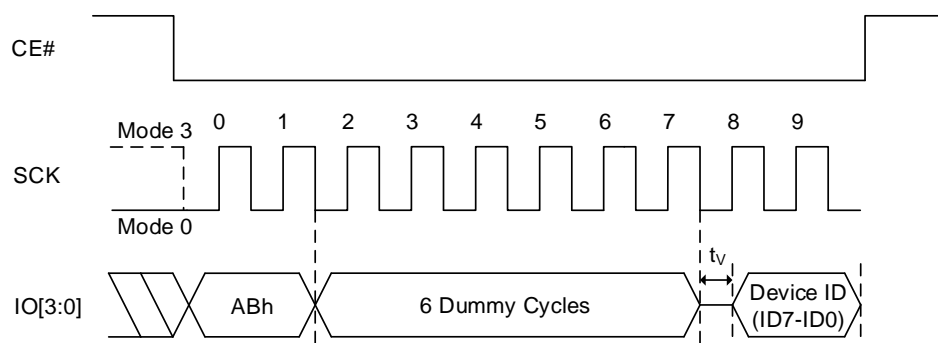
**Table 8.9 Product Identification**

Manufacturer ID		(MF7-MF0)	
ISSI Serial Flash		9Dh	
Instruction	ABh	90h	9Fh
Part Number	Device ID (ID7-ID0)		Memory Type + Capacity (ID15-ID0)
IS25RLP256E	18h		6019h
IS25RWP256E	18h		7019h

**Figure 8.69 Read Product Identification Sequence**



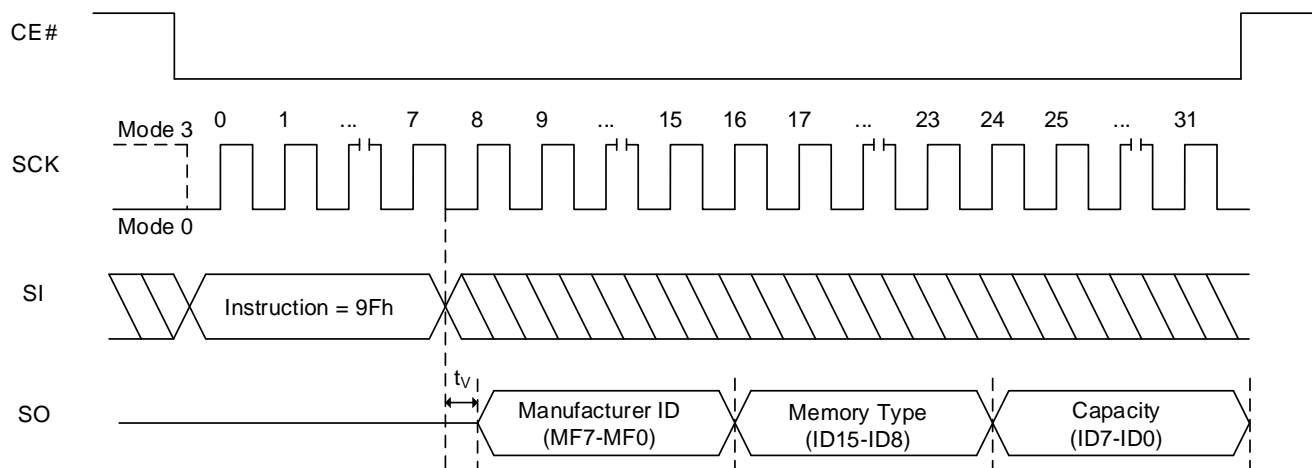
**Figure 8.70 Read Product Identification Sequence In QPI Mode**



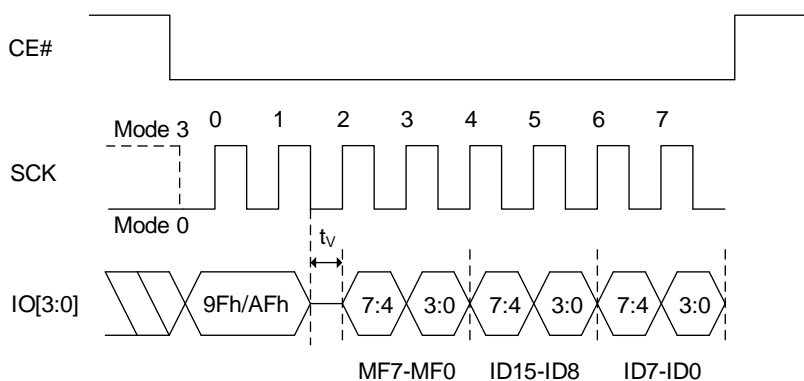
### 8.32 READ PRODUCT IDENTIFICATION BY JEDEC ID OPERATION (RDJDID, 9Fh; RDJDIDQ, AFh)

The JEDEC ID READ instruction allows the user to read the manufacturer and product ID of devices. Refer to Table 8.8 Product Identification for Manufacturer ID and Device ID. After the JEDEC ID READ command (9Fh in SPI mode, AFh in QPI mode) is input, the Manufacturer ID is shifted out MSB first followed by the 2-byte electronic ID (ID15-ID0) that indicates Memory Type and Capacity, one bit at a time. Each bit is shifted out during the falling edge of SCK. If CE# stays low after the last bit of the 2-byte electronic ID, the Manufacturer ID and 2-byte electronic ID will loop until CE# is pulled high.

**Figure 8.71 Read Product Identification by JEDEC ID Read Sequence in SPI mode**



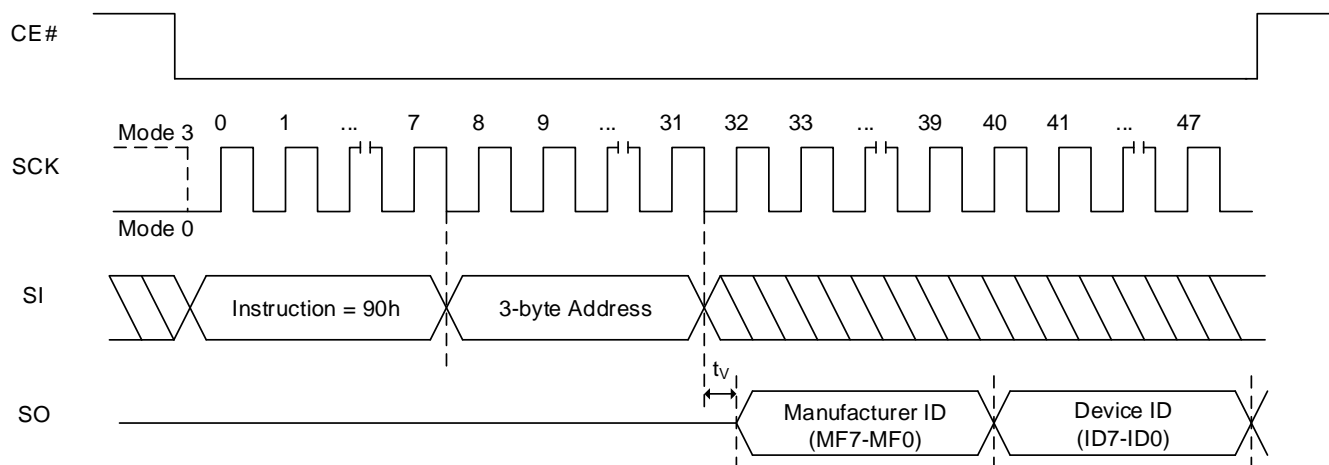
**Figure 8.72 RDJDID and RDJDIDQ (Read JEDEC ID) Sequence In QPI Mode**



### 8.33 READ DEVICE MANUFACTURER AND DEVICE ID OPERATION (RDMDID, 90h)

The Read Device Manufacturer and Device ID (RDMDID) instruction allows the user to read the Manufacturer and product ID of devices. Refer to Table 8.8 Product Identification for Manufacturer ID and Device ID. The RDMDID instruction code is followed by two dummy bytes and one byte address (A7~A0), each bit being latched-in on SI during the rising edge of SCK. If one byte address is initially set as A0 = 0, then the Manufacturer ID is shifted out on SO with the MSB first followed by the device ID (ID7-ID0). Each bit is shifted out during the falling edge of SCK. If one byte address is initially set as A0 = 1, then Device ID7-ID0 will be read first followed by the Manufacturer ID. The Manufacturer and Device ID can be read continuously alternating between the two until CE# is driven high.

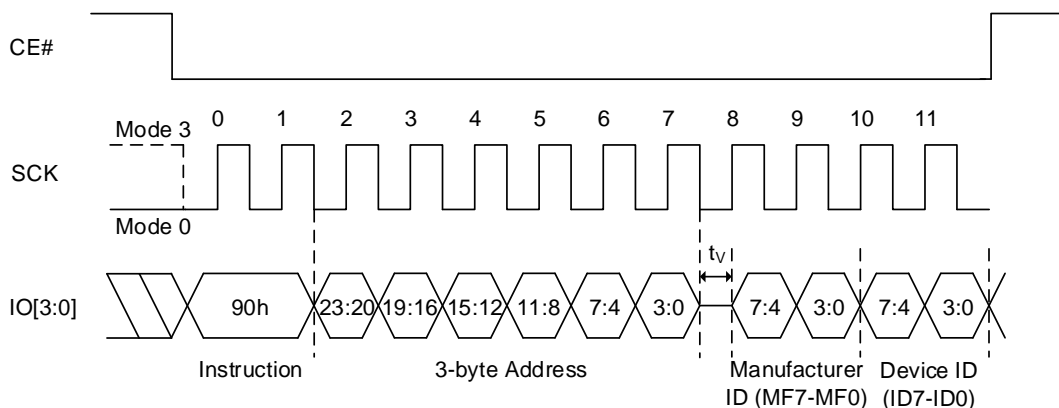
**Figure 8.73 Read Product Identification by RDMDID Sequence In SPI Mode (A0=0)**



**Notes:**

1. ADDRESS A0 = 0, will output the 1-byte Manufacturer ID (MF7-MF0) → 1-byte Device ID (ID7-ID0)  
ADDRESS A0 = 1, will output the 1-byte Device ID (ID7-ID0) → 1-byte Manufacturer ID (MF7-MF0)
2. The Manufacturer and Device ID can be read continuously and will alternate from one to the other until CE# pin is pulled high.

**Figure 8.74 Read Product Identification by RDMDID Sequence In QPI Mode (A0=0)**



**Notes:**

1. ADDRESS A0 = 0, will output the 1-byte Manufacturer ID (MF7-MF0) → 1-byte Device ID (ID7-ID0)  
ADDRESS A0 = 1, will output the 1-byte Device ID (ID7-ID0) → 1-byte Manufacturer ID (MF7-MF0)
2. The Manufacturer and Device ID can be read continuously and will alternate from one to the other until CE# pin is pulled high.

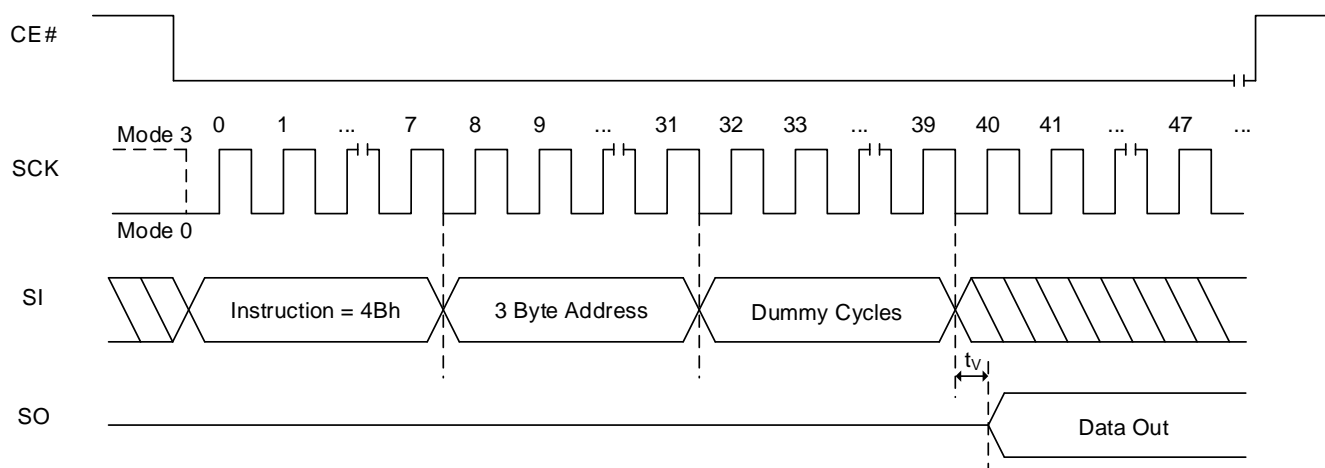
### 8.34 READ UNIQUE ID NUMBER (RDUID, 4Bh)

The Read Unique ID Number (RDUID) instruction accesses a factory-set read-only 16-byte number that is unique to the device. The ID number can be used in conjunction with user software methods to help prevent copying or cloning of a system. The RDUID instruction is instated by driving the CE# pin low and shifting the instruction code (4Bh) followed by 3 address bytes and dummy cycles (configurable, default is 8 clocks). After which, the 16-byte ID is shifted out on the falling edge of SCK as shown below.

As a result, the sequence of RDUID instruction is same as FAST READ. RDUID sequence in QPI mode is also same as FAST READ sequence in QPI mode except for the instruction code. Refer to the FAST READ operation in QPI mode.

**Note:** 16 bytes of data will repeat as long as CE# is low and SCK is toggling.

**Figure 8.75 RDUID Sequence In SPI Mode**



**Note:** Dummy cycles depends on Read Parameter setting. Detailed information in Table 6.11 Read Dummy Cycles.

**Table 8.10 Unique ID Addressing**

A[23:16]	A[15:9]	A[8:4]	A[3:0]
XXh	XXh	00h	0h Byte address
XXh	XXh	00h	1h Byte address
XXh	XXh	00h	2h Byte address
XXh	XXh	00h	⋮
XXh	XXh	00h	Fh Byte address

**Note:** XX means “don’t care”.

### 8.35 READ SFDP OPERATION (RDSFDP, 5Ah)

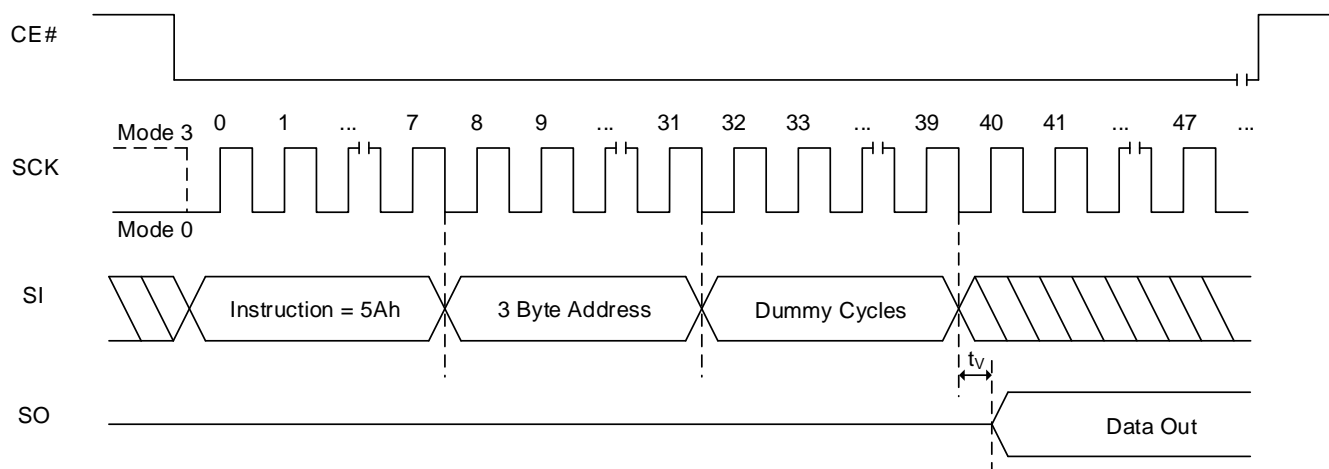
The Serial Flash Discoverable Parameters (SFDP) standard provides a consistent method of describing the functions and features of serial Flash devices in a standard set of internal parameter tables. These parameters can be interrogated by host system software to enable adjustments needed to accommodate divergent features from multiple vendors. For more details please refer to the JEDEC Standard JESD216 (Serial Flash Discoverable Parameters).

The sequence of issuing RDSFDP instruction in SPI mode is:

CE# goes low → Send RDSFDP instruction (5Ah) → Send 3 address bytes on SI pin → 8 dummy cycles on SI pin → Read SFDP code on SO → End RDSFDP operation by driving CE# high at any time during data out. Refer to ISSI's Application note for SFDP table. The data at the addresses that are not specified in SFDP table are undefined.

RDSFDP Sequence in QPI mode, has 8 dummy cycles before SFDP code, too.

**Figure 8.76 RDSFDP (Read SFDP) Sequence**



### 8.36 NO OPERATION (NOP, 00h)

The No Operation command solely cancels a Reset Enable command and has no impact on any other commands. It is available in both SPI and QPI modes. To execute a NOP, the host drives CE# low, sends the NOP command cycle (00h), then drives CE# high.

### 8.37 SOFTWARE RESET (RESET-ENABLE (RSTEN, 66h) AND RESET (RST, 99h)) AND HARDWARE RESET

The Software Reset operation is used as a system reset that puts the device in normal operating mode. During the Reset operation, the value of volatile registers will default back to the value in the corresponding non-volatile register. However, the volatile FREEZE bit and the volatile PPB Lock bit in the PPB Lock Register are not changed by Software Reset. In all other respects, Software Reset is the same as Hardware Reset. This operation consists of two commands: Reset-Enable (RSTEN) and Reset (RST). The operation requires the Reset-Enable command followed by the Reset command. Any command other than the Reset command after the Reset-Enable command will disable the Reset-Enable.

Execute the CE# pin low → sends the Reset-Enable command (66h), and drives CE# high. Next, the host drives CE# low again, sends the Reset command (99h), and pulls CE# high.

For the device with dedicated RESET#, the RESET# pin is always applicable regardless of the QE bit value.

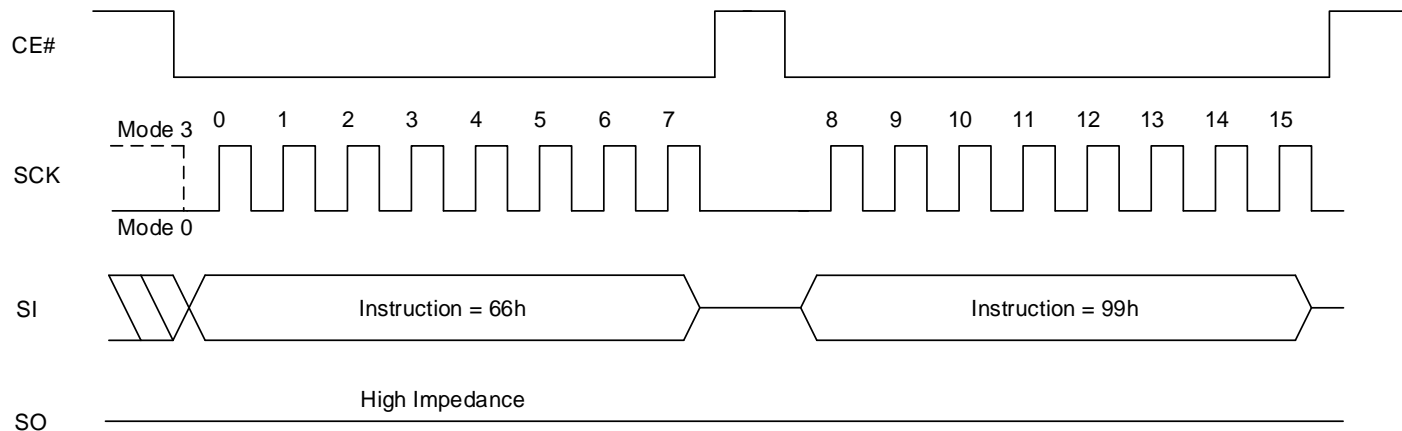
In order to activate Hardware Reset, the RESET# pin (or ball) must be driven low for a minimum period of  $t_{RESET}$  (100ns). Drive RESET# low for a minimum period of  $t_{RESET}$  will interrupt any on-going internal and external operations, release the device from deep power down mode<sup>1</sup>, disable all input signals, force the output pin enter a state of high impedance, and reset all the read parameters.

The required wait time after activating a HW Reset before the device will accept another instruction is  $t_{HWRST}$  of 35us.

The Software/Hardware Reset during an active Program or Erase operation aborts the operation, which can result in corrupting or losing the data of the targeted address range. Depending on the prior operation, the reset timing may vary. Recovery from a Write operation will require more latency than recovery from other operations.

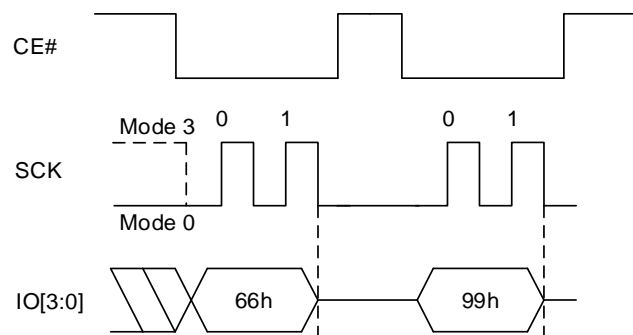
**Note: The Status and Function Registers remain unaffected.**

**Figure 8.77 Software Reset Enable and Software Reset Sequence (RSTEN, 66h + RST, 99h)**





**Figure 8.78 Software Reset Enable and Software Reset Sequence In QPI Mode (RSTEN, 66h + RST, 99h)**



### 8.38 SECURITY INFORMATION ROW

The security Information Row is comprised of an additional 4 x 256 bytes of programmable information. The security bits can be reprogrammed by the user. Any program security instruction issued while an erase, program or write cycle is in progress is rejected without having any effect on the cycle that is in progress.

**Table 8.11 Information Row Valid Address Range**

Address Assignment	A[23:16]	A[15:8]	A[7:0]
IRL0 (Information Row Lock0)	00h	00h	Byte address
IRL1	00h	10h	Byte address
IRL2	00h	20h	Byte address
IRL3	00h	30h	Byte address

Bit 7~4 of the Function Register is used to permanently lock the programmable memory array.

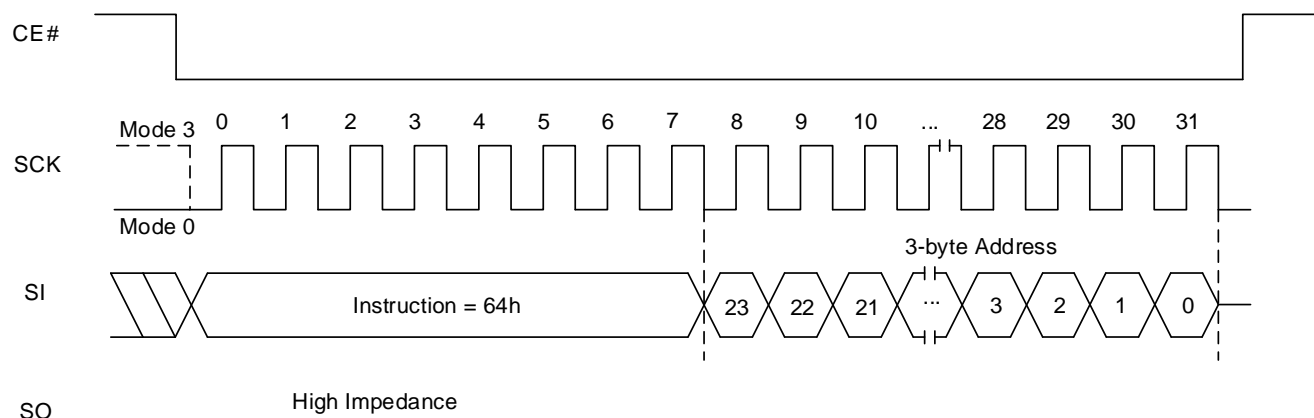
When Function Register bit IRLx = "0", the 256 bytes of the programmable memory array can be programmed.  
When Function Register bit IRLx = "1", the 256 bytes of the programmable memory array function as read only.

### 8.39 INFORMATION ROW ERASE OPERATION (IRER, 64h)

Information Row Erase (IRER) instruction erases the data in the Information Row x (x: 0~3) array. Prior to the operation, the Write Enable Latch (WEL) must be set via a Write Enable (WREN) instruction. The WEL bit is automatically reset after the completion of the operation.

The sequence of IRER operation: Pull CE# low to select the device → Send IRER instruction code → Send three address bytes → Pull CE# high. CE# should remain low during the entire instruction sequence. Once CE# is pulled high, Erase operation will begin immediately. The internal control logic automatically handles the erase voltage and timing.

**Figure 8.79 IRER (Information Row Erase) Sequence**



#### 8.40 INFORMATION ROW PROGRAM OPERATION (IRP, 62h)

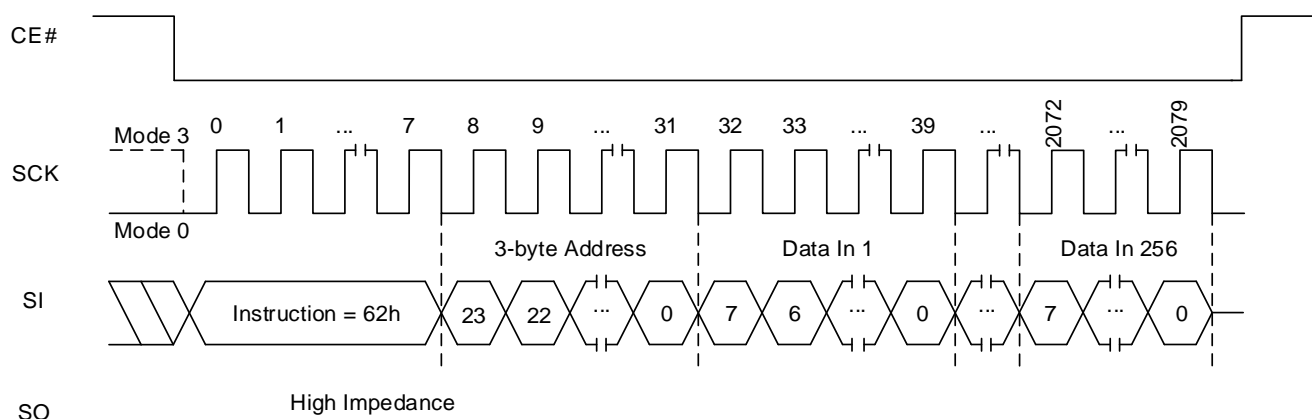
The Information Row Program (IRP) instruction allows up to 256 bytes of data to be programmed into the memory in a single operation. Before the execution of IRP instruction, the Write Enable Latch (WEL) must be enabled through a Write Enable (WREN) instruction.

The IRP instruction code, three address bytes and program data (1 to 256 bytes) should be sequentially input. Three address bytes has to be input as specified in the Table 8.10 Information Row Valid Address Range. Program operation will start once the CE# goes high, otherwise the IRP instruction will not be executed. The internal control logic automatically handles the programming voltages and timing. During a program operation, all instructions will be ignored except the RDSR instruction. The progress or completion of the program operation can be determined by reading the WIP bit. If the WIP bit is “1”, the program operation is still in progress. If WIP bit is “0”, the program operation has completed.

If more than 256 bytes data are sent to a device, the address counter rolls over within the same page. The previously latched data are discarded and the last 256 bytes data are kept to be programmed into the page. The starting byte can be anywhere within the page. When the end of the page is reached, the address will wrap around to the beginning of the same page. If the data to be programmed are less than a full page, the data of all other bytes on the same page will remain unchanged.

**Note:** A program operation can alter “1”s into “0”s, but an erase operation is required to change “0”s back to “1”s. A byte cannot be reprogrammed without first erasing the corresponding Information Row array which is one of IR0~3.

**Figure 8.80 IRP (Information Row Program) Sequence**



#### 8.41 INFORMATION ROW READ OPERATION (IRRD, 68h)

The IRRD instruction is used to read memory data.

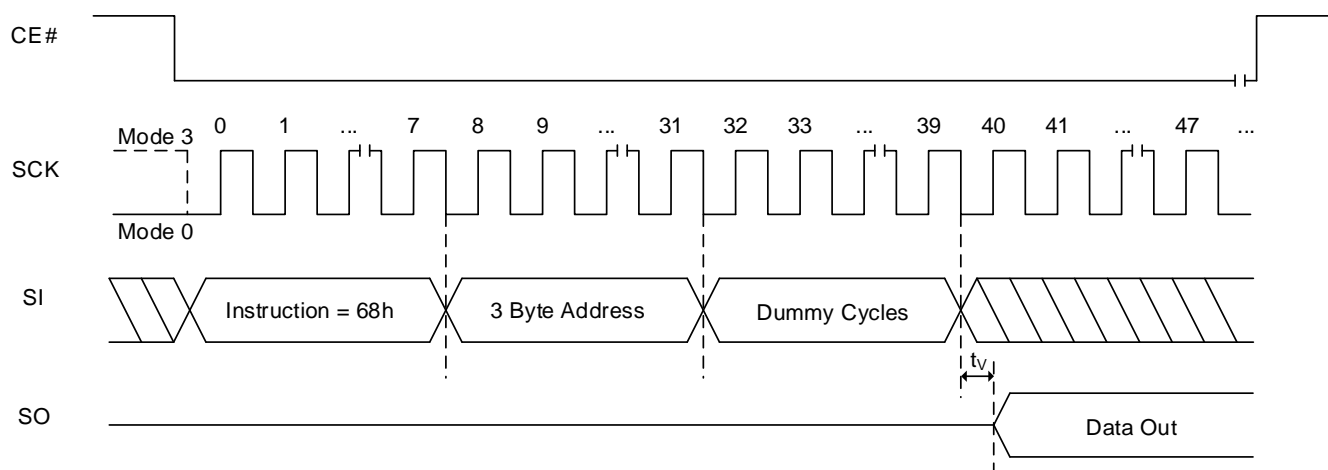
The IRRD instruction code is followed by three address bytes (A23 - A0) and dummy cycles (configurable, default is 8 clocks), transmitted via the SI line, with each bit latched-in during the rising edge of SCK. Then the first data byte addressed is shifted out on the SO line, with each bit shifted out at a maximum frequency  $f_{CT}$ , during the falling edge of SCK.

The address is automatically incremented by one after each byte of data is shifted out. Once the address reaches the last address of each 256 byte Information Row, the next address will not be valid and the data of the address will be garbage data. It is recommended to repeat four times IRRD operation that reads 256 byte with a valid starting address of each Information Row in order to read all data in the 4 x 256 byte Information Row array. The IRRD instruction is terminated by driving CE# high (VIH).

If an IRRD instruction is issued while an Erase, Program or Write cycle is in process (WIP=1) the instruction is ignored and will not have any effects on the current cycle

The sequence of IRRD instruction is same as Fast Read except for the instruction code. IRRD QPI sequence is same as Fast Read QPI except for the instruction code. Refer to the Fast Read QPI operation.

**Figure 8.81 IRRD (Information Row Read) Sequence**



**Note:** Dummy cycles depends on Read Parameter setting. Detailed information in Table 6.11 Read Dummy Cycles.

#### **8.42 FAST READ DTR MODE OPERATION (FRDTR, 0Dh or 4FRDTR, 0Eh)**

The FRDTR/4FRDTR instruction is for doubling the data in and out. Signals are triggered on both rising and falling edge of clock. The address is latched on both rising and falling edge of SCK, and data of each bit shifts out on both rising and falling edge of SCK. The 2-bit address can be latched-in at one clock, and 2-bit data can be read out at one clock, which means one bit at the rising edge of clock, the other bit at the falling edge of clock.

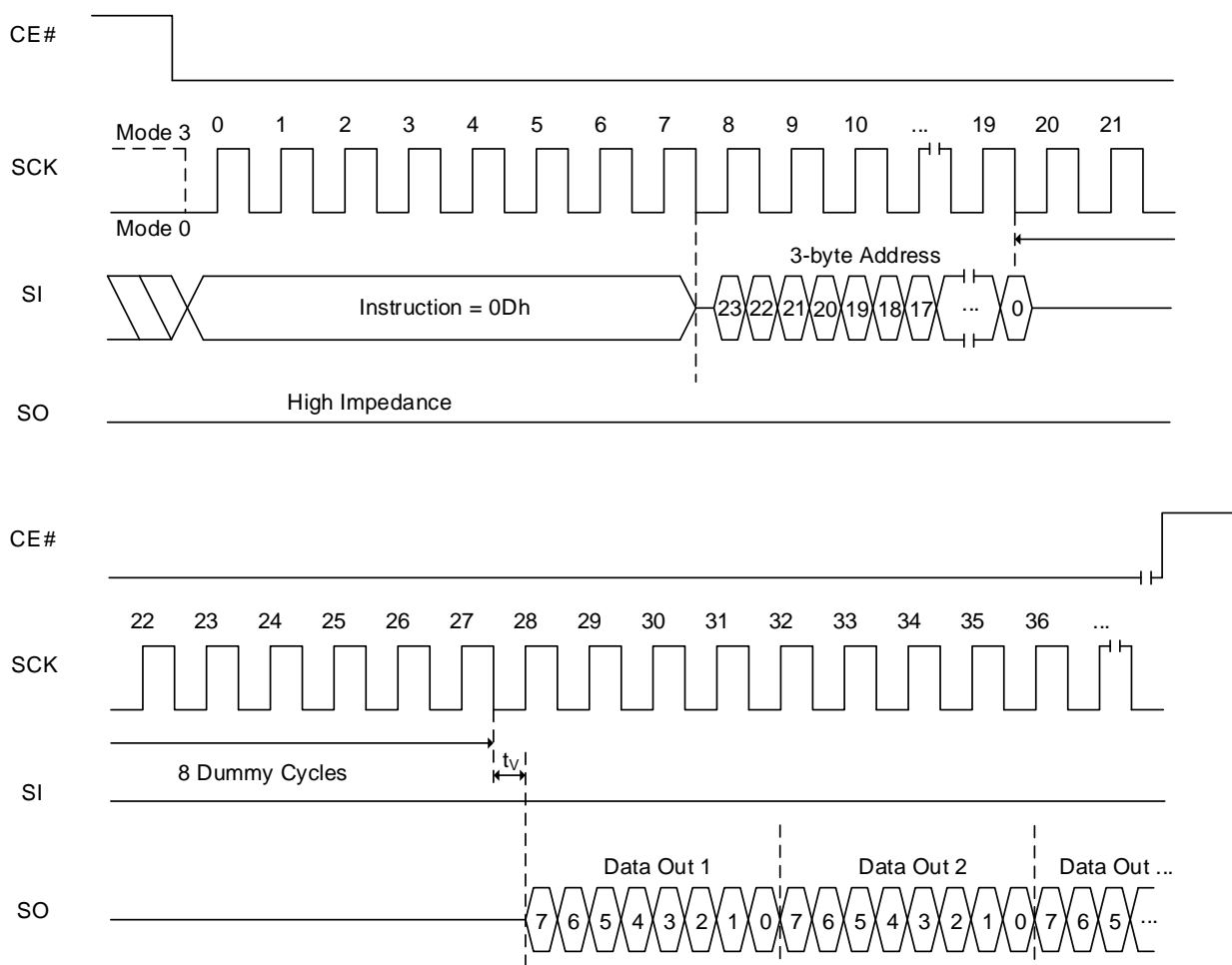
The first address byte can be at any location. The address is automatically increased to the next higher address after each byte of data is shifted out, so the whole memory can be read out in a single FRDTR/4FRDTR instruction. The address counter rolls over to 0 when the highest address is reached.

- 0Dh (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- 0Dh (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- 0Eh is followed by a 4-byte address (A31-A0)

The sequence of issuing FRDTR/4FRDTR instruction is: CE# goes low → Sending FRDTR/4FRDTR instruction code (1bit per clock) → 3-byte or 4-byte address on SI (2-bit per clock) as above → 8 dummy clocks (configurable, default is 8 clocks) on SI → Data out on SO (2-bit per clock) → End FRDTR/4FRDTR operation via driving CE# high at any time during data out.

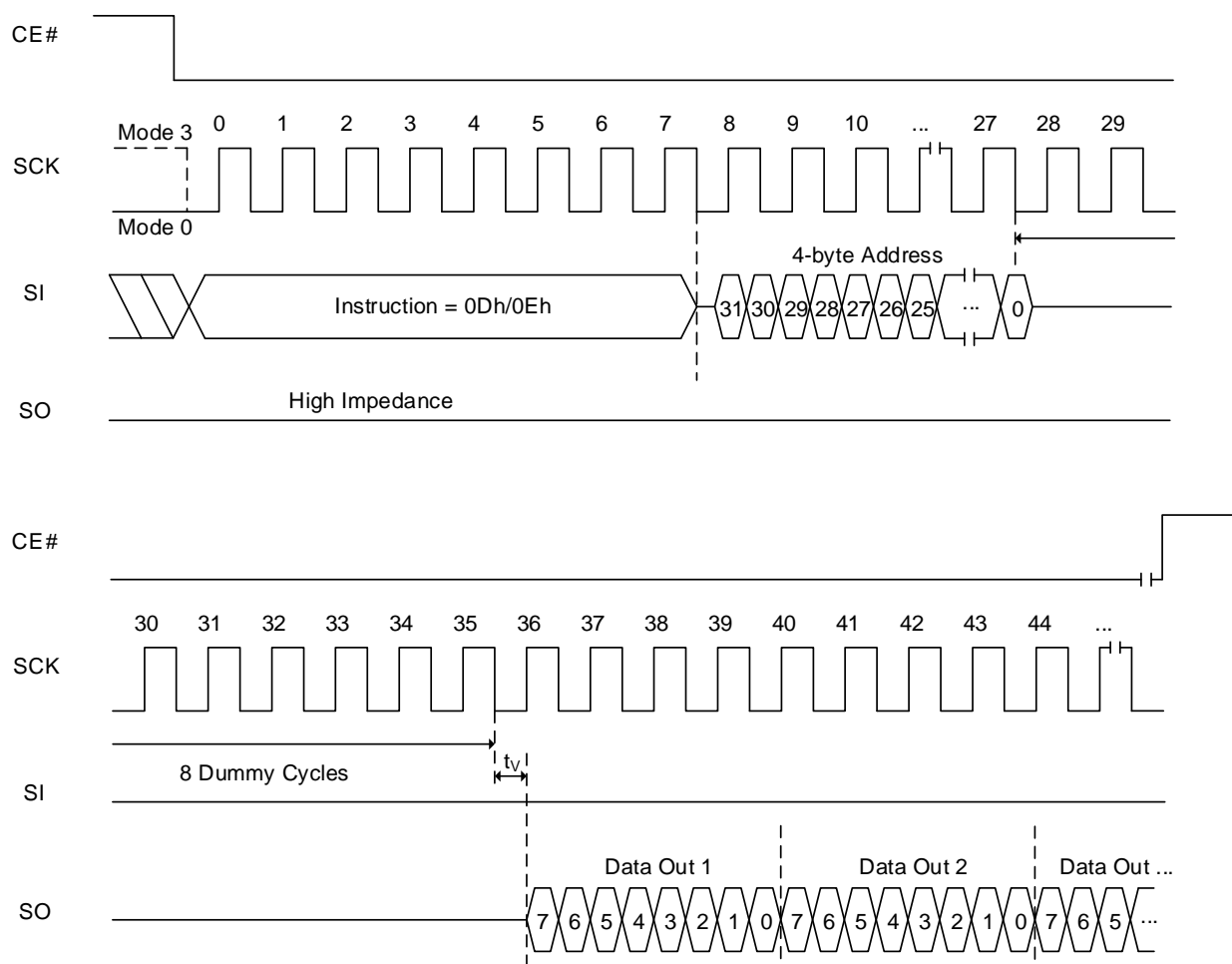
While a Program/Erase/Write Status Register cycle is in progress, FRDTR/4FRDTR instruction will be rejected without any effect on the current cycle.

Figure 8.82 FRDTR Sequence (0Dh [EXTADD=0], 3-byte address)



Note: Dummy cycles depends on Read Parameter setting. Detailed information in Table 6.11 Read Dummy Cycles.

**Figure 8.83 FRDTR Sequence (0Dh [EXTADD=1] or 0Eh, 4-byte address)**



**Note:** Dummy cycles depends on Read Parameter setting. Detailed information in Table 6.11 Read Dummy Cycles.

### FAST READ DTR OPERATION IN QPI MODE (FRDTR, 0Dh or 4FRDTR, 0Eh)

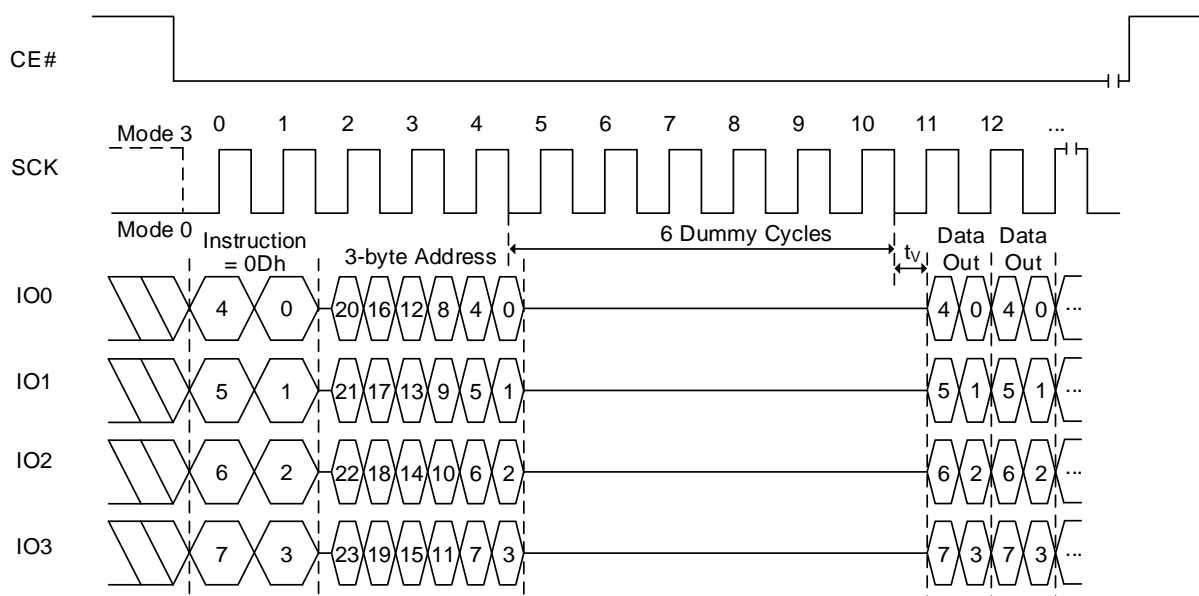
The FRDTR/4FRDTR instruction in QPI mode utilizes all four IO lines to input the instruction code so that only two clocks are required, while the FRDTR/4FRDTR instruction in SPI mode requires that the byte-long instruction code is shifted into the device only via IO0 (SI) line in eight clocks. In addition, subsequent address and data out are shifted in/out via all four IO lines unlike the FRDTR/4FRDTR instruction. Eventually this operation is same as the FRQDTR/4FRQDTR in QPI mode, but the only different thing is that AX mode is not available in the FRDTR/4FRDTR operation in QPI mode.

- 0Dh (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- 0Dh (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- 0Eh is followed by a 4-byte address (A31-A0)

The sequence of issuing FRDTR/4FRDTR instruction in QPI mode is: CE# goes low → Sending FRDTR/4FRDTR QPI instruction (4-bit per clock) → 24-bit or 32-bit address interleave on IO3, IO2, IO1 & IO0 (8-bit per clock) as above → 6 dummy clocks (configurable, default is 6 clocks) → Data out interleave on IO3, IO2, IO1 & IO0 (8-bit per clock) → End FRDTR/4FRDTR operation in QPI mode by driving CE# high at any time during data out.

If the FRDTR/4FRDTR instruction in QPI mode is issued while an Erase, Program or Write cycle is in process is in progress (WIP=1), the instruction will be rejected without any effect on the current cycle.

**Figure 8.84 FRDTR Sequence In QPI Mode (0Dh [EXTADD=0], 3-byte address)**



**Notes:**

1. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
2. Sufficient dummy cycles are required to avoid I/O contention.



Timing diagram for Mode 3 read operation. The diagram shows the relationship between CE#, SCK, and data bus IO[3:0]. CE# is active low. SCK is a clock signal. The data bus shows the instruction (0Dh/0Eh), a 4-byte address, 6 dummy cycles, and the data output. The data output is shown for IO0, IO1, IO2, and IO3. The data output is valid for  $t_{DI}$  after the last dummy cycle.

1. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
2. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

### 8.43 FAST READ DUAL IO DTR MODE OPERATION (FRDDTR, BDh or 4FRDDTR, BEh)

The FRDDTR/4FRDDTR instruction enables Double Transfer Rate throughput on dual I/O of the device in read mode. The address (interleave on dual I/O pins) is latched on both rising and falling edge of SCK, and the data (interleave on dual I/O pins) shift out on both rising and falling edge of SCK. The 4-bit address can be latched-in at one clock, and 4-bit data can be read out at one clock, which means two bits at the rising edge of clock, the other two bits at the falling edge of clock.

The first address byte can be at any location. The address is automatically increased to the next higher address after each byte of data is shifted out, so the whole memory can be read out with a single FRDDTR/4FRDDTR instruction. The address counter rolls over to 0 when the highest address is reached. Once writing FRDDTR/4FRDDTR instruction, the following address/dummy/data out will perform as 4-bit instead of previous 1-bit.

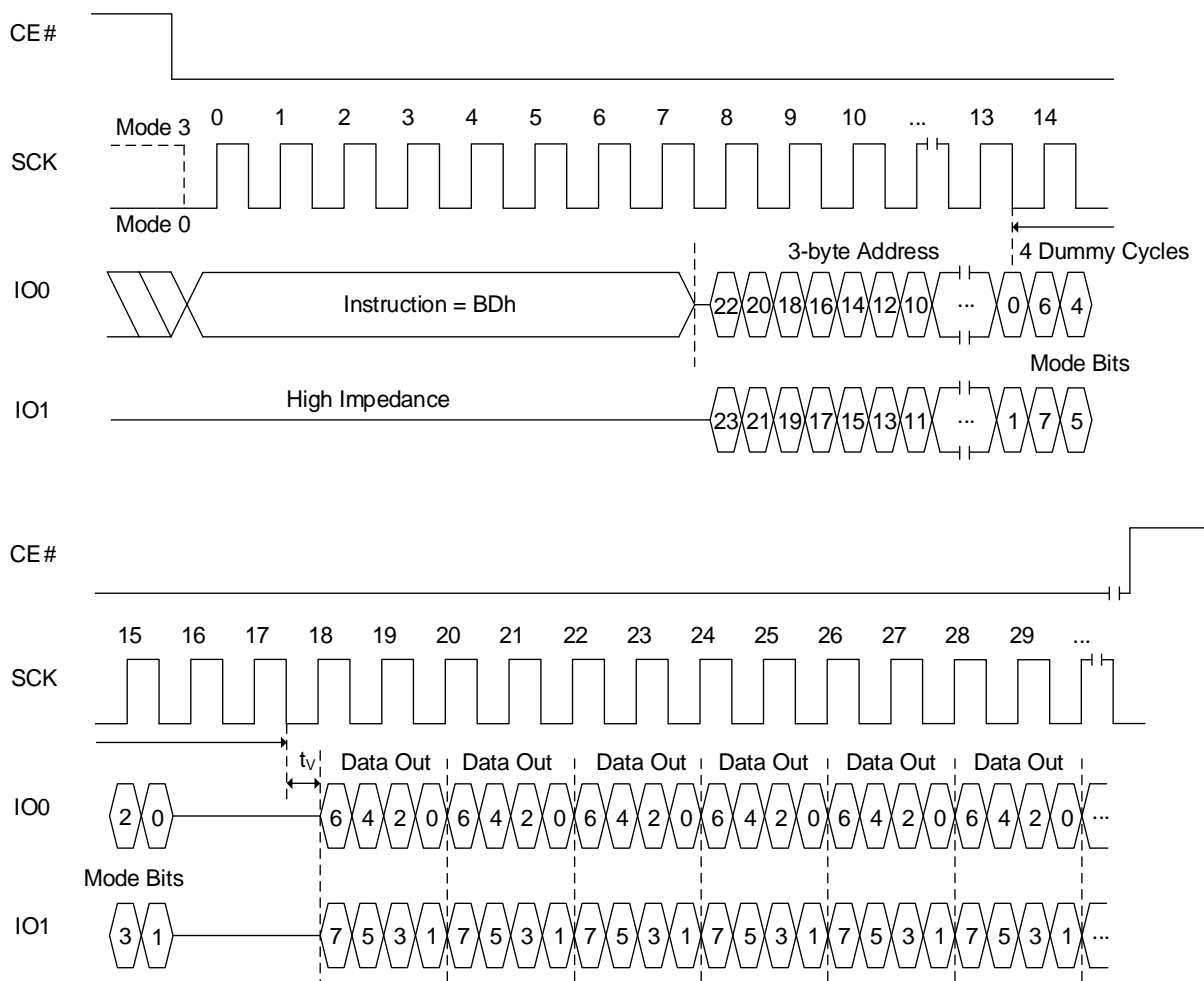
- BDh (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- BDh (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- BEh is followed by a 4-byte address (A31-A0)

The sequence of issuing FRDDTR/4FRDDTR instruction is: CE# goes low → Sending FRDDTR/4FRDDTR instruction (1-bit per clock) → 24-bit or 32-bit address interleave on IO1 & IO0 (4-bit per clock) as above → 4 dummy clocks (configurable, default is 4 clocks) on IO1 & IO0 → Data out interleave on IO1 & IO0 (4-bit per clock) → End FRDDTR/4FRDDTR operation via pulling CE# high at any time during data out (Please refer to Figures 8.86 and 8.87 for 2 x I/O Double Transfer Rate Read Mode Timing Waveform).

If AXh (where X is don't care) is input for the mode bits during dummy cycles, the device will enter AX read operation mode which enables subsequent FRDDTR/4FRDDTR execution skips command code. It saves cycles as described in Figures 8.88 and 8.89. When the code is different from AXh (where X is don't care), the device exits the AX read operation. After finishing the read operation, device becomes ready to receive a new command.

If the FRDDTR/4FRDDTR instruction is issued while an Erase, Program or Write cycle is in process is in progress (WIP=1), the instruction will be rejected without any effect on the current cycle.

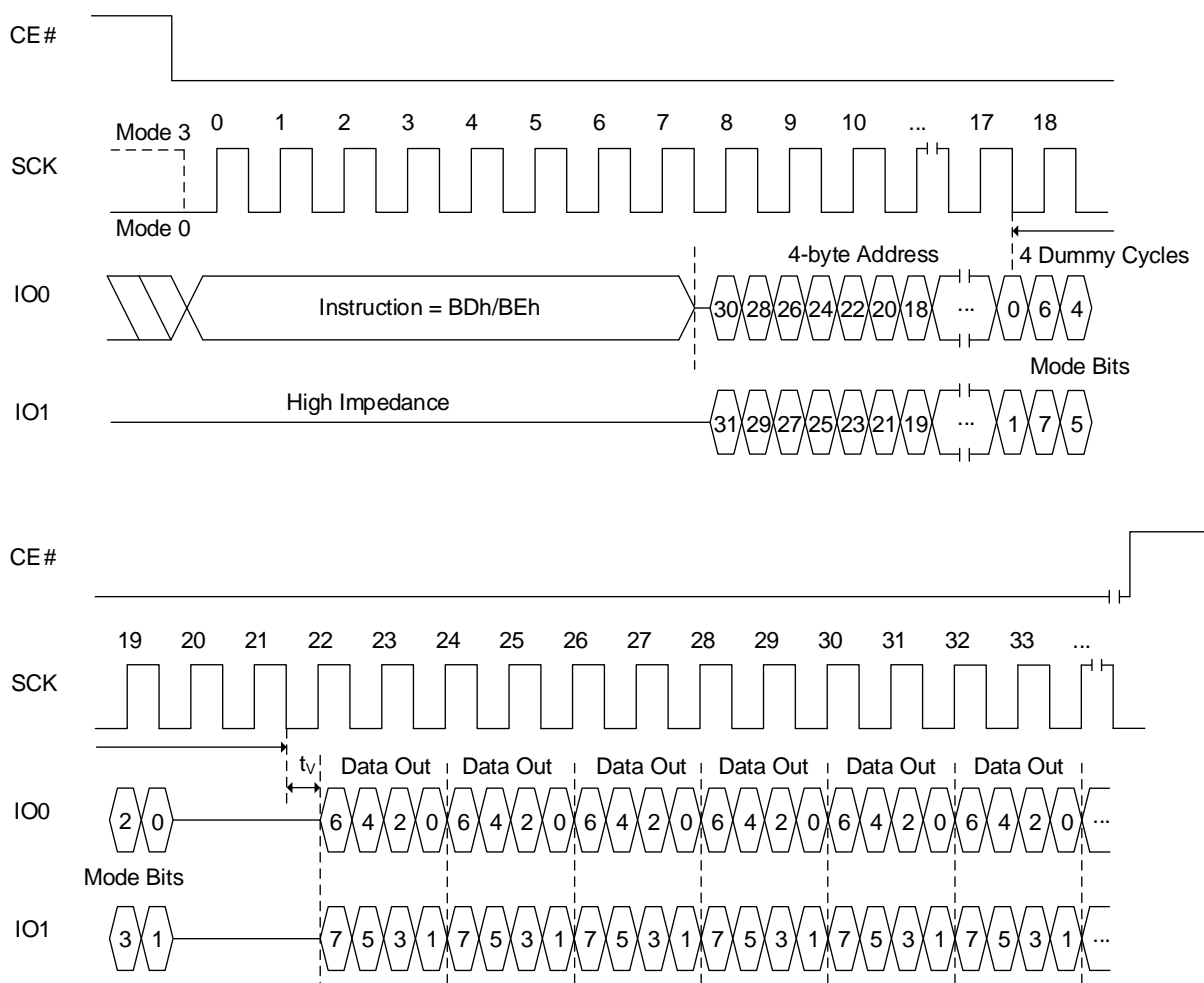
Figure 8.86 FRDDTR Sequence (BDh [EXTADD=0], 3-byte address)



**Notes:**

1. If the mode bits=AXh (where X is don't care), it can execute the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

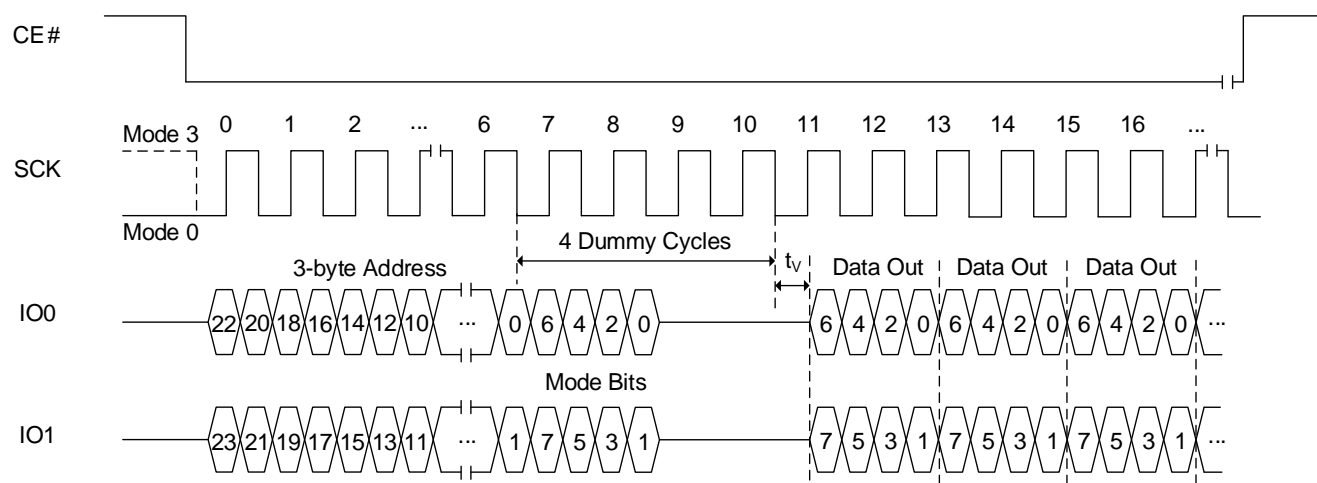
Figure 8.87 FRDDTR Sequence (BDh [EXTADD=1] or BEh, 4-byte address)



**Notes:**

1. If the mode bits=AXh (where X is don't care), it can execute the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

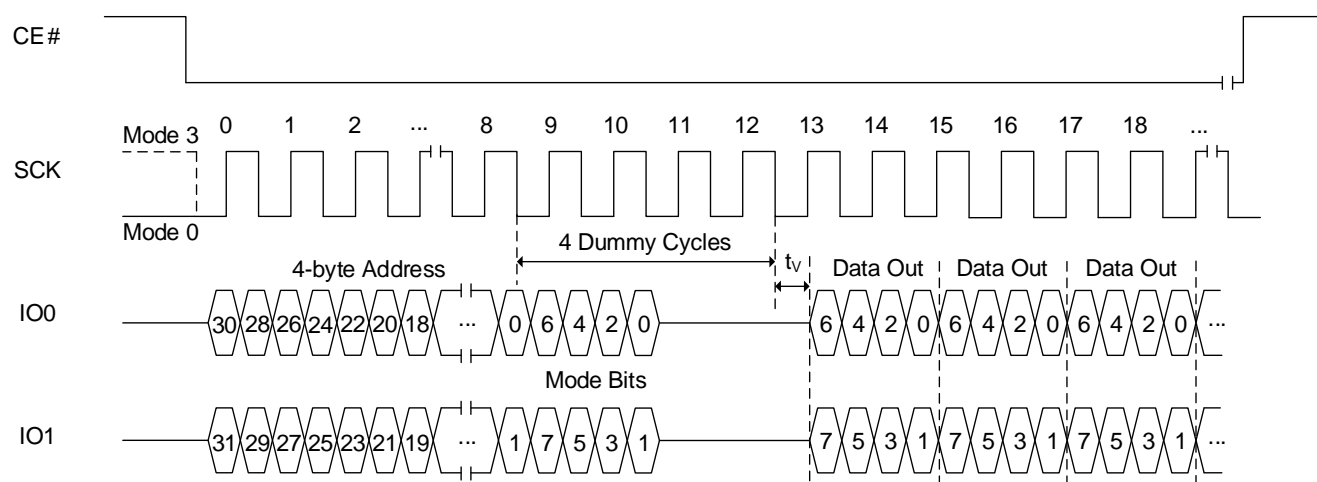
Figure 8.88 FRDDTR AX Read Sequence (BDh [EXTADD=0], 3-byte address)



**Notes:**

1. If the mode bits=AXh (where X is don't care), it will keep executing the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

Figure 8.89 FRDDTR AX Read Sequence (BDh [EXTADD=1] or BEh, 4-byte address)



**Notes:**

1. If the mode bits=AXh (where X is don't care), it will keep executing the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

#### 8.44 FAST READ QUAD IO DTR MODE OPERATION IN SPI MODE (FRQDTR, EDh or 4FRQDTR, EEh)

The FRQDTR/4FRQDTR instruction enables Double Transfer Rate throughput on quad I/O of the device in read mode.

A Quad Enable (QE) bit of Status Register must be set to "1" before sending the Fast Read Quad I/O DTR instruction.

The address (interleave on 4 I/O pins) is latched on both rising and falling edge of SCK, and data (interleave on 4 I/O pins) shift out on both rising and falling edge of SCK. The 8-bit address can be latched-in at one clock, and 8-bit data can be read out at one clock, which means four bits at the rising edge of clock, the other four bits at the falling edge of clock.

The first address byte can be at any location. The address is automatically increased to the next higher address after each byte data is shifted out, so the whole memory can be read out with a single FRQDTR/4FRQDTR instruction. The address counter rolls over to 0 when the highest address is reached. Once writing FRQDTR/4FRQDTR instruction, the following address/dummy/data out will perform as 8-bit instead of previous 1-bit.

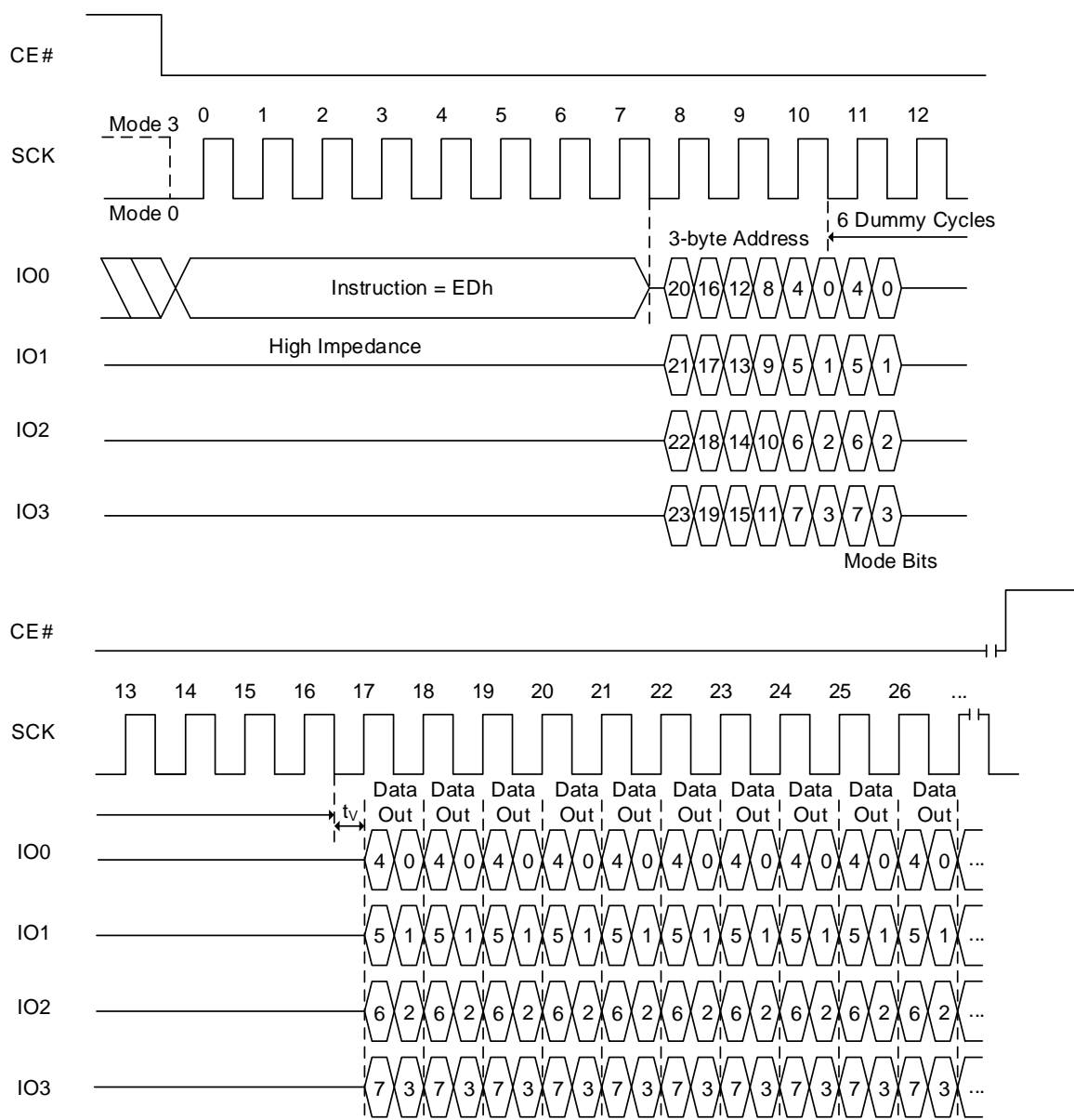
- EDh (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- EDh (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- EEh is followed by a 4-byte address (A31-A0)

The sequence of issuing FRQDTR/4FRQDTR instruction is: CE# goes low → Sending FRQDTR/4FRQDTR instruction (1-bit per clock) → 24-bit or 32-bit address interleave on IO3, IO2, IO1 & IO0 (8-bit per clock) as above → 6 dummy clocks (configurable, default is 6 clocks) → Data out interleave on IO3, IO2, IO1 & IO0 (8-bit per clock) → End FRQDTR/4FRQDTR operation by driving CE# high at any time during data out.

If AXh (where X is don't care) is input for the mode bits during dummy cycles, the device will enter AX read operation mode which enables subsequent FRQDTR/4FRQDTR execution skips command code. It saves cycles as described in Figures 8.92 and 8.93. When the code is different from AXh (where X is don't care), the device exits the AX read operation. After finishing the read operation, device becomes ready to receive a new command.

If the FRQDTR/4FRQDTR instruction is issued while an Erase, Program or Write cycle is in process (WIP=1), the instruction will be rejected without any effect on the current cycle.

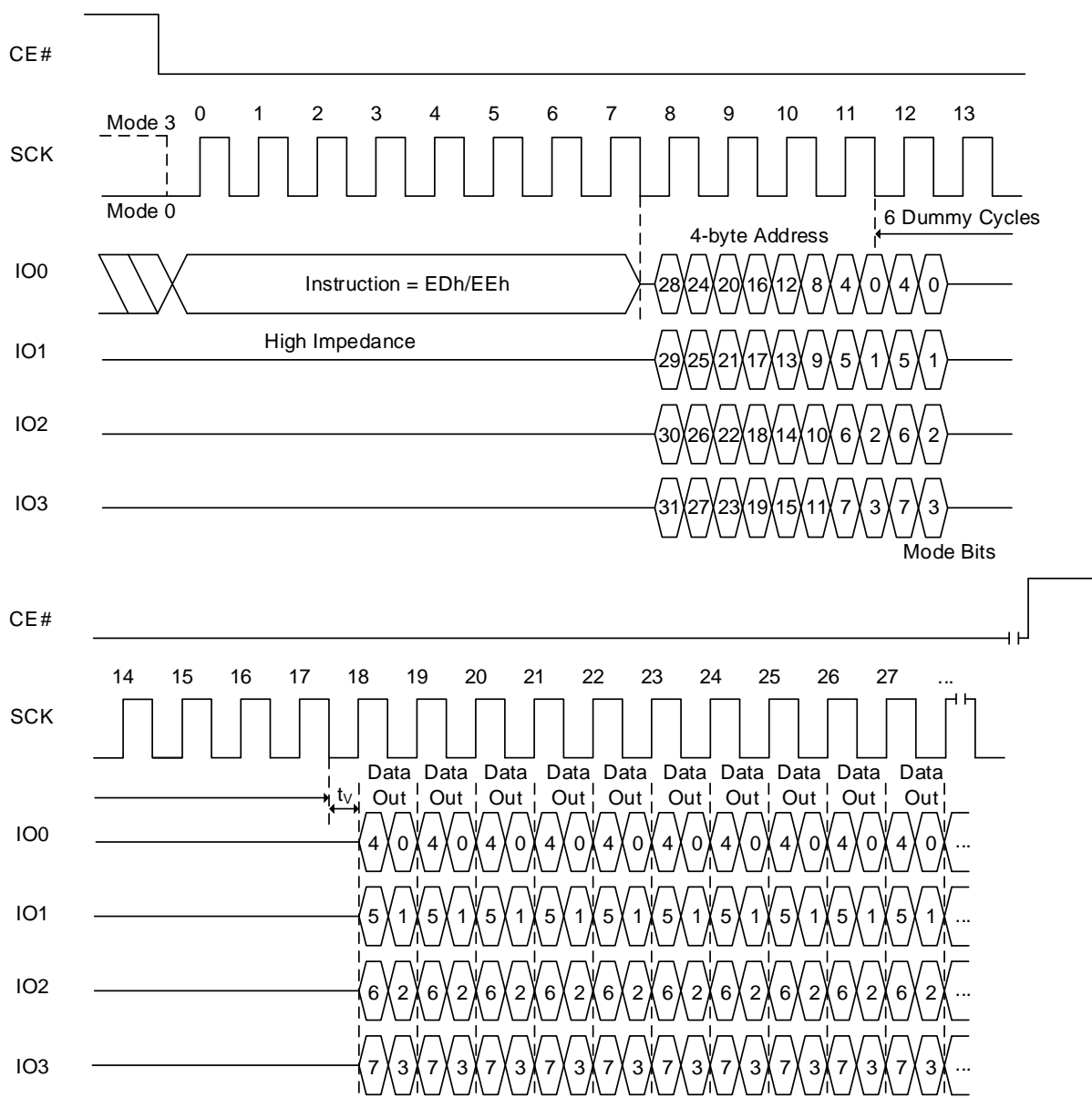
**Figure 8.90 FRQDTR Sequence In SPI Mode (EDh [EXTADD=0], 3-byte address)**



**Notes:**

1. If the mode bits=AXh (where X is don't care), it can execute the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

**Figure 8.91 FRQDTR Sequence In SPI Mode (EDh [EXTADD=1] or EEh, 4-byte address)**

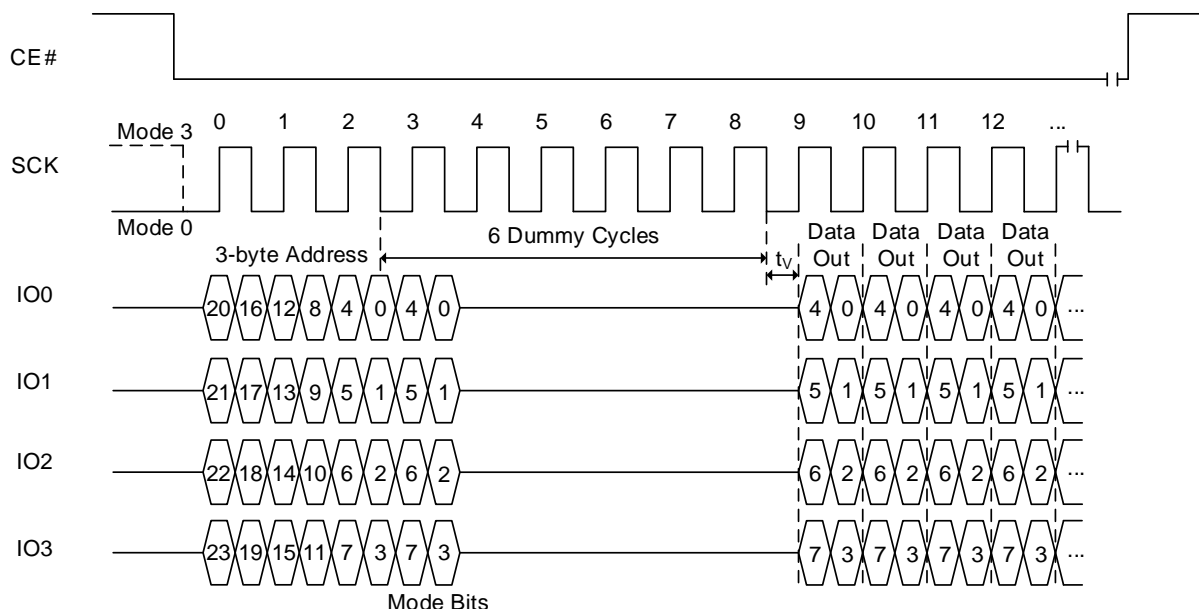


**Notes:**

1. If the mode bits=AXh (where X is don't care), it can execute the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.



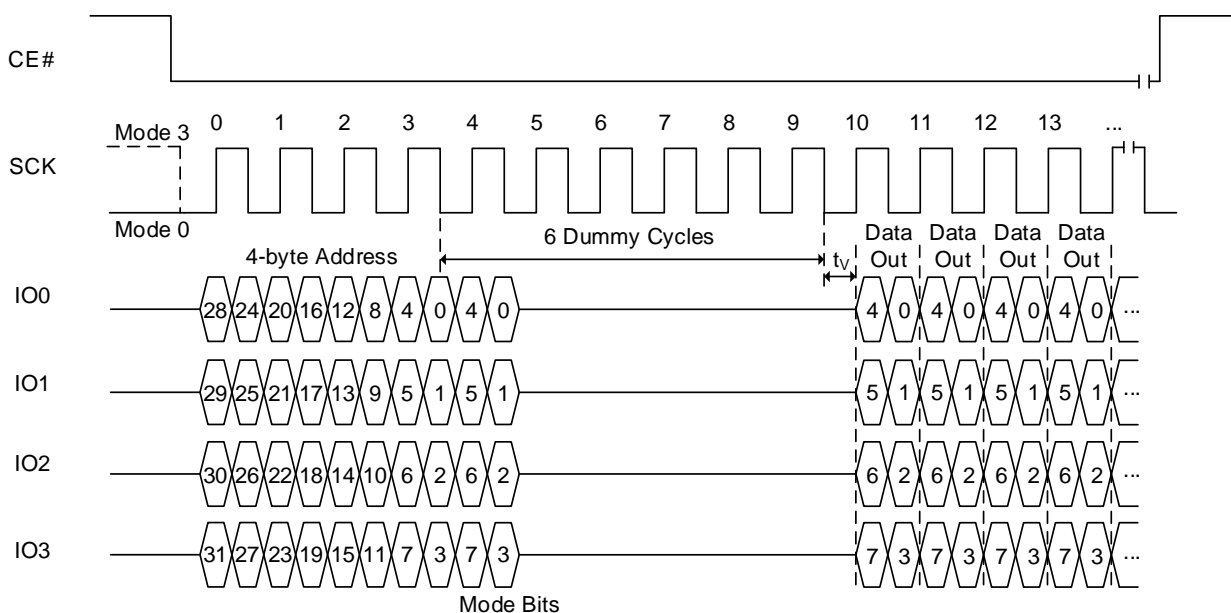
Figure 8.92 FRQDTR AX Read Sequence (EDh [EXTADD=0], 3-byte address)



**Notes:**

1. If the mode bits=AXh (where X is don't care), it will keep executing the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

Figure 8.93 FRQDTR AX Read Sequence (EDh [EXTADD=1] or EEh, 4-byte address)

**Notes:**

1. If the mode bits=AXh (where X is don't care), it will keep executing the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

**FAST READ QUAD IO DTR MODE OPERATION IN QPI MODE (FRQDTR, EDh OR 4FRQDTR, EEh)**

The FRQDTR/4FRQDTR instruction in QPI mode utilizes all four IO lines to input the instruction code so that only two clocks are required, while the FRQDTR/4FRQDTR instruction in SPI mode requires that the byte-long instruction code is shifted into the device only via IO0 line in eight clocks. As a result, 6 command cycles will be reduced by the FRQDTR/4FRQDTR instruction in QPI mode. In addition, subsequent address and data out are shifted in/out via all four IO lines like the FRQDTR/4FRQDTR instruction. In fact, except for the command cycle, the FRQDTR/4FRQDTR operation in QPI mode is exactly same as the FRQDTR/4FRQDTR operation in SPI mode.

It is not required to set QE bit to “1”.before Fast Read Quad I/O DTR instruction in QPI mode.

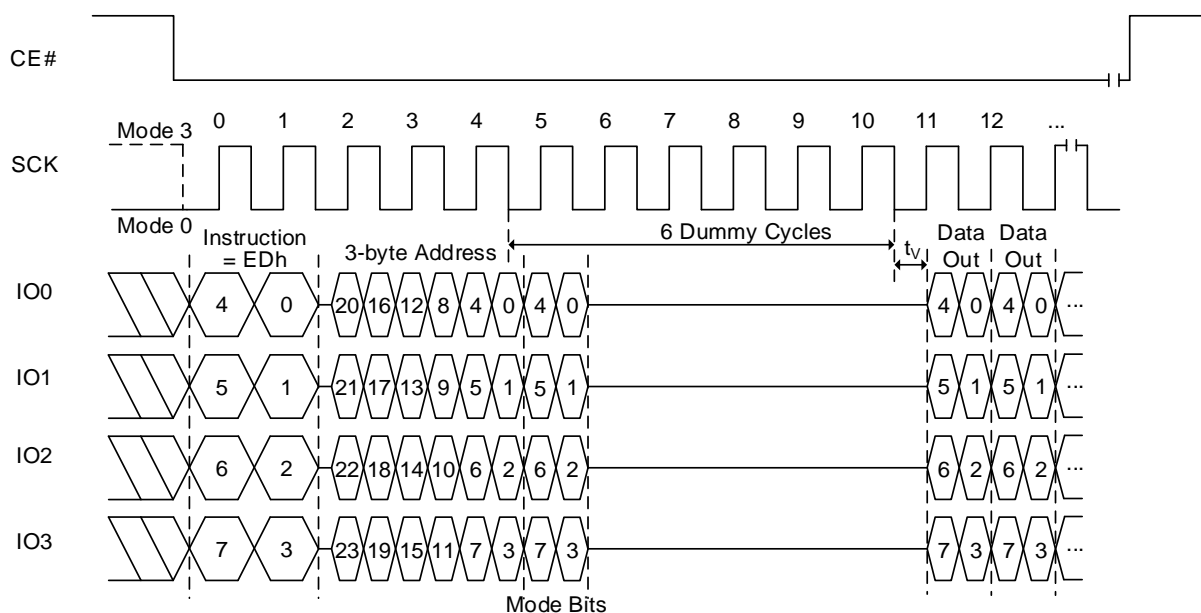
- EDh (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- EDh (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- EEh is followed by a 4-byte address (A31-A0)

The sequence of issuing FRQDTR/4FRQDTR instruction is: CE# goes low → Sending FRQDTR/4FRQDTR instruction (4-bit per clock) → 24-bit or 32-bit address interleave on IO3, IO2, IO1 & IO0 (8-bit per clock) as above → 6 dummy clocks (configurable, default is 6 clocks) → Data out interleave on IO3, IO2, IO1 & IO0 (8-bit per clock) → End FRQDTR/4FRQDTR operation by driving CE# high at any time during data out.

If AXh (where X is don't care) is input for the mode bits during dummy cycles, the device will enter AX read operation mode which enables subsequent FRQDTR/4FRQDTR in QPI mode execution skips command code. It saves cycles as described in Figures 8.92 and 8.93. When the code is different from AXh (where X is don't care), the device exits the AX read operation. After finishing the read operation, device becomes ready to receive a new command.

If the FRQDTR/4FRQDTR instruction in QPI mode is issued while an Erase, Program or Write cycle is in process is in progress (WIP=1), the instruction will be rejected without any effect on the current cycle.

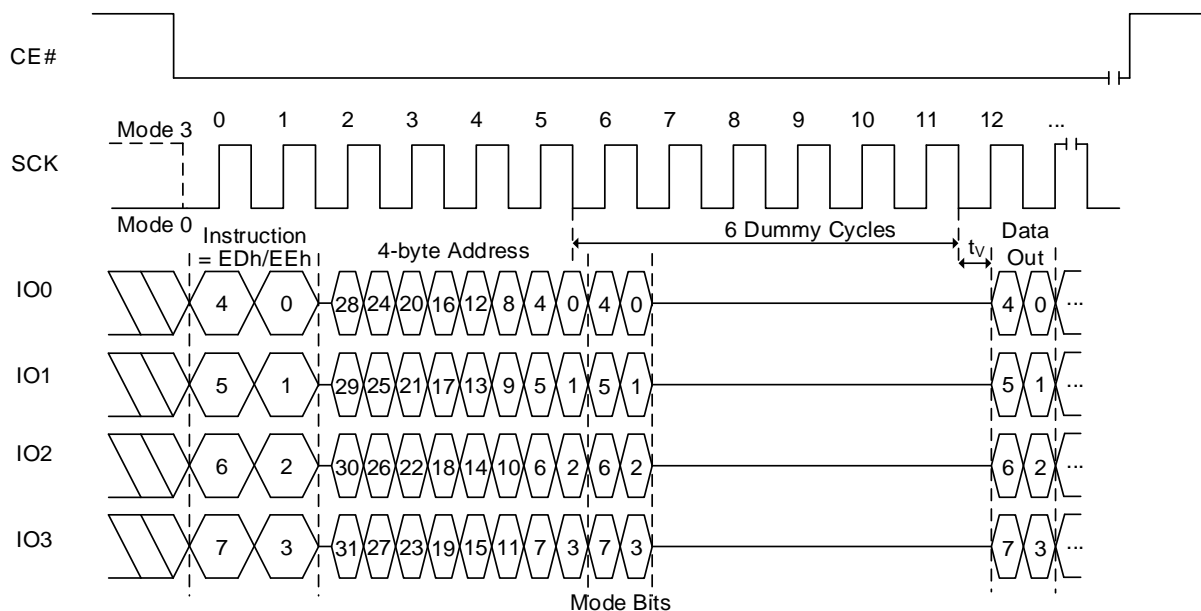
**Figure 8.94 FRQDTR Sequence In QPI Mode (EDh [EXTADD=0], 3-byte address)**



**Notes:**

1. If the mode bits=AXh (where X is don't care), it can execute the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

**Figure 8.95 FRQDTR Sequence In QPI Mode (EDh [EXTADD=1] or EEh, 4-byte address)**



**Notes:**

1. If the mode bits=AXh (where X is don't care), it can execute the AX read mode (without command). When the mode bits are different from AXh, the device exits the AX read operation.
2. Number of dummy cycles depends on clock speed. Detailed information in Table 6.11 Read Dummy Cycles.
3. Sufficient dummy cycles are required to avoid I/O contention. If the number of dummy cycles and AX bits cycles are same, then X should be Hi-Z.

## 8.45 SECTOR LOCK/UNLOCK FUNCTIONS

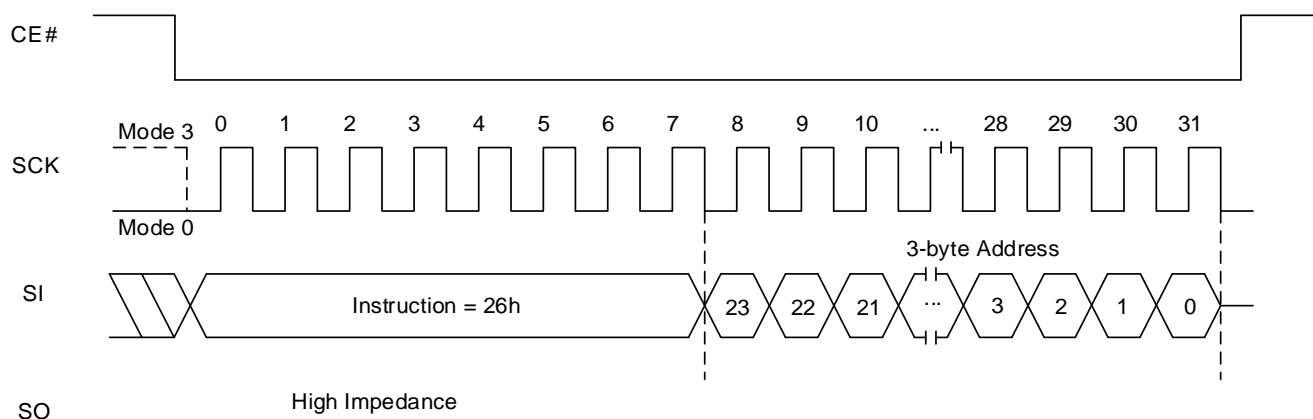
### SECTOR UNLOCK OPERATION (SECUNLOCK, 26h or 4SECUNLOCK, 25h)

The Sector Unlock command allows the user to select a specific sector to allow program and erase operations. This instruction is effective when the blocks are designated as write-protected through the BP0-BP3 bits in the Status Register and TBS bit in the Function Register. Only one sector can be enabled at any time. To enable a different sector, a previously enabled sector must be disabled by executing a Sector Lock command.

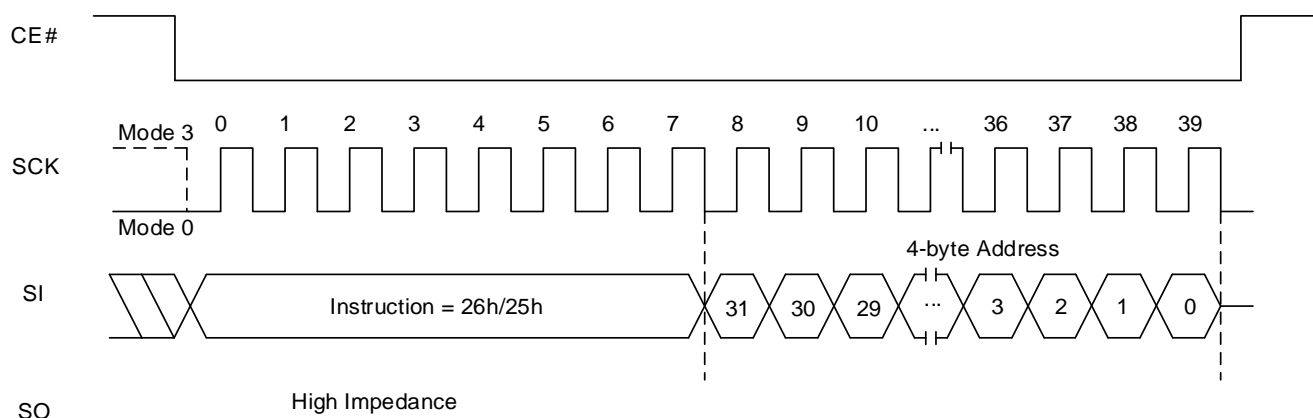
- 26h (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- 26h (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- 25h is followed by a 4-byte address (A31-A0)

The instruction code is followed by a 24-bit or 32-bit address specifying the target sector as above, but A0 through A11 are not decoded. The remaining sectors within the same block remain as read-only.

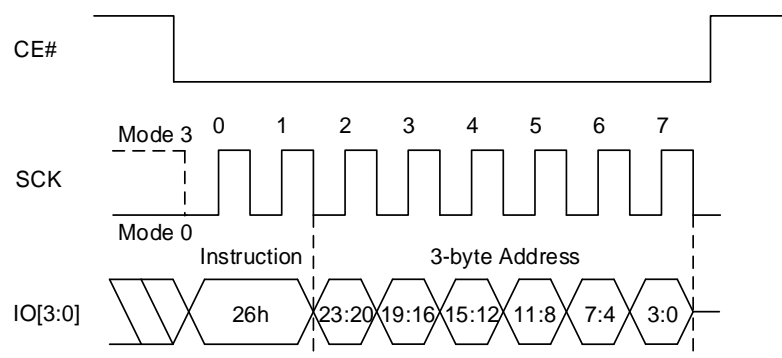
**Figure 8.96 Sector Unlock Sequence In SPI Mode (26h [EXTADD=0], 3-byte address)**



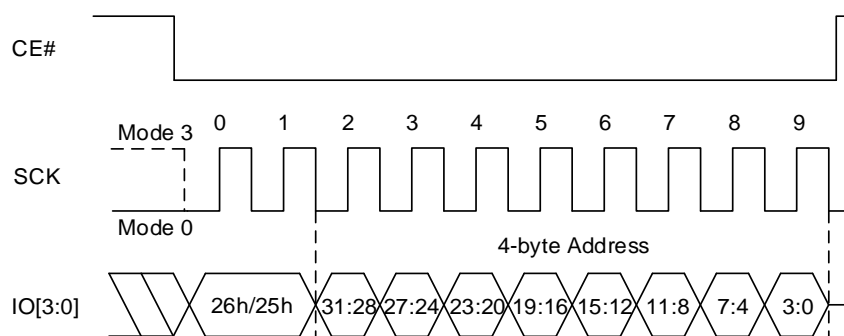
**Figure 8.97 Sector Unlock Sequence In SPI Mode (26h [EXTADD=1] or 25h, 4-byte address)**



**Figure 8.98 Sector Unlock Sequence In QPI Mode (26h [EXTADD=0], 3-byte address)**



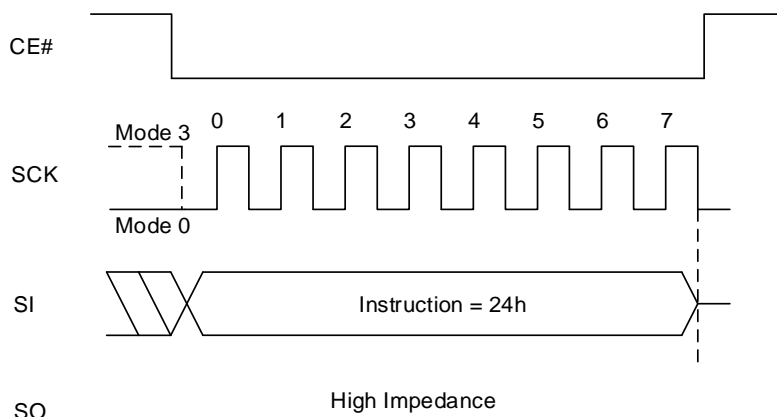
**Figure 8.99 Sector Unlock Sequence In QPI Mode (26h [EXTADD=1] or 25h, 4-byte address)**



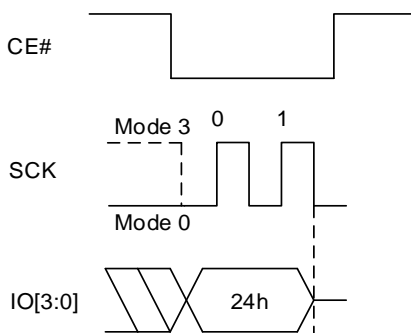
### SECTOR LOCK OPERATION (SELOCK, 24h)

The Sector Lock command relocks a sector that was previously unlocked by the Sector Unlock command. The instruction code does not require an address to be specified, as only one sector can be enabled at a time. The remaining sectors within the same block remain in read-only mode.

**Figure 8.100 Sector Lock Sequence In SPI mode**



**Figure 8.101 Sector Lock Sequence In QPI mode**





## 8.46 AUTOBOOT

SPI devices normally require 32 or more cycles of command and address shifting to initiate a read command. And, in order to read boot code from an SPI device, the host memory controller or processor must supply the read command from a hardwired state machine or from some host processor internal ROM code.

Parallel NOR devices need only an initial address, supplied in parallel in a single cycle, and initial access time to start reading boot code.

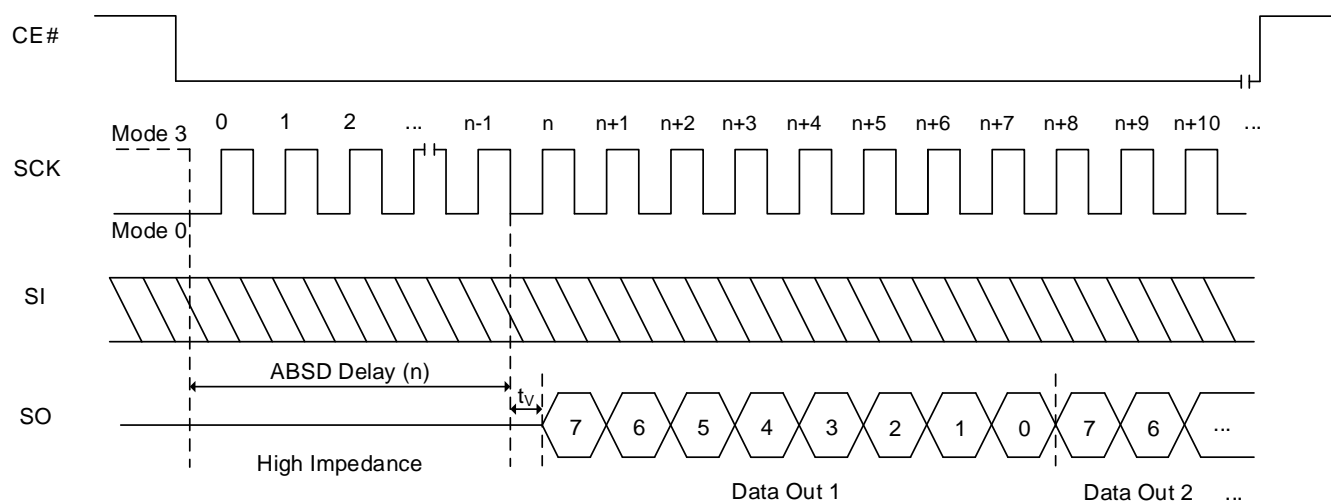
The AutoBoot feature allows the host memory controller to take boot code from the device immediately after the end of reset, without having to send a read command. This saves 32 or more cycles and simplifies the logic needed to initiate the reading of boot code.

- As part of the Power-up Reset, Hardware Reset, or Software Reset process the AutoBoot feature automatically starts a read access from a pre-specified address. At the time the reset process is completed, the device is ready to deliver code from the starting address. The host memory controller only needs to drive CE# signal from high to low and begin toggling the SCK signal. The device will delay code output for a pre-specified number of clock cycles before code streams out.
  - The Auto Boot Start Delay (ABSD) field of the AutoBoot register specifies the initial delay if any is needed by the host.
  - The host cannot send commands during this time.
  - If QE bit (Bit 6) in the Status Register is set to “1”, Fast Read Quad I/O operation will be selected and initial delay is the same as dummy cycles of Fast Read Quad I/O Read operation. If it is set to “0”, Fast Read operation will be applied and initial delay is the same as dummy cycles of Fast Read operation.
- The starting address of the boot code is selected by the value programmed into the AutoBoot Start Address (ABSA) field of the AutoBoot Register.
  - Data will continuously shift out until CE# returns high.
- At any point after the first data byte is transferred, when CE# returns high, the SPI device will reset to standard SPI mode; able to accept normal command operations.
  - A minimum of one byte must be transferred.
  - AutoBoot mode will not initiate again until another power cycle or a reset occurs.
- An AutoBoot Enable bit (ABE) is set to enable the AutoBoot feature.

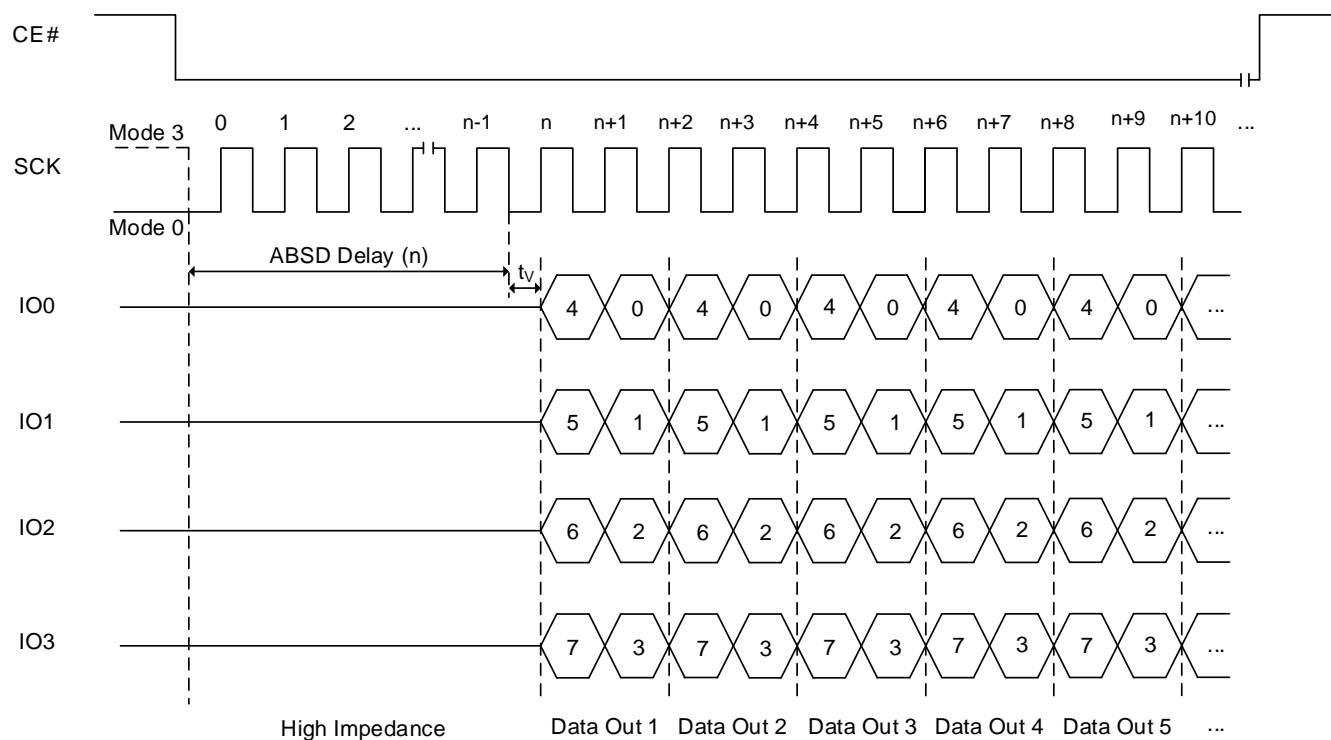
The AutoBoot register bits are non-volatile and provide:

- The starting address set by the AutoBoot Start Address (ABSA).
- The number of initial delay cycles, set by the AutoBoot Start Delay (ABSD) 4-bit count value.

**Figure 8.102 AutoBoot Sequence (QE = 0)**



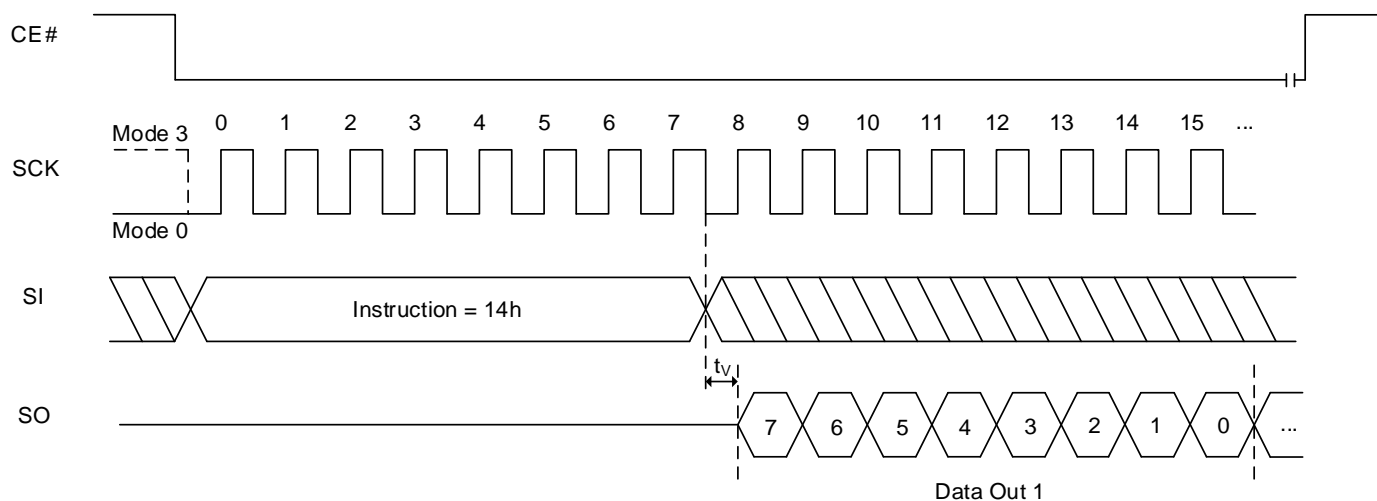
**Figure 8.103 AutoBoot Sequence (QE = 1)**



**AUTOBOOT REGISTER READ OPERATION (RDABR, 14h)**

The AutoBoot Register Read command is shifted in. Then the 32bit AutoBoot Register is shifted out, least significant byte first, most significant bit of each byte first. It is possible to read the AutoBoot Register continuously by providing multiples of 32bits.

RDABR operation is valid only at SPI mode only.

**Figure 8.104 RDABR Sequence In SPI Mode**

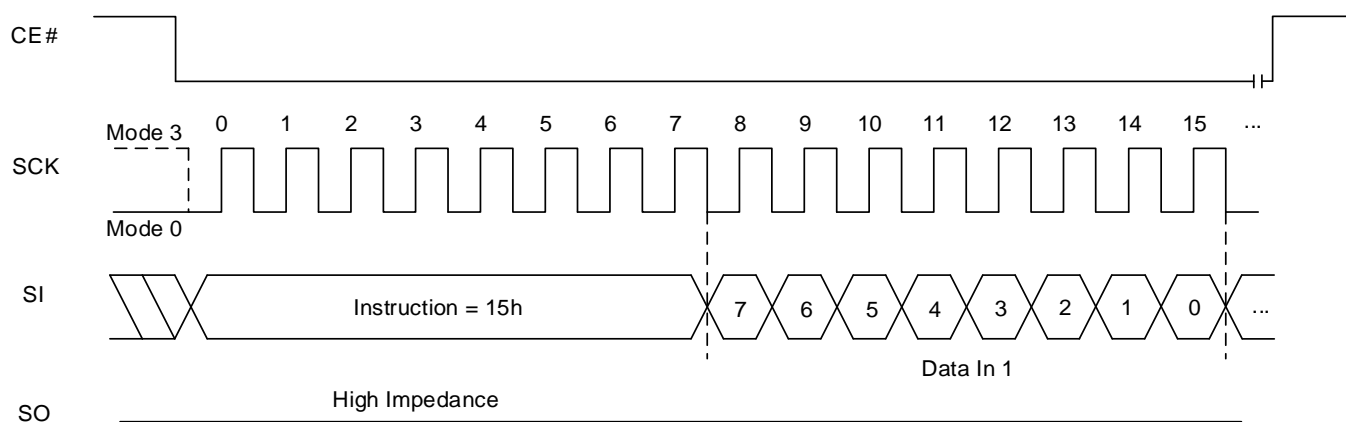
### AUTOBOOT REGISTER WRITE OPERATION (WRABR, 15h)

Before the WRABR command can be accepted, a Write Enable (WREN) command must be issued and decoded by the device, which sets the Write Enable Latch (WEL) in the Status Register to enable any write operations.

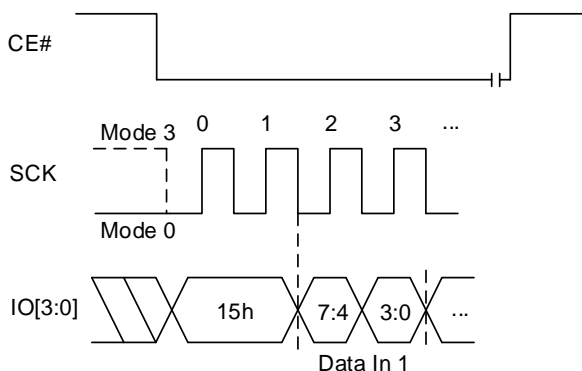
The WRABR command is entered by shifting the instruction and the data bytes, least significant byte first, most significant bit of each byte first. The WRABR data is 32bits in length.

CE# must be driven high after the 32nd bit of data has been latched. If not, the WRABR command is not executed. As soon as CE# is driven high, the WRABR operation is initiated. While the WRABR operation is in progress, Status Register or Extended Read Register may be read to check the value of the Write In Progress (WIP) bit. The WIP bit is "1" during the WRABR operation, and is "0" when it is completed. When the WRABR cycle is completed, the WEL is set to "0".

**Figure 8.105 WRABR Sequence In SPI Mode**



**Figure 8.106 WRABR Sequence In QPI Mode**



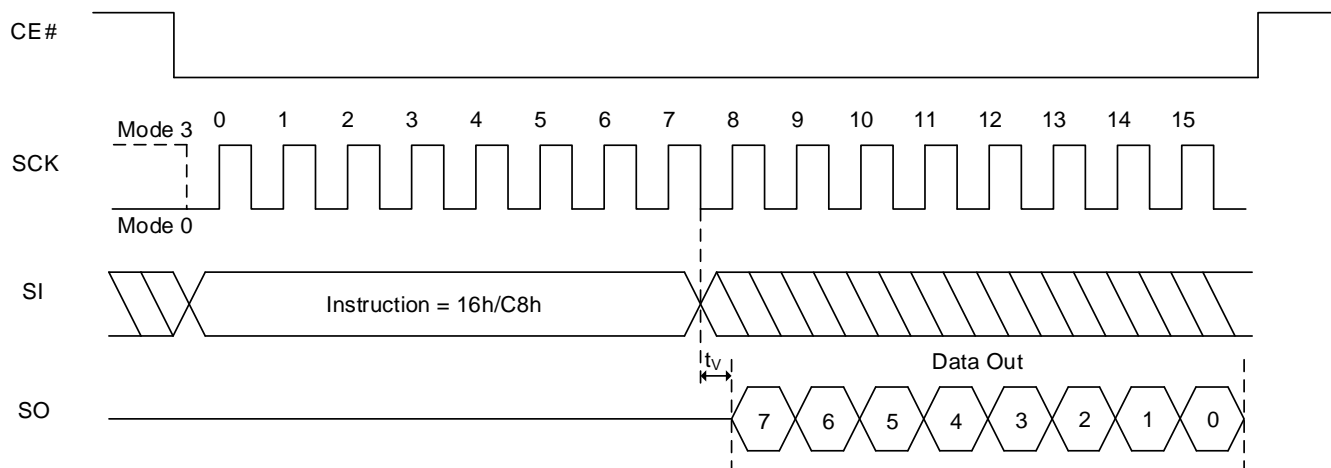
#### 8.47 READ BANK ADDRESS REGISTER OPERATION (RDBR: 16h/C8h)

The Read Bank Address Register (RDBR) instruction allows the Bank Address Register contents to be read. RDBR is used to read only a volatile Bank Address Register.

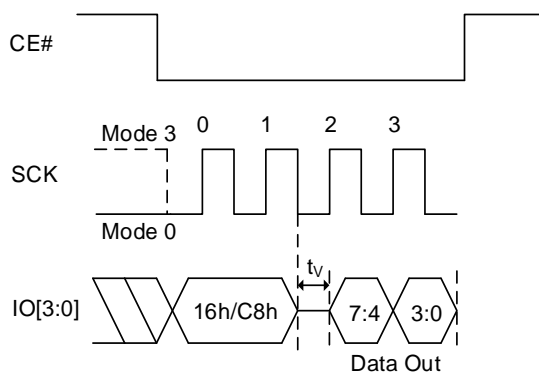
The instruction code is first shifted in. Then the 8-bit Bank Register is shifted out. It is possible to read the Bank Address Register continuously by providing multiples of eight bits.

Data is shifted in from SI and data is shifted out from SO in SPI sequence whereas data in and out is via four pins (IO0-IO3) in QPI sequence.

**Figure 8.107 Read Bank Address Register Sequence In SPI Mode**



**Figure 8.108 Read Bank Address Register Sequence In QPI Mode**



#### 8.48 WRITE BANK ADDRESS REGISTER OPERATION (WRBRNV: 18h, WRBRV: 17h/C5h)

The Write Bank Address Register (WRBRNV and WRBRV) instruction is used to write address bits above A23, into the Bank Address Register (BAR). WRBRNV is used to write the non-volatile Bank Address Register and WRBRV is used to write the volatile Bank Address Register. The instruction is also used to write the Extended Address Control bit (EXTADD) that is also in BAR [7]. BAR provides the high order addresses needed by devices having more than 128Mbits (16Mbytes), when using 3-byte address commands without extended addressing enabled (BAR [7] EXTADD = 0).

To write non-volatile Bank Address Register, a standard Write Enable (06h) instruction must previously have been executed for the device to accept WRBRNV(18h) instruction (Status Register bit WEL must equal "1").

To write volatile Bank Address Register, C5h or 17h command can be used.

When using C5h instruction, a standard Write Enable (06h) instruction must previously have been executed for the device to accept C5h instruction (Status Register bit WEL must equal "1").

But 17h instruction does not require a standard Write Enable (06h) operation. (Status Register bit WEL remains "0").

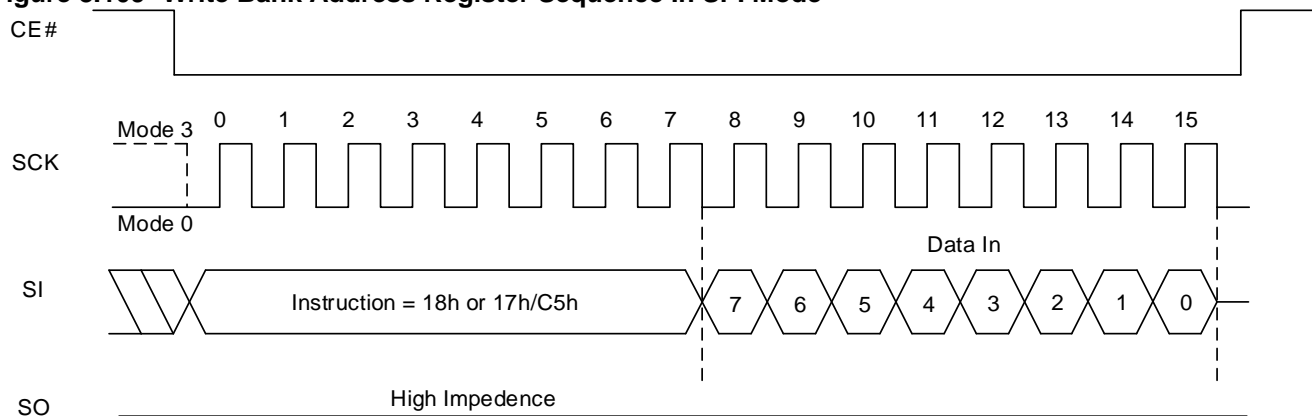
The WRBRNV/WRBRV instructions are followed by the data byte. The Bank Address Register is one data byte in length. The Write In Progress (WIP) bit is "1" during WRBRNV/WRBRV operation, and is "0" when it is completed. Any bank address bit reserved for the future should always be written as "0". Data is shifted in from SI and in SPI whereas data is shifted in via four pins (IO0-IO3) in QPI.

Bit 7 (EXTADD) of volatile Bank Address Register is also writable with EN4B (B7h)/EX4B (29h) instruction.

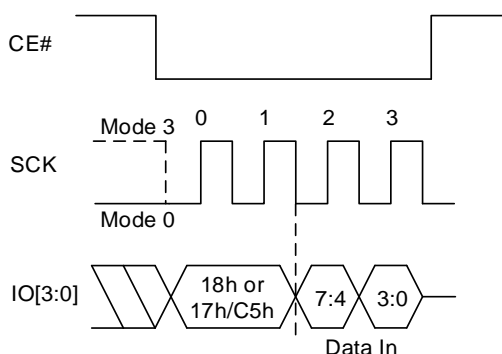
But B7h/29h instruction does not require a standard Write Enable (06h) (Status Register bit WEL remains 0).

**Note: When WRBRNV is executed, the volatile Bank Address Register is set as well as the non-volatile Bank Address Register.**

**Figure 8.109 Write Bank Address Register Sequence In SPI Mode**



**Figure 8.110 Write Bank Address Register Sequence In QPI Mode**



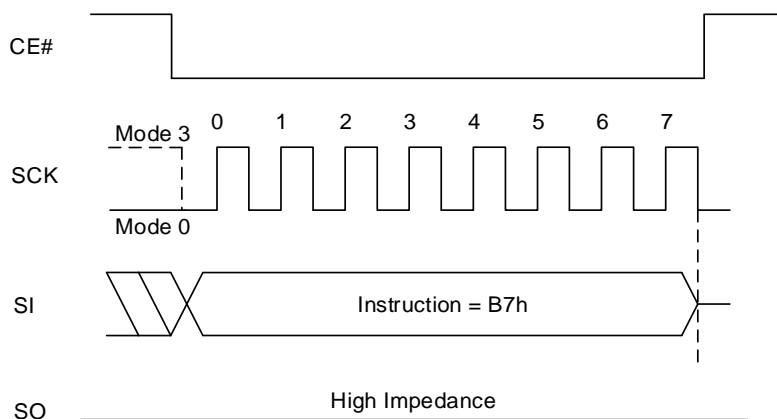
### 8.49 ENTER 4-BYTE ADDRESS MODE OPERATION (EN4B, B7h)

The Enter 4-byte Address Mode instruction allows 32bit address (A31-A0) to be used to access the memory array beyond 128Mb. To execute EN4B operation, the host drives CE# low, sends the instruction code and then drives CE# high. The Exit 4-byte Address Mode instruction can be used to exit the 4-byte address mode.

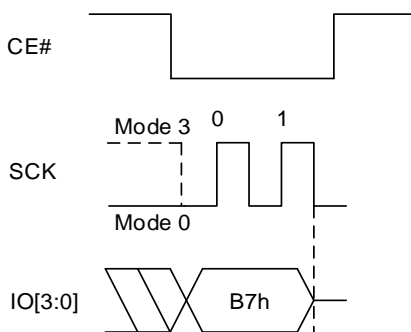
A Write Enable (WREN, 06h) command is not required prior to EN4B (B7h) command.

**Note:** The EN4B instruction will set the Bit 7 (EXTADD) of the volatile Bank Address Register to “1”, but will not change the non-volatile Bank Address Register.

**Figure 8.111 Enter 4-byte Address Mode Sequence In SPI Mode**



**Figure 8.112 Enter 4-byte Address Mode Sequence In QPI Mode**



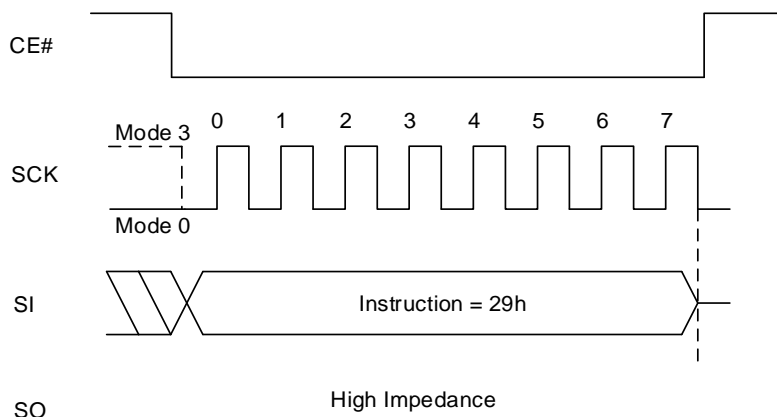
### 8.50 EXIT 4-BYTE ADDRESS MODE OPERATION (EX4B, 29h)

In order to be backward compatible, the Exit 4-byte Address Mode instruction allows 24bit address (A23-A0) to be used to access the memory array up to 128Mb. The Bank Address Register must be used to access the memory array beyond 128Mb. To execute EX4B operation, the host drives CE# low, sends the instruction code and then drives CE# high.

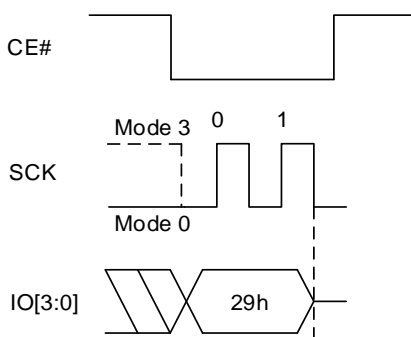
A Write Enable (WREN, 06h) command is not required prior to EX4B (29h) command.

**Note: The EX4B instruction will reset the Bit 7 (EXTADD) of the volatile Bank Address Register to “0” , but will not change the non-volatile Bank Address Register.**

**Figure 8.113 Exit 4-byte Address Mode Sequence In SPI Mode**



**Figure 8.114 Exit 4-byte Address Mode Sequence In QPI Mode**





### 8.51 READ DYB OPERATION (RDDYB, FAh or 4RDDYB, E0h)

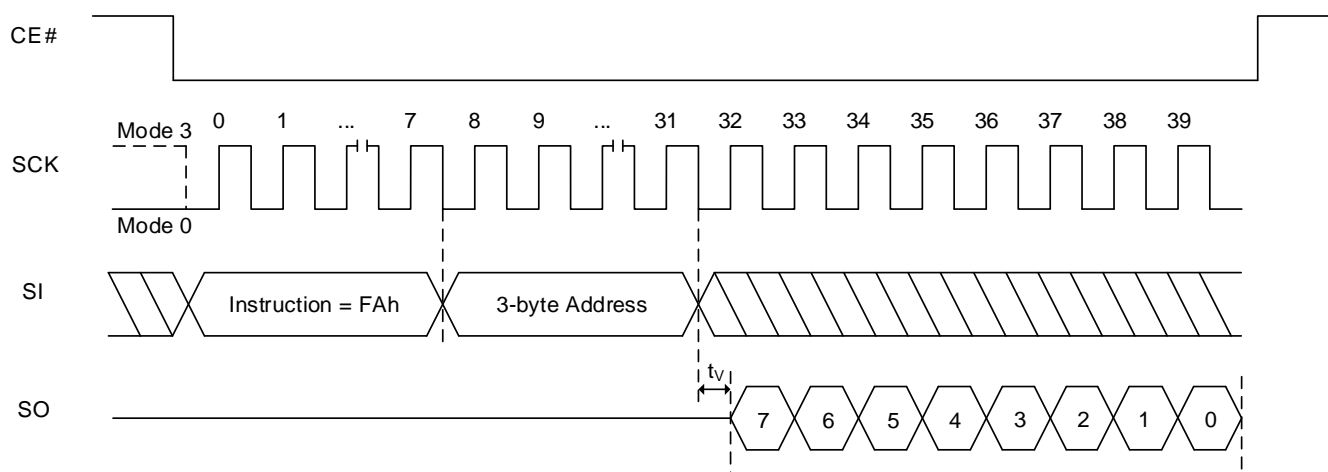
FAh (EXTADD=0) is followed by a 3-byte address (A23-A0) or

- FAh (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- E0h is followed by a 4-byte address (A31-A0)

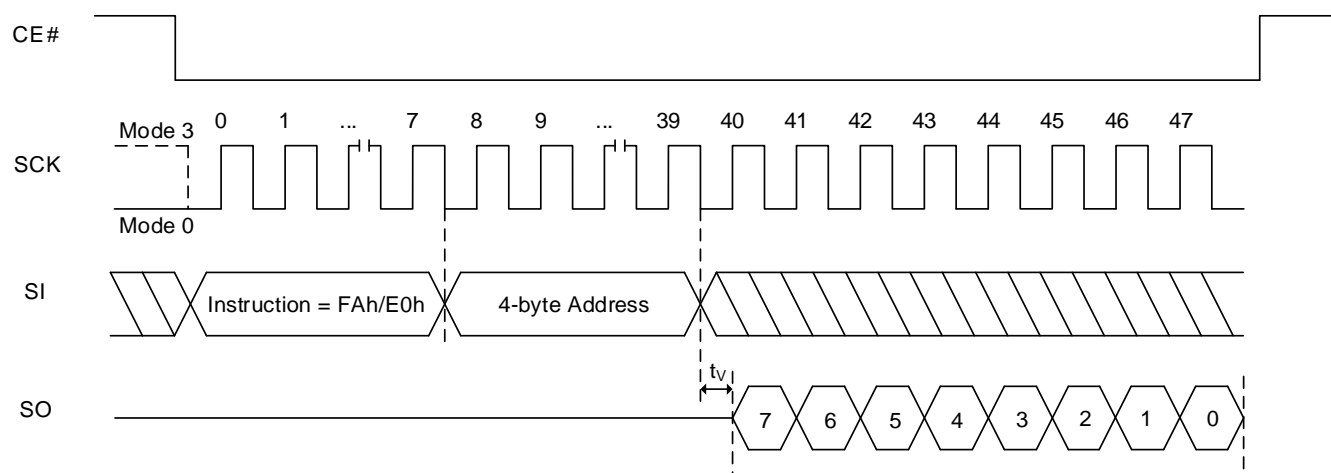
The instruction is used to read Dynamic Protection Bit (DYB) status of the given sector/block. The instruction code is entered first, followed by the 24-bit or 32-bit address selecting location zero within the desired sector/block as above. Then the 8-bit DYB access register contents are shifted out. Each bit (SPI) is shifted out at the SCK frequency by the falling edge of the SCK signal. It is possible to read the same DYB access register continuously by providing multiples of eight bits. The address of the DYB register does not increment so this is not a means to read the entire DYB array. Each location must be read with a separate Read DYB instruction.

**Note: Data must be either 00h (protected) or FFh (unprotected).**

**Figure 8.115 Read DYB Sequence In SPI Mode (FAh [EXTADD=0], 3-byte address)**



**Figure 8.116 Read DYB Sequence In SPI Mode(FAh [EXTADD=1] or E0h, 4-byte address)**



**8.52 WRITE DYB OPERATION (WRDYB, FBh or 4WRDYB, E1h)**

Before the WRDYB/4WRDYB command can be accepted by the device, a standard Write Enable (06h) instruction must previously have been executed for the device to accept Write DYB instruction (Status Register bit WEL must equal 1).

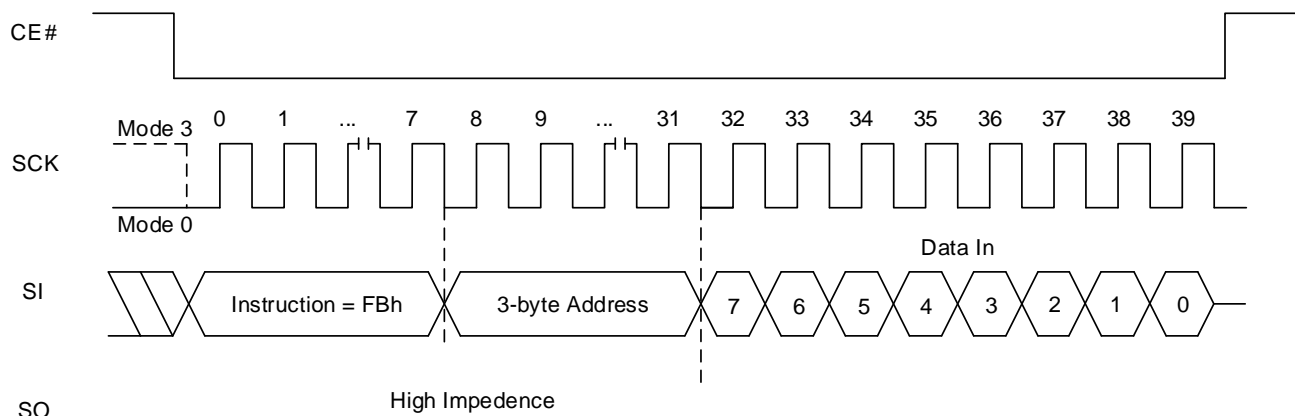
- FBh (EXTADD=0) is followed by a 3-byte address (A23-A0) or
- FBh (EXTADD=1) is followed by a 4-byte address (A31-A0) or
- E1h is followed by a 4-byte address (A31-A0)

The WRDYB/4WRDYB command is entered by driving CE# low, followed by the instruction code, the 24-bit or 32-bit address selecting location zero within the desired sector/block as above, then the data byte. The DYB Access Register is one data byte in length.

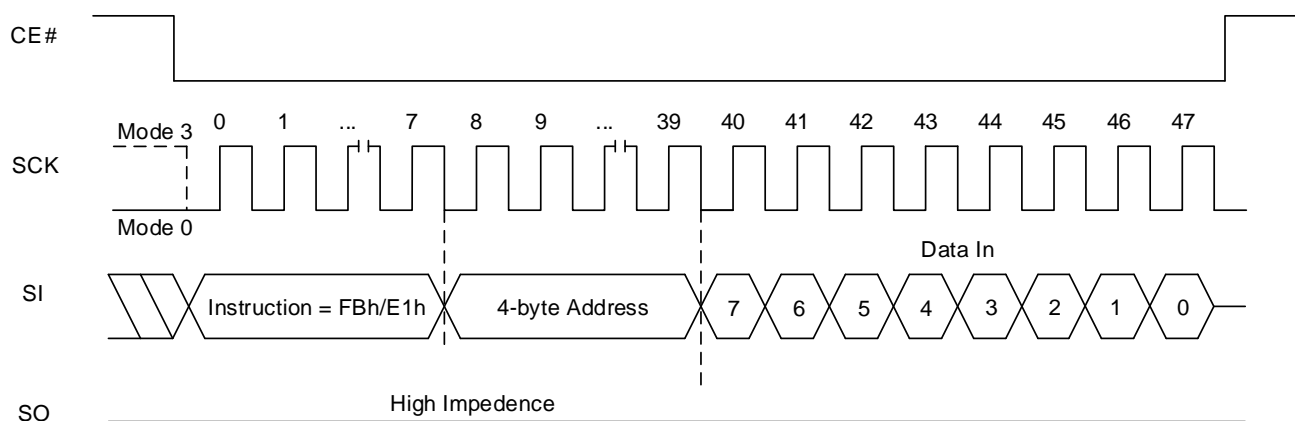
CE# must be driven high after the eighth bit of data has been latched in. As soon as CE# is driven high, the WRDYB/4WRDYB operation is initiated.

**Note: Data must be either 00h (protected) or FFh (unprotected).**

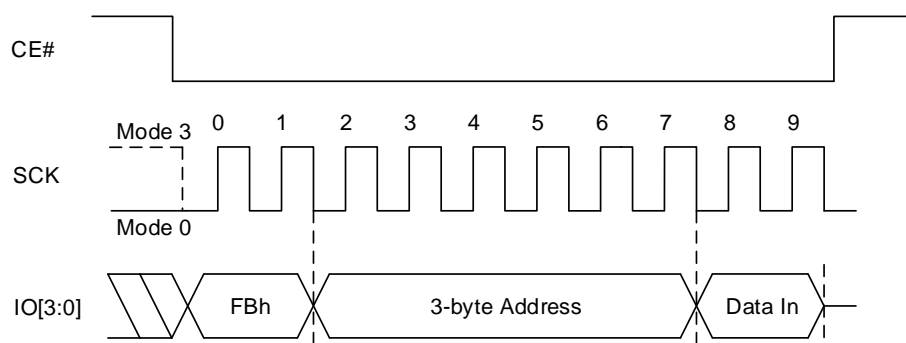
**Figure 8.117 Write DYB Sequence In SPI Mode (FBh [EXTADD=0], 3-byte address)**



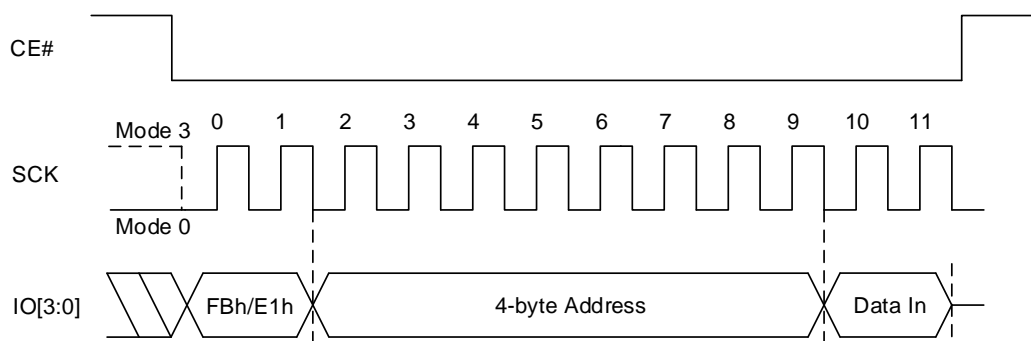
**Figure 8.118 Write DYB Sequence In SPI Mode (FBh [EXTADD=1] or E1h, 4-byte address)**



**Figure 8.119 Write DYB Sequence In QPI Mode (FBh [EXTADD=0], 3-byte address)**



**Figure 8.120 Write DYB Sequence In QPI Mode (FBh [EXTADD=1] or E1h, 4-byte address)**



### 8.53 READ PPB OPERATION (RDPPB, FCh or 4RDPPB, E2h)

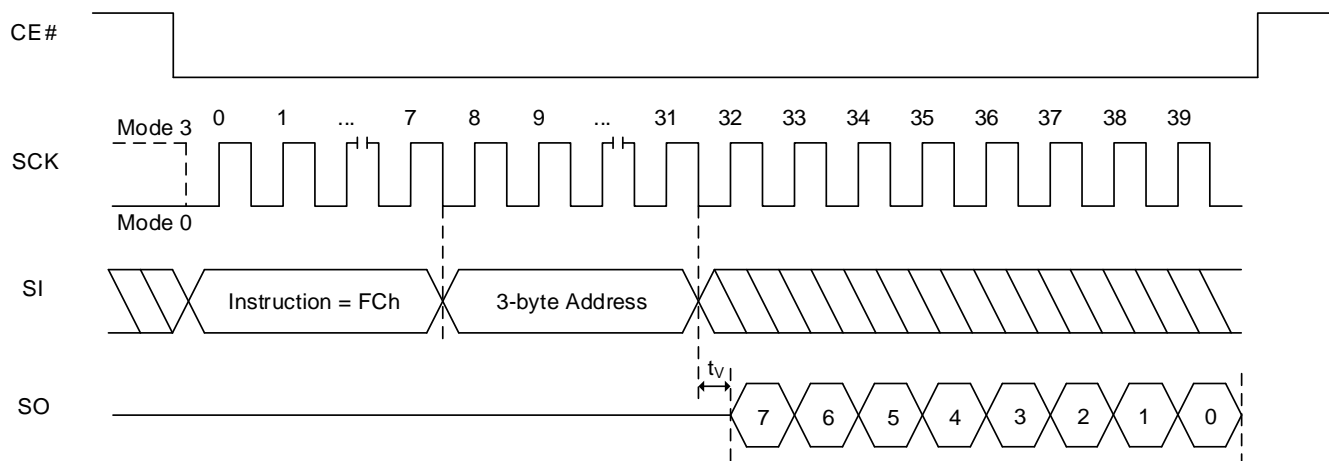
FCh (EXTADD=0) is followed by a 3-byte address (A23-A0) or  
 FCh (EXTADD=1) is followed by a 4-byte address (A31-A0) or  
 E2h is followed by a 4-byte address (A31-A0)

The instruction code is shifted into SI by the rising edges of the SCK signal, followed by the 24-bit or 32-bit address selecting location zero within the desired sector/block as above. Then the 8-bit PPB Access Register contents are shifted out on SO. The RDPPB/4RDPPB is supporting only SPI, not supporting QPI.

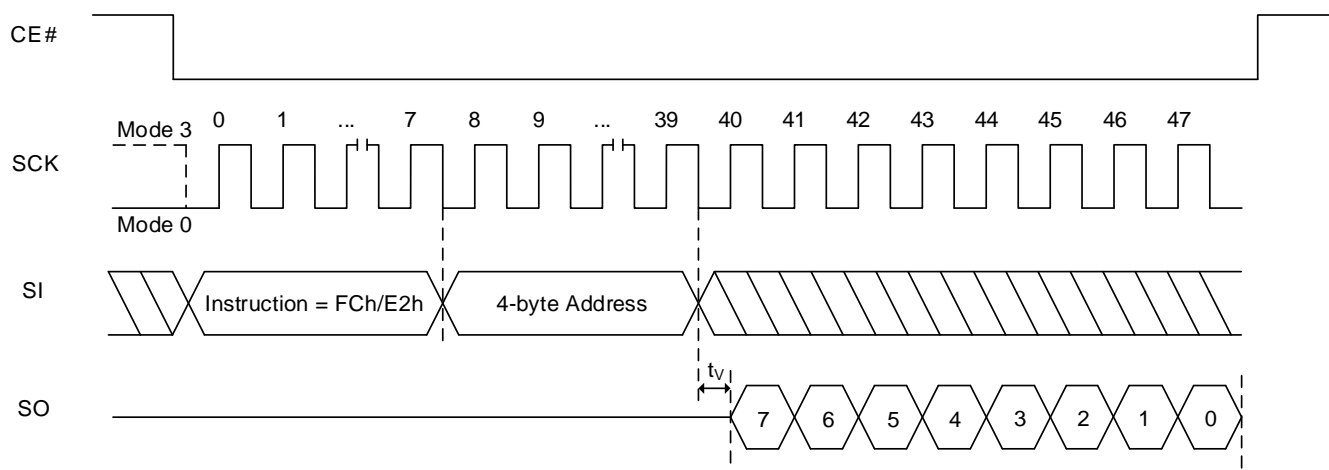
It is possible to read the same PPB Access Register continuously by providing multiples of eight bits. The address of the PPB Access Register does not increment so this is not a means to read the entire PPB array. Each location must be read with a separate Read PPB command.

**Note:** Data must be either 00h (protected) or FFh (unprotected).

**Figure 8.121 Read PPB Sequence (FCh [EXTADD=0], 3-byte address)**



**Figure 8.122 Read PPB Sequence (FCh [EXTADD=1] or E2h, 4-byte address)**



### 8.54 PROGRAM PPB OPERATION (PGPPB, FDh or 4PGPPB, E3h)

Before the Program PPB (PGPPB/4PGPPB) command is sent, a Write Enable (WREN) command must be issued. After the WREN command has been decoded, the device will set the Write Enable Latch (WEL) in the Status Register.

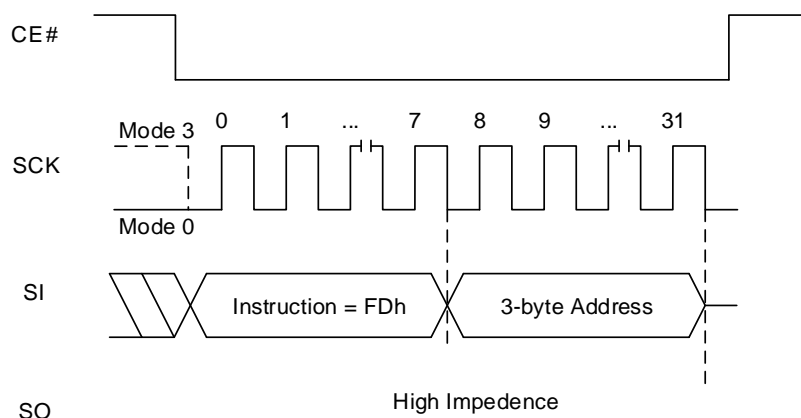
FDh (EXTADD=0) is followed by a 3-byte address (A23-A0) or  
FDh (EXTADD=1) is followed by a 4-byte address (A31-A0) or  
E3h is followed by a 4-byte address (A31-A0)

The PGPPB/4PGPPB command is entered by driving CE# low, followed by the instruction code, followed by the 24-bit or 32-bit address selecting location zero within the desired sector/block as above.

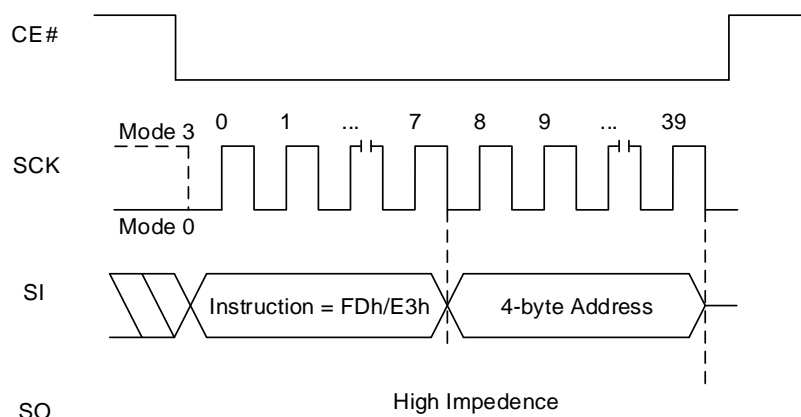
The PGPPB/4PGPPB command affects the WIP bit in the same manner as any other programming operation. CE# must be driven high after the last bit of address has been latched in. As soon as CE# is driven high, the PGPPB/4PGPPB operation is initiated. While the PGPPB/4PGPPB operation is in progress, the Status Register or Extended Read Register may be read to check the value of the Write In Progress (WIP) bit. The WIP bit is "1" during the PGPPB/4PGPPB operation, and is "0" when it is completed. When the PGPPB/4PGPPB operation is completed, the WEL is set to "0".

**Note:** Data must be either 00h (protected) or FFh (unprotected).

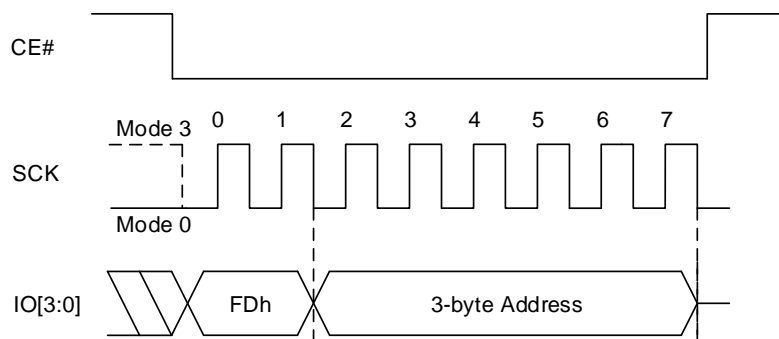
**Figure 8.123 Program PPB Sequence In SPI Mode (FDh [EXTADD=0], 3-byte address)**



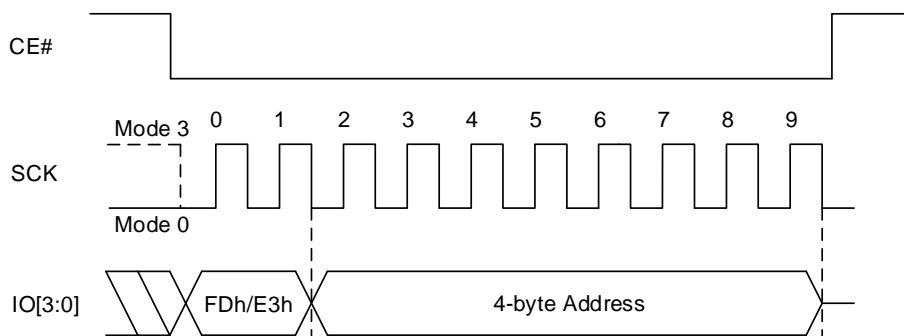
**Figure 8.124 Program PPB Sequence In SPI Mode (FDh [EXTADD=1] or E3h, 4-byte address)**



**Figure 8.125 Program PPB Sequence In QPI Mode (FDh [EXTADD=0], 3-byte address)**



**Figure 8.126 Program PPB Sequence In QPI Mode (FDh [EXTADD=1] or E3h, 4-byte address)**



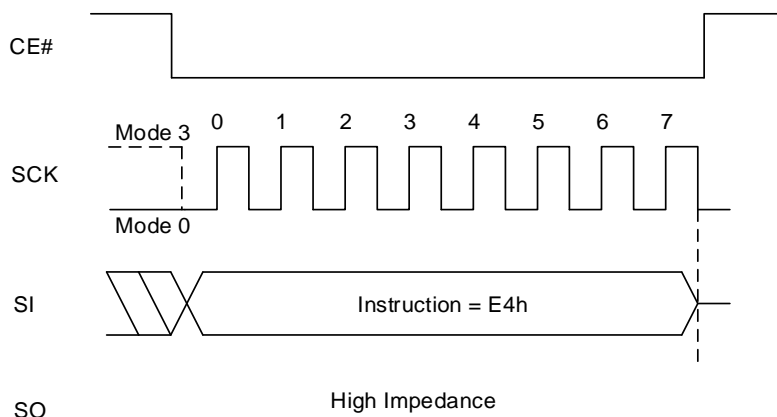
### 8.55 ERASE PPB OPERATION (ERPPB, E4h)

The Erase PPB (ERPPB) command sets all PPB bits to “1”. Before the ERPPB command can be accepted by the device, a Write Enable (WREN) command must be issued and decoded by the device, which sets the Write Enable Latch (WEL) in the Status Register to enable any write operations.

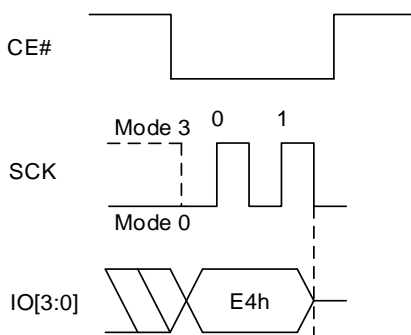
The instruction code is shifted in by the rising edges of the SCK signal. CE# must be driven high after the eighth bit of the instruction byte has been latched in. This will initiate the beginning of internal erase cycle, which involves the pre-programming and erase of the entire PPB memory array. Without CE# being driven high after the eighth bit of the instruction, the PPB erase operation will not be executed.

With the internal erase cycle in progress, the user can read the value of the Write In Progress (WIP) bit to check if the operation has been completed. The WIP bit will indicate “1” when the erase cycle is in progress and “0” when the erase cycle has been completed. When the ERPPB operation is completed, the WEL is set to “0”. Erase suspend is not allowed during PPB Erase.

**Figure 8.127 Erase PPB Sequence In SPI Mode**



**Figure 8.128 Erase PPB Sequence In QPI Mode**

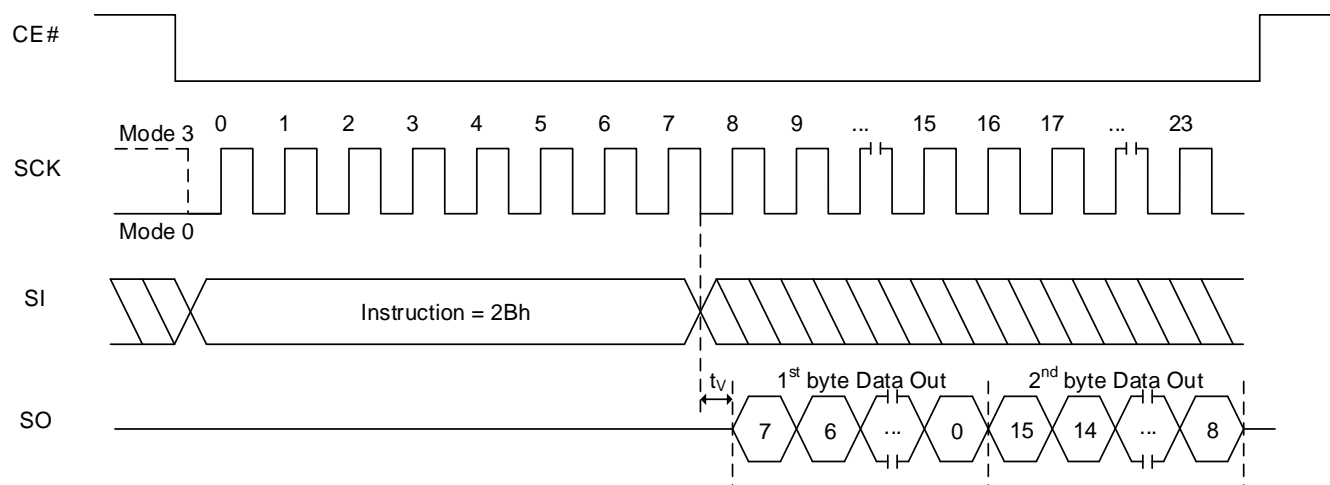




### 8.56 READ ASP OPERATION (RDASP, 2Bh)

The RDASP instruction code is shifted in by the rising edge of the SCK signal. Then the 16-bit ASP register contents is shifted out, least significant byte first, most significant bit of each byte first. Each bit is shifted out at the SCK frequency by the falling edge of the SCK signal. It is possible to read the ASP register continuously by providing multiples of 16 bits.

**Figure 8.129 Read ASP Sequence In SPI Mode**



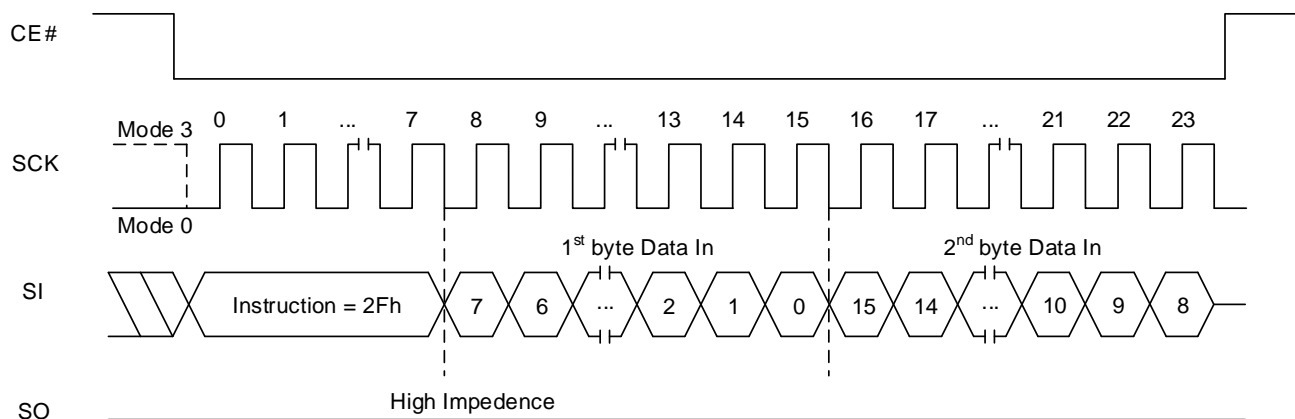
### 8.57 PROGRAM ASP OPERATION (PGASP, 2Fh)

Before the Program ASP (PGASP) command can be accepted by the device, a Write Enable (WREN) command must be issued. After the WREN command has been decoded, the device will set the Write Enable Latch (WEL) in the Status Register to enable any write operations.

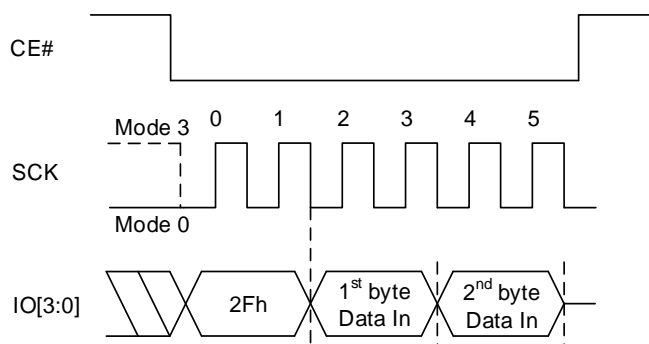
The PGASP command is entered by driving CE# low, followed by the instruction code and two data bytes, least significant byte first, most significant bit of each byte first. The ASP Register is two data bytes in length. The PGASP command affects the Write In Progress (WIP) bit in the same manner as any other programming operation.

CE# input must be driven high after the sixteenth bit of data has been latched in. If not, the PGASP command is not executed. As soon as CE# is driven high, the PGASP operation is initiated. While the PGASP operation is in progress, the Status Register or the Extended Read Register may be read to check the value of WIP bit. The WIP bit is "1" during the PGASP operation, and is "0" when it is completed. When the PGASP operation is completed, the WEL is set to "0".

**Figure 8.130 Program ASP Sequence In SPI Mode**



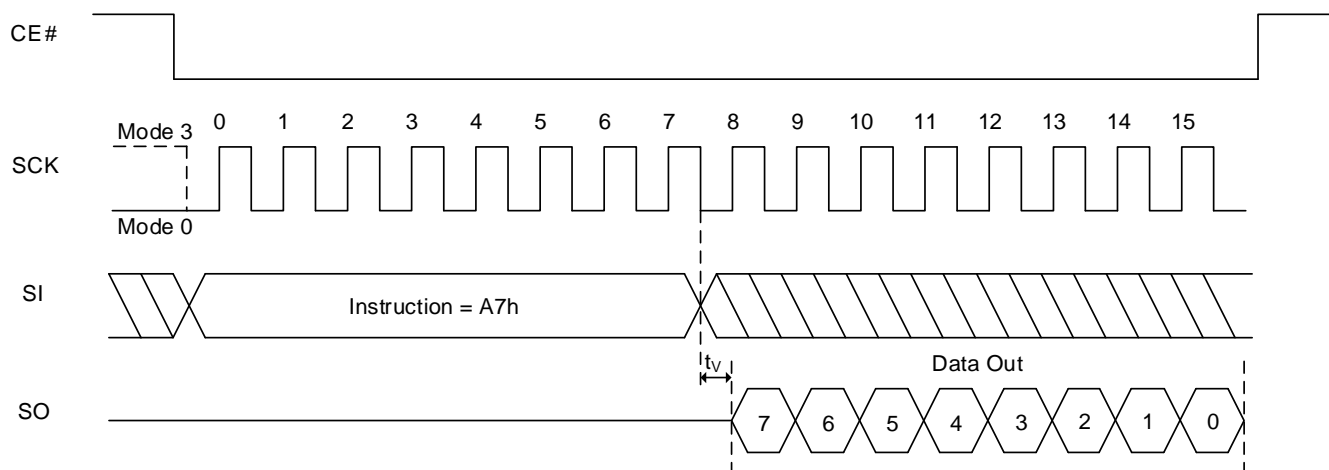
**Figure 8.131 Program ASP Sequence In QPI Mode**



### 8.58 READ PPB LOCK BIT OPERATION (RDPLB, A7h)

The Read PPB Lock Bit (RDPLB) command allows the PPB Lock Register contents to be read. It is possible to read the PPB Lock Register continuously by providing multiples of eight bits. The PPB Lock Register contents may only be read when the device is in standby state with no other operation in progress. It is recommended to check the Write In Progress (WIP) bit before issuing a new command to the device. RDPLB operation is valid only at SPI mode only.

**Figure 8.132 Read PPB Lock Bit Sequence In SPI Mode**

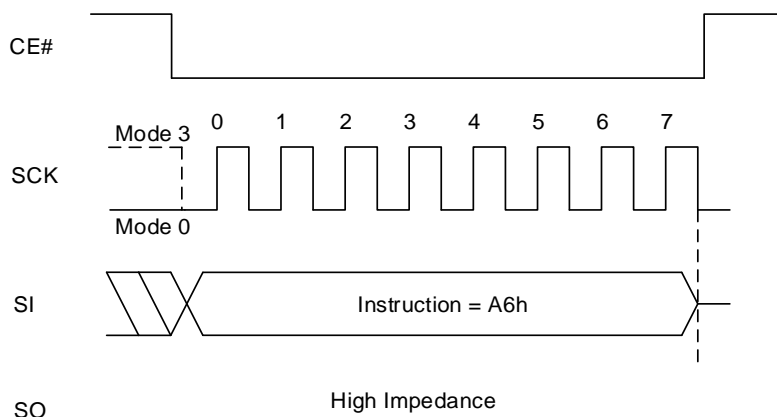


### 8.59 WRITE PPB LOCK BIT OPERATION (WRPLB, A6h)

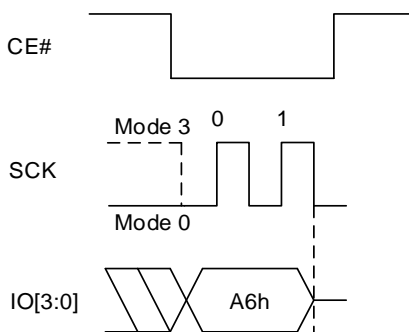
The Write PPB Lock Bit (WRPLB) command clears the PPB Lock (PPBLK) bit to zero. Before the WRPLB command can be accepted by the device, a Write Enable (WREN) command must be issued and decoded by the device, which sets the Write Enable Latch (WEL) in the Status Register to enable any write operations.

The WRPLB command is entered by driving CE# low, followed by the instruction code. CE# must be driven high after the eighth bit of instruction has been latched in. If not, the WRPLB command is not executed. As soon as CE# is driven high, the WRPLB operation is initiated. While the WRPLB operation is in progress, the Status Register or Extended Read Register may still be read to check the value of the Write In Progress (WIP) bit. The WIP bit is "1" during the WRPLB operation, and is "0" when it is completed. When the WRPLB operation is completed, the WEL is set to "0".

**Figure 8.133 Write PPB Lock Bit Sequence In SPI Mode**



**Figure 8.134 Write PPB Lock Bit Sequence In QPI Mode**

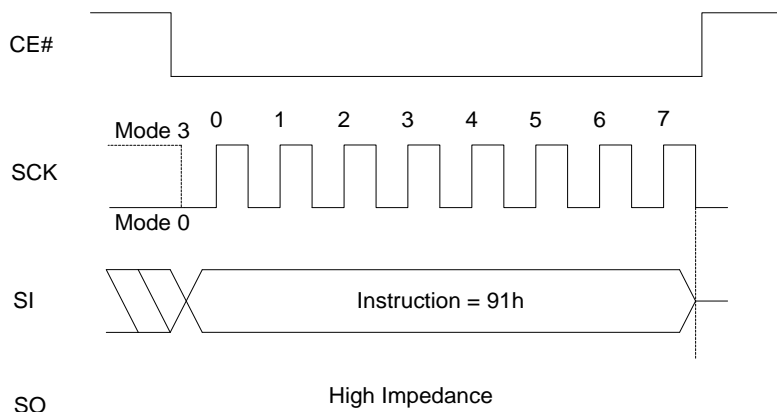


### 8.60 SET FREEZE BIT OPERATION (SFRZ, 91h)

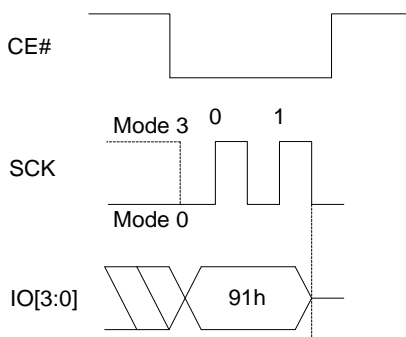
The Set FREEZE Bit (SFRZ) command sets FREEZE (PPB Lock Register bit7) to one. Please refer to the section 6.6.3 PPB Lock Register for more detail. Before the SFRZ command can be accepted by the device, a Write Enable (WREN) command must be issued and decoded by the device, which sets the Write Enable Latch (WEL) in the Status Register to enable any write operations.

The SFRZ command is entered by driving CE# low, followed by the instruction code. CE# must be driven high after the eighth bit of instruction has been latched in. If not, the SFRZ command is not executed. As soon as CE# is driven high, the SFRZ operation is initiated. While the SFRZ operation is in progress, the Status Register or Extended Read Register may still be read to check the value of the Write In Progress (WIP) bit. The WIP bit is "1" during the SFRZ operation, and is "0" when it is completed. When the SFRZ operation is completed, the WEL is set to "0".

**Figure 8.135 Set FREEZE Bit Sequence In SPI Mode**



**Figure 8.136 Set FREEZE Bit Sequence In QPI Mode**



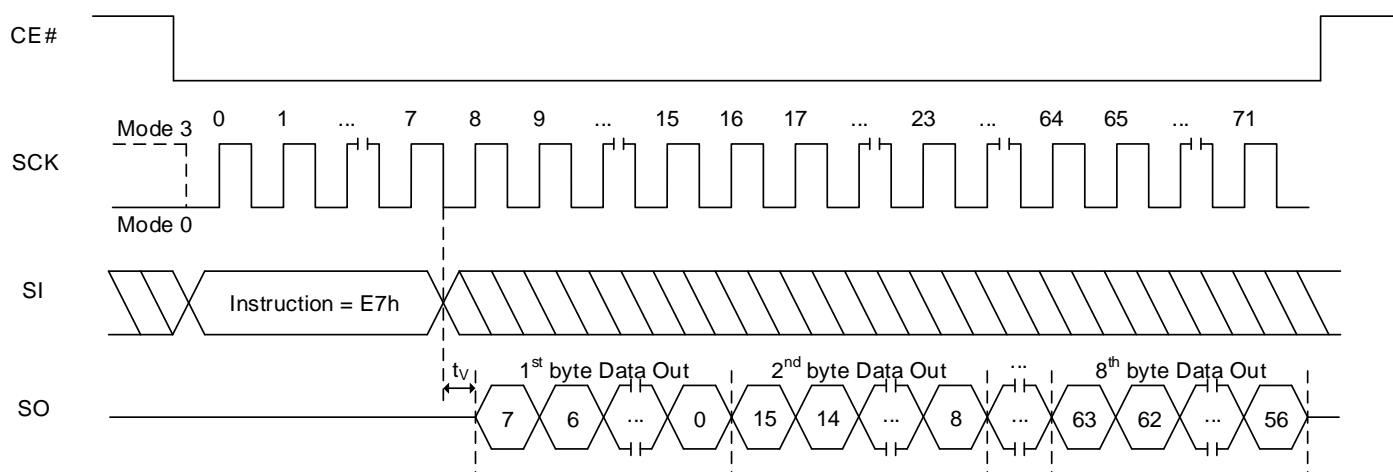
**8.61 READ PASSWORD OPERATION (RDPWD, E7h)**

The correct password value may be read only after it is programmed and before the Password Mode has been selected by programming the Password Protection Mode bit to “0” in the ASP Register (ASP[2]). After the Password Protection Mode is selected the RDPWD command is ignored.

The RDPWD command is shifted in. Then the 64-bit Password is shifted out, least significant byte first, most significant bit of each byte first. Each bit is shifted out at the SCK frequency by the falling edge of the SCK signal. It is possible to read the Password continuously by providing multiples of 64bits.

RDPWD operation is valid only at SPI mode only.

**Figure 8.137 Read password Sequence In SPI Mode**



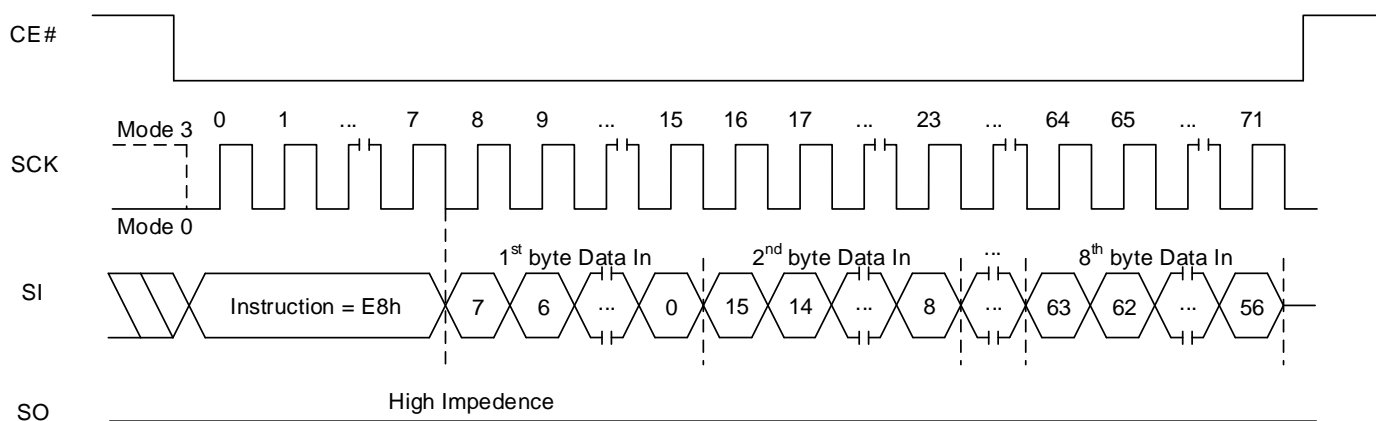
## 8.62 PROGRAM PASSWORD OPERATION (PGPWD, E8h)

Before the Program Password (PGPWD) command can be accepted by the device, a Write Enable (WREN) command must be issued and decoded by the device which sets the Write Enable Latch (WEL) to enable the PGPWD operation. The password can only be programmed before the Password Mode is selected by programming the Password Protection Mode bit to “0” in the ASP Register (ASP[2]). After the Password Protection Mode is selected the PGPWD command is ignored.

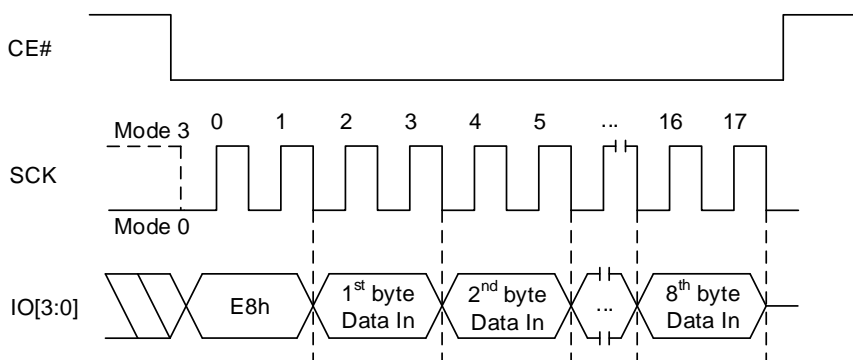
The PGPWD command is entered by driving CE# low, followed by the instruction code and the password data bytes, least significant byte first, most significant bit of each byte first. The password is 64bits in length.

CE# must be driven high after the 64<sup>th</sup> bit of data has been latched. If not, the PGPWD command is not executed. As soon as CE# is driven high, the PGPWD operation is initiated. While the PGPWD operation is in progress, the Status Register or Extended Read Register may be read to check the value of the Write In Progress (WIP) bit. The WIP bit is “1” during the PGPWD operation, and is “0” when it is completed. When the PGPWD operation is completed, the Write Enable Latch (WEL) is set to “0”.

**Figure 8.138 Program Password Sequence In SPI Mode**



**Figure 8.139 Program Password Sequence In QPI Mode**



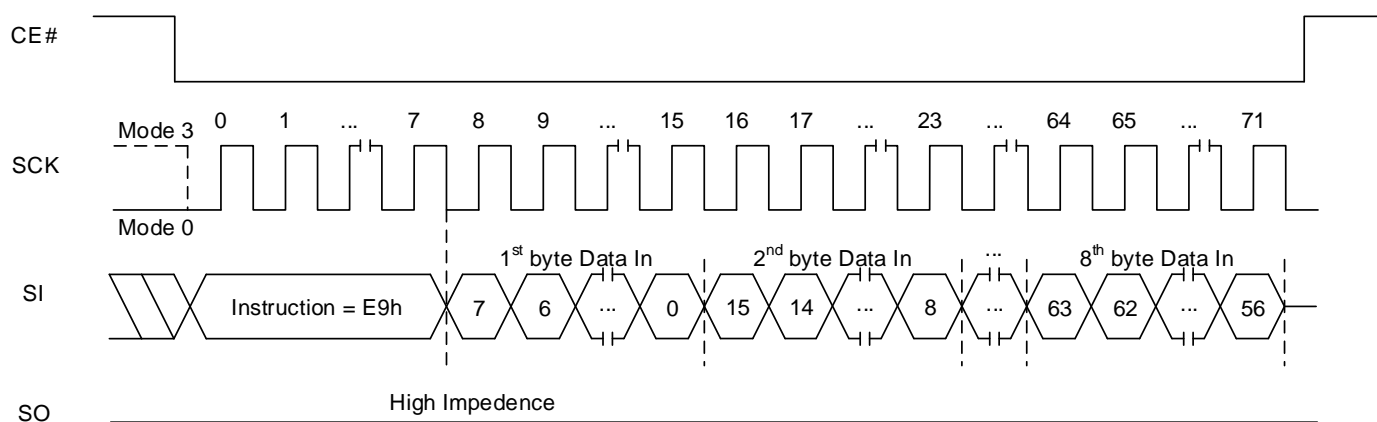
### 8.63 UNLOCK PASSWORD OPERATION (UNPWD, E9h)

The UNPWD command is entered by driving CE# low, followed by the instruction code and the password data bytes, least significant byte first, most significant bit of each byte first. The password is 64bits in length.

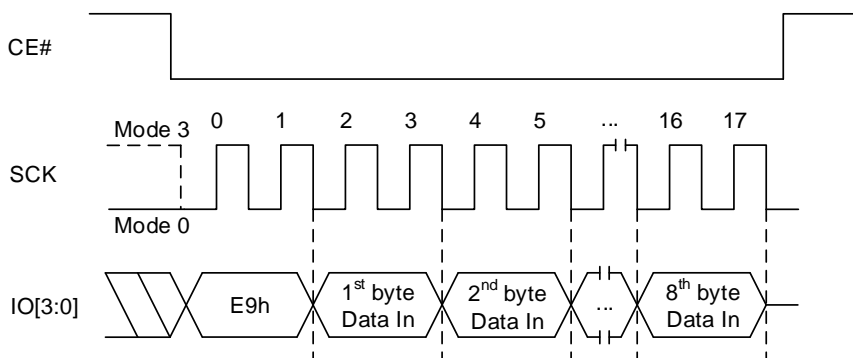
CE# must be driven high after the 64<sup>th</sup> bit of data has been latched. If not, the UNPWD command is not executed. As soon as CE# is driven high, the UNPWD operation is initiated. While the UNPWD operation is in progress, the Status Register or Extended Read Register may be read to check the value of the Write In Progress (WIP) bit. The WIP bit is “1” during the UNPWD operation, and is “0” when it is completed.

If the UNPWD command supplied password does not match the hidden password in the Password Register, the UNPWD command is ignored. This returns the device to standby state, ready for a new command such as a retry of the UNPWD command. If the password does match, the PPB Lock bit is set to “1”.

**Figure 8.140 Unlock Password Sequence In SPI Mode**



**Figure 8.141 Unlock Password Sequence In QPI Mode**





### 8.64 GANG SECTOR/BLOCK LOCK OPERATION (GBLK, 7Eh)

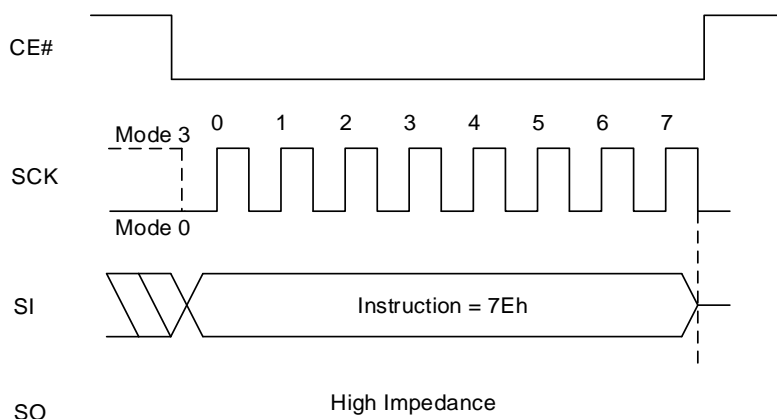
The Gang Sector/Block Lock (GBLK) instruction provides a quick method to set all DYB (Dynamic Protection Bit) bits to “0” at once.

Before the GBLK (7Eh) command can be accepted by the device, a Write Enable (WREN) command must be issued and decoded by the device, which sets the Write Enable Latch (WEL) in the Status Register to enable any write operations.

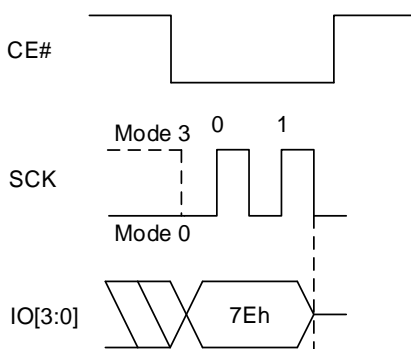
The sequence of issuing GBLK instruction is: drive CE# low → send GBLK instruction code → drive CE# high. The instruction code will be shifted into the device on the rising edge of SCK.

The GBLK command is accepted in both SPI and QPI mode. The CE# must go high exactly at the byte boundary, otherwise, the instruction will be ignored. While the GBLK operation is in progress, the Status Register or Extended Read Register may be read to check the value of the Write In Progress (WIP) bit. The WIP bit is “1” during the GBLK operation, and is “0” when it is completed.

**Figure 8.142 Gang Sector/Block Lock Sequence In SPI Mode**



**Figure 8.143 Gang Sector/Block Lock Sequence In QPI Mode**



### 8.65 GANG SECTOR/BLOCK UNLOCK OPERATION (GBUN, 98h)

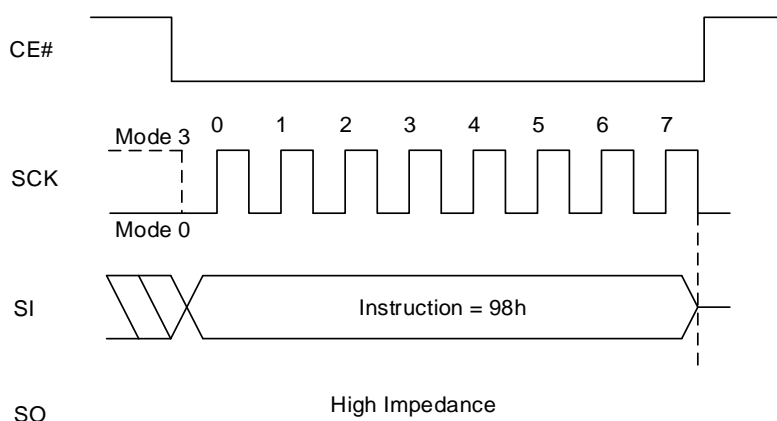
The Gang Sector/Block Unlock (GBUN) instruction provides a quick method to clear all DYB (Dynamic Protection Bit) bits to “1” at once.

Before the GBUN (98h) command can be accepted by the device, a Write Enable (WREN) command must be issued and decoded by the device, which sets the Write Enable Latch (WEL) in the Status Register to enable any write operations.

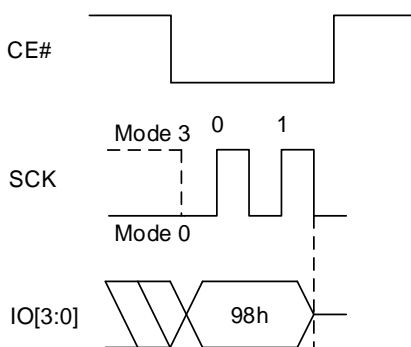
The sequence of issuing GBUN instruction is: drive CE# low → send GBUN instruction code → drive CE# high. The instruction code will be shifted into the device on the rising edge of SCK.

The GBUN command is accepted in both SPI and QPI mode. The CE# must go high exactly at the byte boundary, otherwise, the instruction will be ignored and not be executed. While the GBUN operation is in progress, the Status Register or Extended Read Register may be read to check the value of the Write In Progress (WIP) bit. The WIP bit is “1” during the GBUN operation, and is “0” when it is completed.

**Figure 8.144 Gang Sector/Block Unlock Sequence In SPI Mode**



**Figure 8.145 Gang Sector/Block Unlock Sequence In QPI Mode**



### 8.66 DATA LEARNING PATTERN (DLP)

The Data Learning Pattern is a preamble, and it can help host controller to determine the phase shift from SCK to data edges so that controller can capture data at the center of the data eye at high frequency DTR operation.

DLP can be enabled or disabled by setting the bit EB4 of Extended Read register (DLPEN bit).

When enabled, the device drives all IO bus during last 4 dummy cycles immediately before data is output in DTR operation. (No DLP operation in STR operation). The host must be sure to stop driving the IO bus at least 1 cycle prior to the time that the memory starts outputting the DLP during dummy cycles. That is why, when using DDR IO commands with DLP enabled, 5 or more dummy cycles must be selected.

If there are mode bits for AX read operation, still more than 5 clock cycles are required for DLP operation between last mode bit and first read data.

DLP operation is valid in DTR operation with valid dummy cycles in SPI mode and QPI mode with below command; FRDTR (Fast Read DTR) command in SPI and QPI mode.

FRDDTR (Fast Read Dual IO DTR) command in SPI mode only

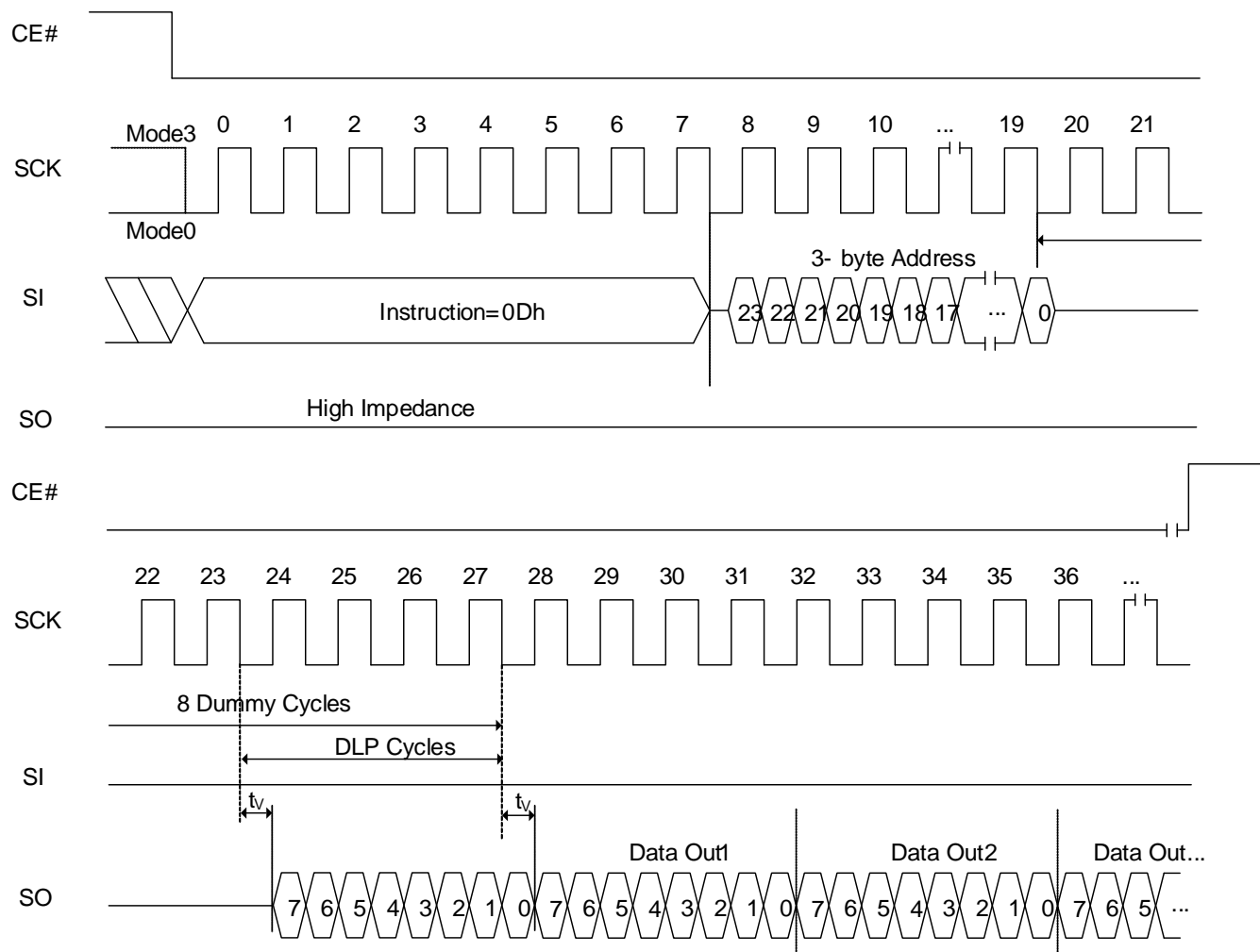
FRQDTR (Fast Read Quad IO DTR) command in SPI and QPI mode

Data Learning Pattern is programmable, and fixed length of 8-bit in DTR operation.

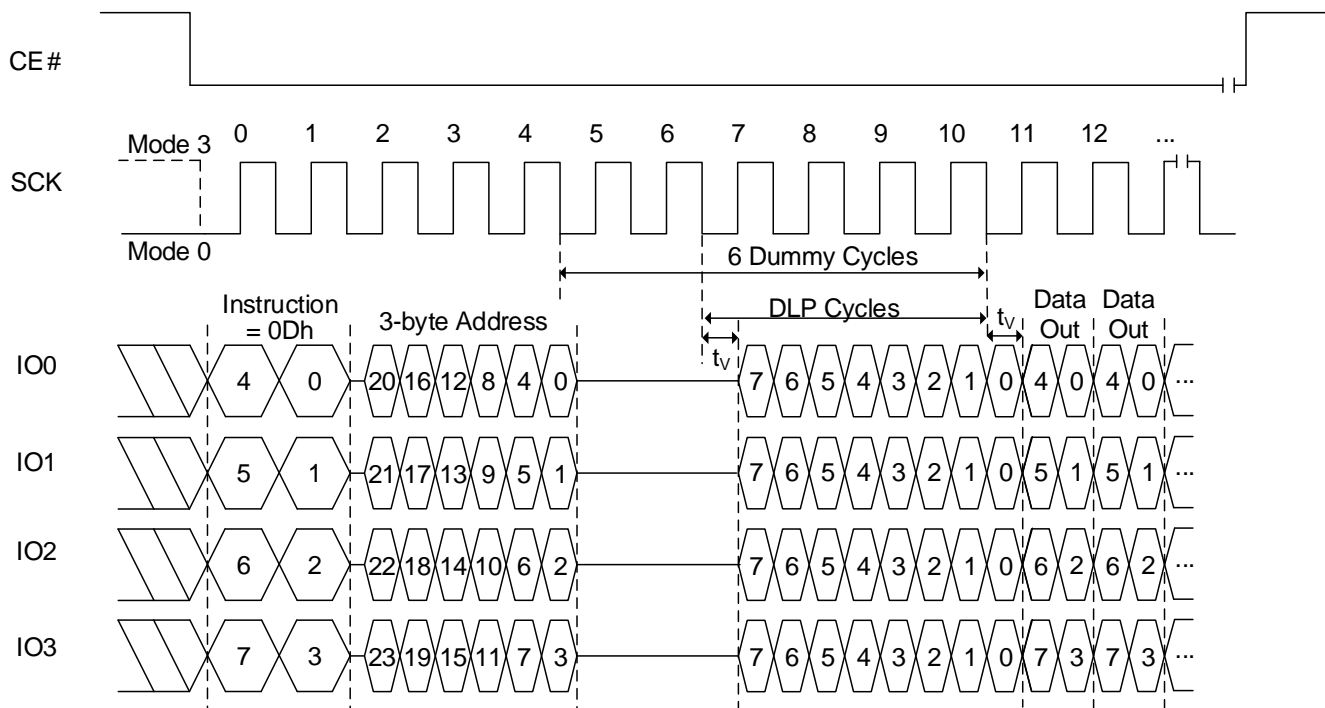
For example, DLP 34h is programmed, then (0011 0100) data will start output from all IOs before data is output during 4 clock cycles. All IOs will output 0 at the first clock edge, subsequently, all IOs will output 0, the 3<sup>rd</sup> will output 1, etc.

The Data Learning Pattern (DLP) resides in an 8-bit Non-Volatile Data Learning Register (NVDLR) as well as an 8-bit Volatile Data Learning Register (VDLR). NVDLR can be programmed with Program NVDLR (PNVDLR 43h) command, and VDLR can be written by Write VDLR (WRVDLR 4Ah) command. Also DLP register value can be read out by Data Learning Pattern Read (RDDLP) command. All 3 commands (PNVDLR, WRVDLR, RDDLP) are valid in SPI mode and QPI mode.

**Figure 8.146 FRDTR Sequence with DLP In SPI Mode (0Dh [EXTADD=0], 3-byte address)**



**Figure 8.147 FRDTR Sequence with DLP In QPI Mode (0Dh [EXTADD=0], 3-byte address)**



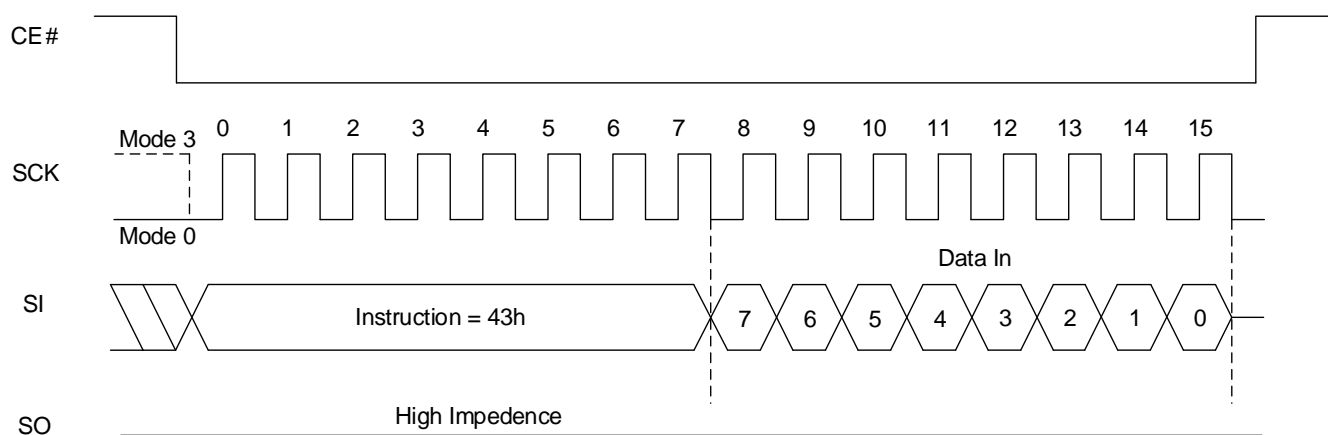
### 8.67 PROGRAM NVDLR OPERATION (PNVDLR, 43h)

Before the Program NVDLR (PNVDLR) command can be accepted by the device, a Write Enable (WREN) command must be issued and decoded by the device. After the Write Enable (WREN) command has been decoded successfully, the device will set the Write Enable Latch (WEL) to enable the PNVDLR operation.

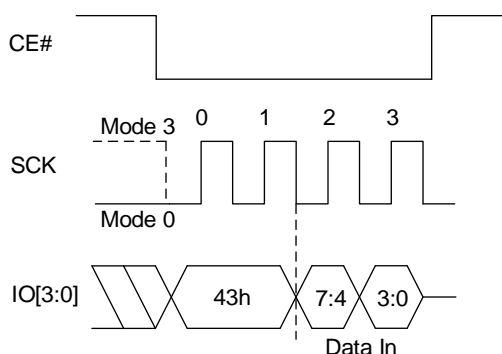
The PNVDLR command (43h) is entered by shifting the instruction and data byte on SI.

CE# must be driven to the logic high state after the eighth (8<sup>th</sup>) bit of data has been latched. If not, the PNVDLR command is not executed. As soon as CE# is driven to the logic high state, the self-timed PNVDLR operation is initiated. While the PNVDLR operation is in progress, the Status Register may be read to check the value of the Write-In Progress (WIP) bit. The Write-In Progress (WIP) bit is a 1 during the PNVDLR cycle, and is 0 when it is completed. The PNVDLR operation can report a program error in the P\_ERR bit of the Extended Read Register. When the PNVDLR operation is completed, the Write Enable Latch (WEL) is set to 0.

**Figure 8.148 PNVDLR Sequence In SPI Mode**



**Figure 8.149 PNVDLR Sequence In QPI Mode**



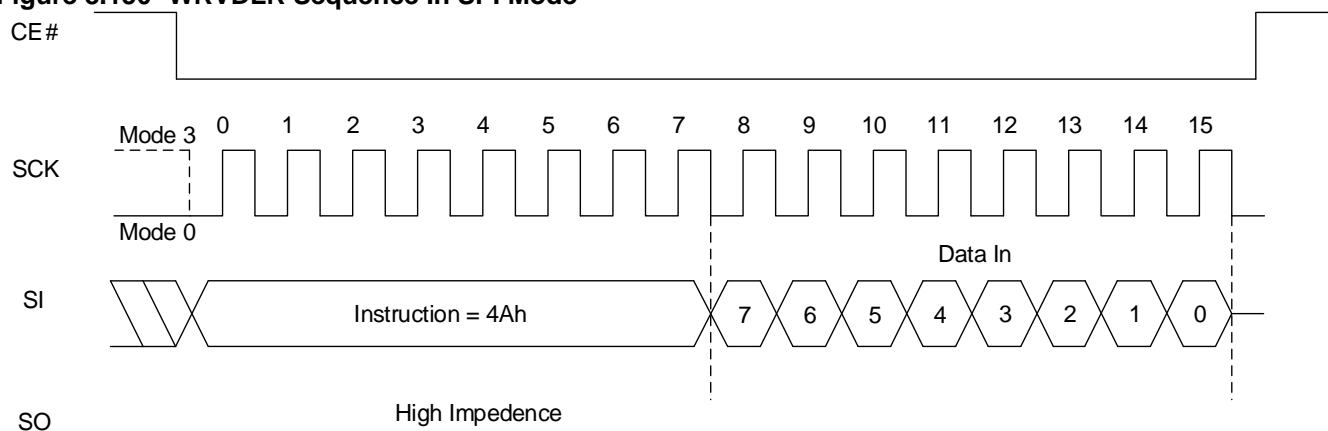
### 8.68 WRITE VDLR OPERATION (WRVDLR, 4Ah)

Before the Write VDLR (WRVDLR) command can be accepted by the device, a Write Enable (WREN) command must be issued and decoded by the device. After the Write Enable (WREN) command has been decoded successfully, the device will set the Write Enable Latch (WEL) to enable the WRVDLR operation.

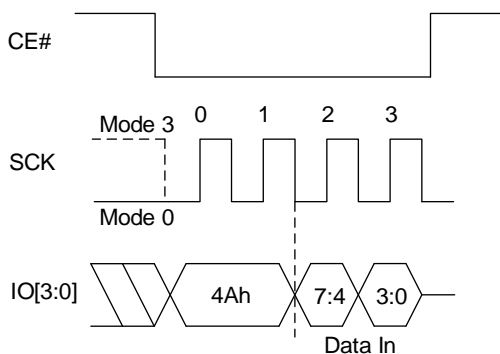
The WRVDLR command (4Ah) is entered by shifting the instruction and data byte on SI.

CE# must be driven to the logic high state after the eighth (8<sup>th</sup>) bit of data has been latched. If not, the WRVDLR command is not executed. As soon as CE# is driven to the logic high state, the WRVDLR operation is initiated with no delays.

**Figure 8.150 WRVDLR Sequence In SPI Mode**



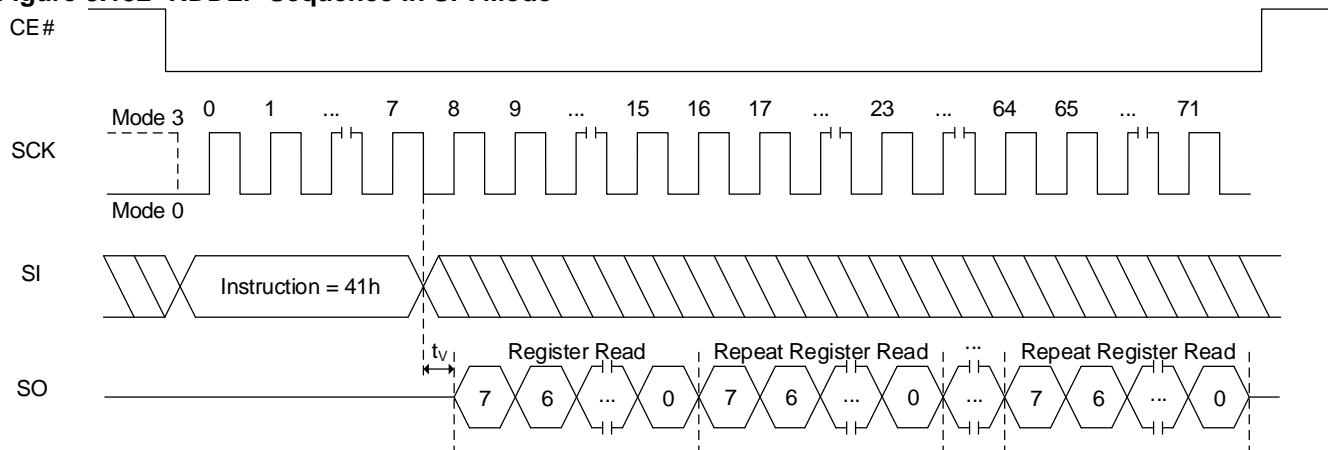
**Figure 8.151 WRVDLR Sequence In QPI Mode**



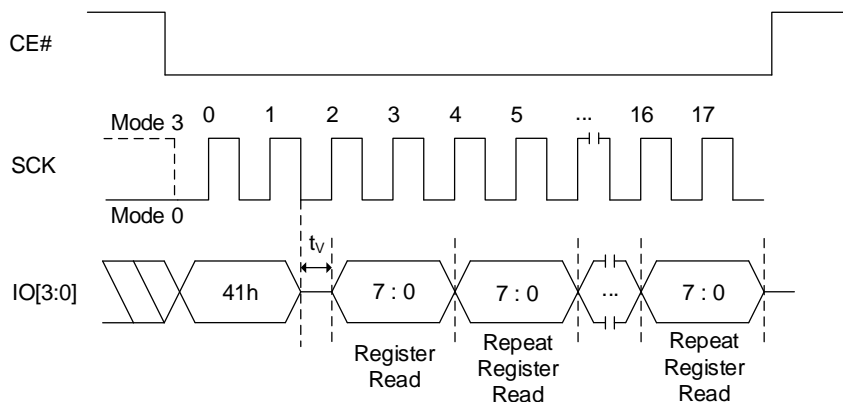
## 8.69 READ DLP OPERATION (RDDLP, 41h)

The instruction is shifted on SI, then the 8-bit DLP from VDLR is shifted out on SO. It is possible to read the DLP continuously by providing multiples of eight clock cycles.

**Figure 8.152 RDDLP Sequence In SPI Mode**



**Figure 8.153 RDDLP Sequence In QPI Mode**





### 8.70 WRITE ROOT KEY OPERATION (9Bh+00h)

Write Root Key operation is to initialize the Root Key Register corresponding to the received Counter Address with the received Root Key.

After command is issued on the interface, the device must ensure that the received transaction is error free. This includes following conditions:

- Payload size is correct (including OP1 is 64 bytes)
- Counter Address falls within the range of supported counters.
- The Root Key Register corresponding to the requested Counter Address was previously uninitialized. [Root\_Key\_Reg\_Init\_State[Counter\_Address]=0xFFh]
- Truncated signature field is the same as least significant 224 bits of HMAC-SHA-256 based signature computed based on received input parameters.

If the received transaction is error free, the device executes the operation successfully, and posts “successful completion” in the RPMC Status Register. This command must be executed to ensure power cycling in the middle of command execution is properly handled. This requires that internal state tracking the root key register initialization is written as the last operation of the command execution.

[Root\_Key\_Reg\_Init\_State[Counter\_Address] = 0]

Root Key Register Write with root key is = 256'hFF...FF is used as a temporary key. When this request is received error free Root\_Key\_Reg\_Init\_State[Counter\_Address] is not affected. Instead, only the corresponding Monotonic Counter is initialized to “0” if previously not initialized.

This state is tracked as separate state using MC\_Init\_State[Counter\_Address]. This state is used to leave the monotonic counters at the current value when an error free Root Key Register operation is received. (Both 256'hFF..FF and non 256'hFF..FF)

Once this command is successfully executed with a non 256'hFF..FF Root Key, the device will not accept the “Write Root Key Register” command any more, and the Root Key value cannot be read out any more.

### **8.71 UPDATE HMAC KEY OPERATION (9Bh+01h)**

Update HMAC Key operation is to update the HMAC Key Register corresponding to the received Counter Address with a new HMAC Key calculated based on received input.

This command must be issued only on every power cycle event on the interface. This allows the HMAC key storage to be implemented using volatile memory. Status registers busy indication is expected to indicate busy for double the amount of Read\_Counter\_Polling\_Delay specified in SFDP table since this command performs two distinct HMAC-SHA-256 computations.

After the command is issued on the interface, the device must ensure that the received transaction is error free. This includes checking following conditions:

- Payload size is correct (including OP1 is 40 bytes)
  - Counter Address falls within the range of supported counters.
  - The Monotonic Counter corresponding to the requested Counter Address was previously initialized.
  - Signature matches the HMAC-SHA-256 based signature computed based on received input parameters.
- This command performs two HMAC-SHA-256 operations.

If the received transaction is error free, the device executes the operation successfully, and posts “successful completion” in the RPMC Status Register.

If the received transaction has errors, the device does not execute the operation, and posts the corresponding error in the RPMC Status Register.

## **8.72 INCREMENT MONOTONIC COUNTER OPERATION (9Bh+02h)**

This command is used to increment Monotonic counter by 1 inside the device.

After the command is issued on the interface, the device must ensure that the received transaction is error free. This includes checking following conditions:

- Payload size is correct (including OP1 is 40 bytes)
- Counter Address falls within the range of supported counters.
- The Monotonic Counter corresponding to the requested Counter Address was previously initialized.
- The HMAC Key Register corresponding to the requested Counter Address was previously initialized.
- The requested Signature matches the HMAC-SHA-256 based signature computed based on received input parameters.
- The received Counter Data matches the current value of the counter read from the device.

If the received transaction is error free, the device executes the operation successfully, and posts “successful completion” in the RPMC Status Register.

If the received transaction has errors, the device does not execute the operation, and posts the corresponding error in the RPMC Status Register. The increment counter implementation should make sure that the counter increment operation is performed in a power glitch aware manner.

### **8.73 REQUEST MONOTONIC COUNTER OPERATION (9Bh+03h)**

This command is used to request the Monotonic counter value inside the device.

After the command is issued on the interface, the device must ensure that the received transaction is error free. This includes checking following conditions:

- Payload size is correct (including OP1 is 48 bytes)
- Counter Address falls within the range of supported counters.
- The Monotonic Counter corresponding to the requested Counter Address was previously initialized.
- The HMAC Key Register corresponding to the requested Counter Address was previously initialized.
- The requested Signnnature matches the HMAC-SHA-256 based signature computed based on received input parameters.

If the received transaction is error free, the device executes the operation successfully, and posts “successful completion” in the RPMC Status Register.

If the received transaction has errors, the device does not execute the operation, and posts “successful completion” in the RPMC Status Register. In response to this command, the device reads the monotonic counter addressed by counter address. It calculates HMAC-SHA-256 signatures the second time, based on the following parameters:

- HMAC Message [127:0] = Tag [95], Counter\_Data\_Read [31:0]
- HMAC Key [255:0] = HMAC\_Key\_Register [Counter\_Address][255:0]

### **8.74 RESERVED RPMC OPERATION (9Bh+04h~FFh)**

If the host controller issues any of the reserved command-types, the device must return Error status in the RPMC Status Register. It asserts bit 2 to indicate that a reserved command-type was issued:

### 8.75 READ RPMC STATUS OPERATION (96h)

This command is used to read from RPMC Status Register for the RPMC status of any previously issued OP1 command.

When previous command was 9Bh+00h, 9Bh+01h or 9Bh+02h, the device will output the RPMC\_Status [7:0] value continuously until CE# terminates the operation.

When previous command was Request Monotonic Counter and if the device returns “successful completion (BUSY=0 in the RPMC Status Register), then it must also return valid data (Tag, Counter, and Signature) with RPMC Status.

If there are other error flags, the values returned in Tag, Counter, and Signature field are invalid. The host controller may abort the read prematurely prior to reading the entire payload. This may occur when the controller wants simply read the RPMC status or when it observes an error being returned in the RPMC status field. The host may also continue reading past the defined payload size of 49 bytes. Since this is an error condition, the device may return any data past the defined payload size. The host controller must ignore the data.

When BUSY=1. From Byte 3 and beyond, the device will output the RPMC\_Status [7:0] value continuously until CE# terminates the operation. The device will not output Tag, CounterData & Signature fields when BUSY=1. Once BUSY becomes 0, another OP2 command must be issued to readout the correct Tag, Data & Signature fields.

- 96h after 9Bh+00h, 9Bh+01h or 9Bh+02h: Read RPMC\_Status [7:0] repeatedly until CE# HIGH.
- 96h after 9Bh+03h
  - BUSY=0: Read RPMC Status, Data and Signature
  - BUSY=1: Read RPMC\_Status [7:0] repeatedly until CE# HIGH.

When BUSY = 1 to 0, then another 96h: Read RPMC Status, Data and Signature

**Table 8.12 RPMC Status Register Definition**

RPMC Status [7:0]	Applicable Command Type (OP1 command)	Description
0000 0000	-	Power On State (OP2 issued directly after power-up)
1000 0000	00, 01, 02, 03	Successful completion (no errors) of OP1 command.
0XXX XXX1	00, 01, 02, 03, 04-0FF	Bit [0] is set to “1” when device is busy executing OP1 command, and reset to “0” when OP1 command execution is completed.
0XXX XX1X	00, 01	Bit [1] is set to “1” when the correct payload size is received. When command type is 00, this error bit is set to “1” on Root Key Register Overwrite or Counter Address out of range or Truncated Signature mismatch error. When command type is 01, this error bit is set to “1” when the corresponding monotonic counter is uninitialized.
0XXX X1XX	00, 01, 02, 03	Bit [2] is set to “1” on Signature Mismatch, Counter Address out of range when correct payload size is received; or Command type is out of range; or incorrect payload size is received.
0XXX X1XX	00, 01, 02, 03	Bit [2] is set to “1” on Signature Mismatch, Counter Address out of range when correct payload size is received; or Command type is out of range; or incorrect payload size is received.
0XXX 1XXX	02, 03	Bit [3] is set to “1” on HMAC Key Register (or monotonic counter) uninitialized and correct payload size is received.
0XX1 XXXX	02	Bit [4] is set to “1” on Counter Data Mismatch and correct payload size is received.
0X1X XXXX	-	Bit [5] is set to “1” on Fatal Error, e.g. program fail, no valid counter found after initialization.
Current Value	-	RPMC Status register will naturally not be updated until first 8 bits of OP1 is received. However it is expected that if an error is found or bit [0] needs to be set, it is set within 8 clocks after OP1 is received.

## 9. ELECTRICAL CHARACTERISTICS

### 9.1 ABSOLUTE MAXIMUM RATINGS <sup>(1)</sup>

Storage Temperature		-65°C to +150°C
Input Voltage with Respect to Ground on All Pins		-0.5V to $V_{CC} + 0.5V$
All Output Voltage with Respect to Ground		-0.5V to $V_{CC} + 0.5V$
$V_{CC}$	IS25RLP	-0.5V to +6.0V
	IS25RWP	-0.5V to +2.5V

**Notes:**

1. Applied conditions greater than those listed in “Absolute Maximum Ratings” may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions above those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect reliability.
2. ANSI/ESDA/JEDEC JS-001

### 9.1 OPERATING RANGE

Operating Temperature	Industrial Grade I	-40°C to 85°C
$V_{CC}$ Power Supply	IS25RLP	2.3V (V <sub>MIN</sub> ) – 3.6V (V <sub>MAX</sub> ); 3.0V (Typ)
	IS25RWP	1.70V (V <sub>MIN</sub> ) – 1.95V (V <sub>MAX</sub> ); 1.8V (Typ)

## 9.2 DC CHARACTERISTICS

(Under operating range)

Symbol	Parameter	Condition	Min	Typ <sup>(2)</sup>	Max	Units
I <sub>CC1</sub>	V <sub>CC</sub> Active Read current <sup>(3)</sup>	NORD at 80MHz		8	10	mA
		FRD Single at 166MHz		10	12	
		FRD Quad at 166MHz		15	18	
		FRD Single at 133MHz		9	11	
		FRD Quad at 133MHz		13	16	
		FRD Quad at 83MHz		11	12	
		FRD Quad at 104MHz		12	14	
		FRD Single DTR at 80MHz		9	12	
		FRD Quad DTR at 80MHz		13	16	
I <sub>CC2</sub>	V <sub>CC</sub> Program Current	CE# = V <sub>CC</sub>		25	30	mA
I <sub>CC3</sub>	V <sub>CC</sub> WRSR Current	CE# = V <sub>CC</sub>		25	30	
I <sub>CC4</sub>	V <sub>CC</sub> Erase Current (SER/4SER/BER32/4BER32/ BER64/4BER64)	CE# = V <sub>CC</sub>		25	30	
I <sub>CC5</sub>	V <sub>CC</sub> Erase Current (CE)	CE# = V <sub>CC</sub>		25	30	
I <sub>CC6</sub>	RPMC OP1 Write Only	CE# = V <sub>CC</sub>		10	20	
I <sub>CC7</sub>	RPMC OP1 & Array Read	CE# = V <sub>CC</sub>		25	40	
I <sub>CC8</sub>	RPMC OP1 & Array PROGRAM/ERASE	CE# = V <sub>CC</sub>		35	50	
I <sub>CC9</sub>	RPMC OP2 Read Only	CE# = V <sub>CC</sub>		5	10	
I <sub>CC10</sub>	RPMC OP2 & Array PROGRAM/ERASE	CE# = V <sub>CC</sub>		30	40	
I <sub>SB1</sub>	V <sub>CC</sub> Standby Current CMOS	CE# = V <sub>CC</sub> , V <sub>IN</sub> = GND or V <sub>CC</sub> <sup>(4)</sup>		35	100	μA
I <sub>SB2</sub>	Deep power down current	IS25RLP		13	40	μA
		IS25RWP		6	40	
I <sub>LI</sub>	Input Leakage Current	V <sub>IN</sub> = 0V to V <sub>CC</sub>			±4	μA
I <sub>LO</sub>	Output Leakage Current	V <sub>IN</sub> = 0V to V <sub>CC</sub>			±4	μA
V <sub>IL</sub> <sup>(1)</sup>	Input Low Voltage		-0.5		0.3V <sub>CC</sub>	V
V <sub>IH</sub> <sup>(1)</sup>	Input High Voltage		0.7V <sub>CC</sub>		V <sub>CC</sub> + 0.3	V
V <sub>OL</sub>	Output Low Voltage	I <sub>OL</sub> = 100 μA			0.2	V
V <sub>OH</sub>	Output High Voltage	I <sub>OH</sub> = -100 μA	V <sub>CC</sub> - 0.2			V

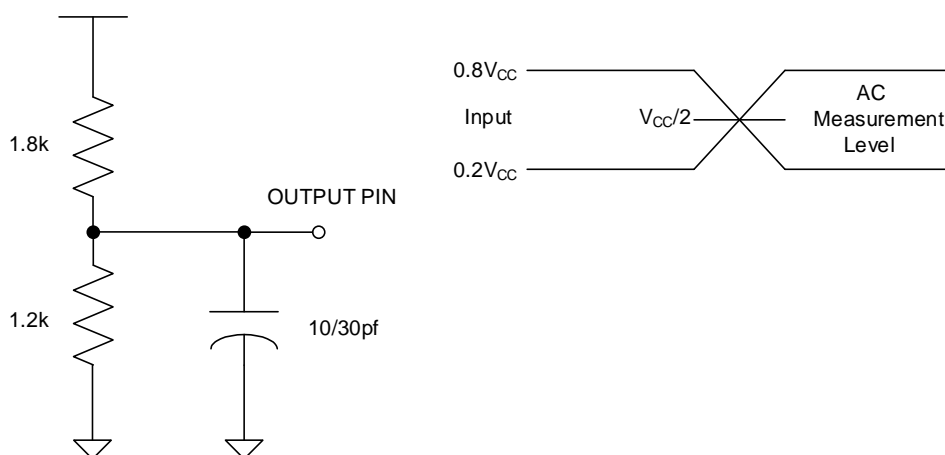
### Notes:

- Maximum DC voltage on input or I/O pins is V<sub>CC</sub> + 0.5V. During voltage transitions, input or I/O pins may overshoot V<sub>CC</sub> by +2.0V for a period of time not to exceed 20ns. Minimum DC voltage on input or I/O pins is -0.5V. During voltage transitions, input or I/O pins may undershoot GND by -2.0V for a period of time not to exceed 20ns.
- Typical values are included for reference only and are not guaranteed or tested. Typical values are measured at V<sub>CC</sub> = V<sub>CC</sub> (Typ), TA=25°C.
- Outputs are unconnected during reading data so that output switching current is not included.
- V<sub>IN</sub> = V<sub>CC</sub> for the dedicated RESET# pin.

### 9.3 AC MEASUREMENT CONDITIONS

Symbol	Parameter	Min	Max	Units
CL	Load Capacitance		30	pF
			10	pF
TR,TF	Input Rise and Fall Times		5	ns
VIN	Input Pulse Voltages	0.2V <sub>CC</sub> to 0.8V <sub>CC</sub>		V
VREFI	Input Timing Reference Voltages	0.3V <sub>CC</sub> to 0.7V <sub>CC</sub>		V
VREFO	Output Timing Reference Voltages	0.5V <sub>CC</sub>		V

**Figure 9.1 Output test load & AC measurement I/O Waveform**



### 9.4 PIN CAPACITANCE

(TA = 25°C, VCC=3V for IS25RLP, VCC=1.8V for IS25RWP, 1MHz)

Symbol	Parameter	Test Condition	IS25RLP		IS25RWP		Units
			Min	Max	Min	Max	
C <sub>IN</sub>	Input Capacitance (CE#, SCK)	V <sub>IN</sub> = 0V	-	12	-	12	pF
C <sub>IN/OUT</sub>	Input/Output Capacitance (other pins)	V <sub>IN/OUT</sub> = 0V	-	14	-	18	pF



## 9.5 AC CHARACTERISTICS

(Under operating range,)

Symbol	Parameter				Min	Typ <sup>(2)</sup>	Max	Units
f <sub>CT</sub>	Clock Frequency except for fast read DTR and read (03h)	IS25RLP	VCC = 2.7V to 3.6V		0		166	MHz
			VCC = 2.3V to 3.6V		0		133	
		IS25RWP	VCC = 1.7V to 1.95V		0		166 <sup>(4, 5)</sup>	
	Clock Frequency for fast read DTR: SPI DTR, Dual DTR, Dual I/O DTR, Quad I/O DTR, and QPI DTR.				0		80	
	Clock Frequency for RPMC (OP1, OP2)				0		104	
f <sub>C</sub>	Clock Frequency for read (03h)	IS25RLP		0		80		
		IS25RWP		0		50		
t <sub>CLCH</sub> <sup>(1)</sup>	SCK Rise Time (peak to peak)				0.1			V/ns
t <sub>CHCL</sub> <sup>(1)</sup>	SCK Fall Time ( peak to peak)				0.1			V/ns
t <sub>CKH</sub>	SCK High Time		For read (03h)		0.45 x 1/f <sub>Cmax</sub>			ns
			For others		0.45 x 1/f <sub>CTmax</sub>			
t <sub>CKL</sub>	SCK Low Time		For read (03h)		0.45 x 1/f <sub>Cmax</sub>			ns
			For others		0.45 x 1/f <sub>CTmax</sub>			
t <sub>CEH</sub>	CE# High Time				7			ns
t <sub>CS</sub>	CE# Setup Time				5			ns
t <sub>CH</sub>	CE# Hold Time				5			ns
t <sub>CHSL</sub>	CE# Not Active Hold Time				5			ns
t <sub>SHCH</sub>	CE# Not Active Setup Time				5			ns
t <sub>DS</sub>	Data In Setup Time		STR		2			ns
			DTR		1.5			
t <sub>DH</sub>	Data in Hold Time		STR		2			ns
			DTR		1.5			
			RPMC		4.0			
t <sub>V</sub>	Output Valid Time	IS25RLP	2.7~3.6V	@10pF			5.5 <sup>(4)</sup>	ns
				@30pF			7.0 <sup>(4)</sup>	
			2.3~3.6V	@10pF			6.0	
				@30pF			7.5	
		IS25RWP	1.70~1.95V	@10pF			5.5	
				@30pF			7.0	
t <sub>V1</sub>	Output Valid Time of OP2 operation			@10pF			7.0	ns
				@30pF			8.0	
t <sub>OH</sub>	Output Hold Time				2			ns
t <sub>DIS</sub> <sup>(1)</sup>	Output Disable Time						8	ns
t <sub>WHS</sub> <sup>(3)</sup>	Write Protect Setup Time				20			ns
t <sub>SHWL</sub> <sup>(3)</sup>	Write Protect Hold Time				100			ns
t <sub>HLCH</sub>	HOLD Active Setup Time relative to SCK				2			ns
t <sub>CHHH</sub>	HOLD Active Hold Time relative to SCK				2			ns

Symbol	Parameter	Min	Typ <sup>(2)</sup>	Max	Units
t <sub>HCH</sub>	HOLD Not Active Setup Time relative to SCK	2			ns
t <sub>CHHL</sub>	HOLD Not Active Hold Time relative to SCK	2			ns
t <sub>LZ</sub> <sup>(1)</sup>	HOLD to Output Low Z			8	ns
t <sub>HZ</sub> <sup>(1)</sup>	HOLD to Output High Z			8	ns
t <sub>EC</sub>	Sector Erase Time (4Kbyte)		100	300	ms
	Block Erase Time (32Kbyte)		0.14	0.5	s
	Block Erase time (64Kbyte)		0.17	1.0	s
	Block Erase time (256Kbyte)		0.68	4.0	s
	Chip Erase Time		70	180	s
t <sub>PP</sub>	Page Program Time	256byte	0.3	1.0	ms
		512byte	0.6	2.0	
t <sub>RES1</sub> <sup>(1)</sup>	Release deep power down		100	200	μs
t <sub>DP</sub> <sup>(1)</sup>	Deep power down		16	32	μs
t <sub>W</sub>	Write Status Register time		2	15	ms
t <sub>SUS</sub> <sup>(1)</sup>	Suspend to read ready			100	μs
t <sub>SRST</sub> <sup>(1)</sup>	Software Reset recovery time			35	μs
t <sub>RESET</sub> <sup>(1)</sup>	RESET# pin low pulse width	150			ns
t <sub>HWRST</sub> <sup>(1)</sup>	Hardware Reset recovery time			800	μs
t <sub>PDRDY_RPMC</sub>	Power ready RPMC access time			2	ms
t <sub>KEY</sub>	RPMC Write Root Key		500	550	μs
t <sub>HMAC</sub>	RPMC Update HMAC Key Register		100	445	μs
t <sub>INC1</sub>	RPMC Increment Monotonic Counter			300	ms
t <sub>REQ</sub>	RPMC Request Monotonic Counter		80	105	μs

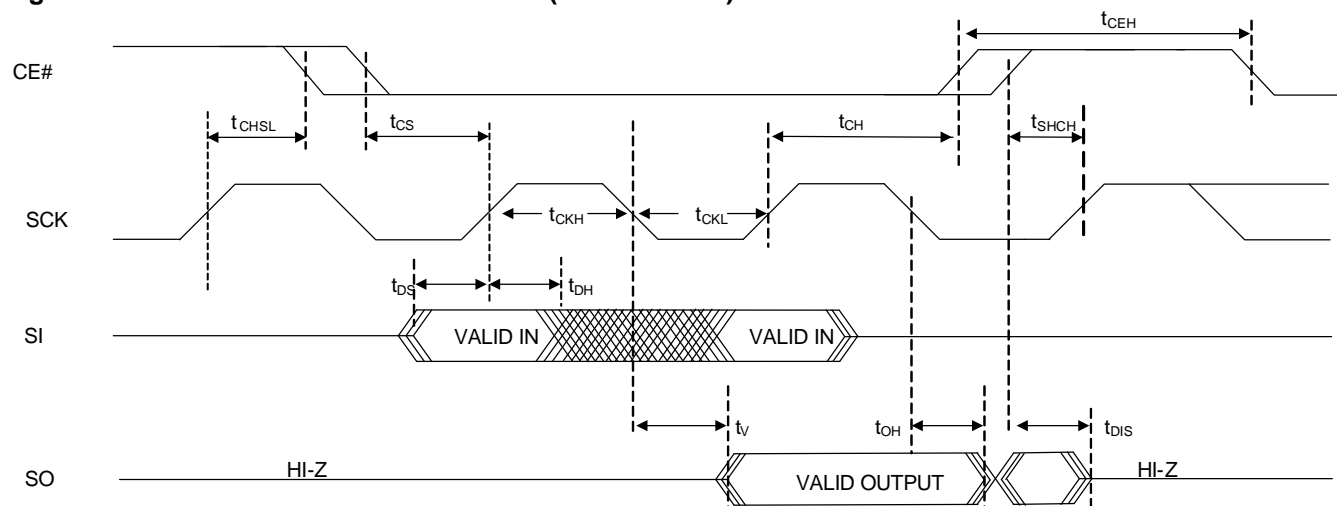
**Notes:**

1. These parameters are characterized and not 100% tested.
2. Typical values are included for reference only and are not guaranteed or tested. Typical values are measured at V<sub>CC</sub> = V<sub>CC</sub> (Typ), TA=25°C.
3. Only applicable as a constraint for a WRITE STATUS REGISTER command when SRWD is set at 1.
4. Values are guaranteed by characterization and not 100% tested in production.
5. RDSR(05h), RDID(ABh), RDJDID(9Fh), RDJDIDQ(AFh, QPI mode only), RDRP(61h), RDFR(48h) operation frequency for 1.8V device is like below:

SPI Mode	QPI Mode
166MHz	133MHz

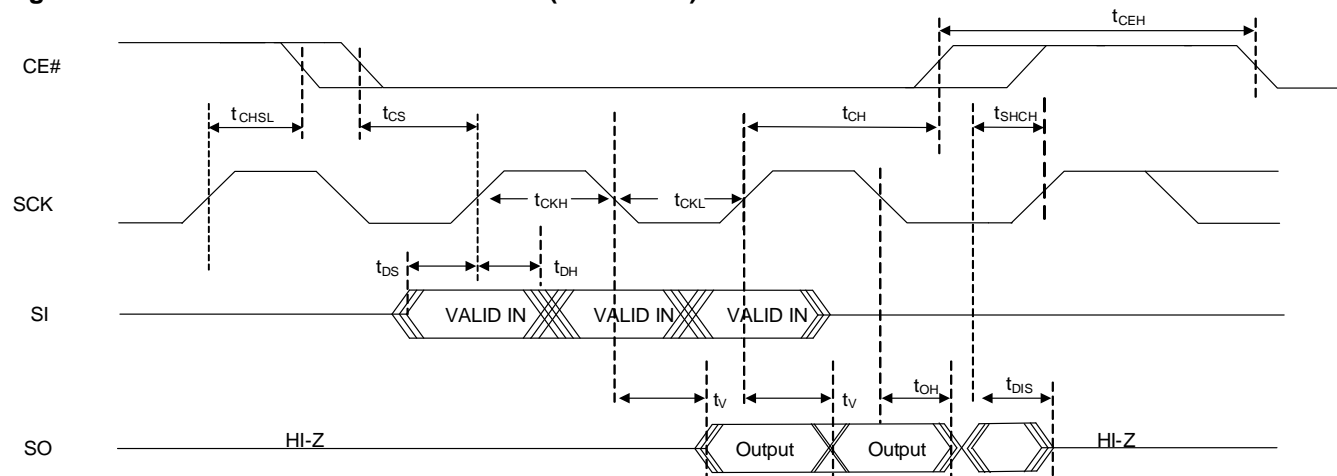
## 9.6 SERIAL INPUT/OUTPUT TIMING

**Figure 9.2 SERIAL INPUT/OUTPUT TIMING (Normal Mode) <sup>(1)</sup>**



**Note1: For SPI Mode 0 (0,0)**

**Figure 9.3 SERIAL INPUT/OUTPUT TIMING (DTR Mode) <sup>(1)</sup>**



**Note1: For SPI Mode 0 (0,0)**

Figure 9.4 HOLD TIMING

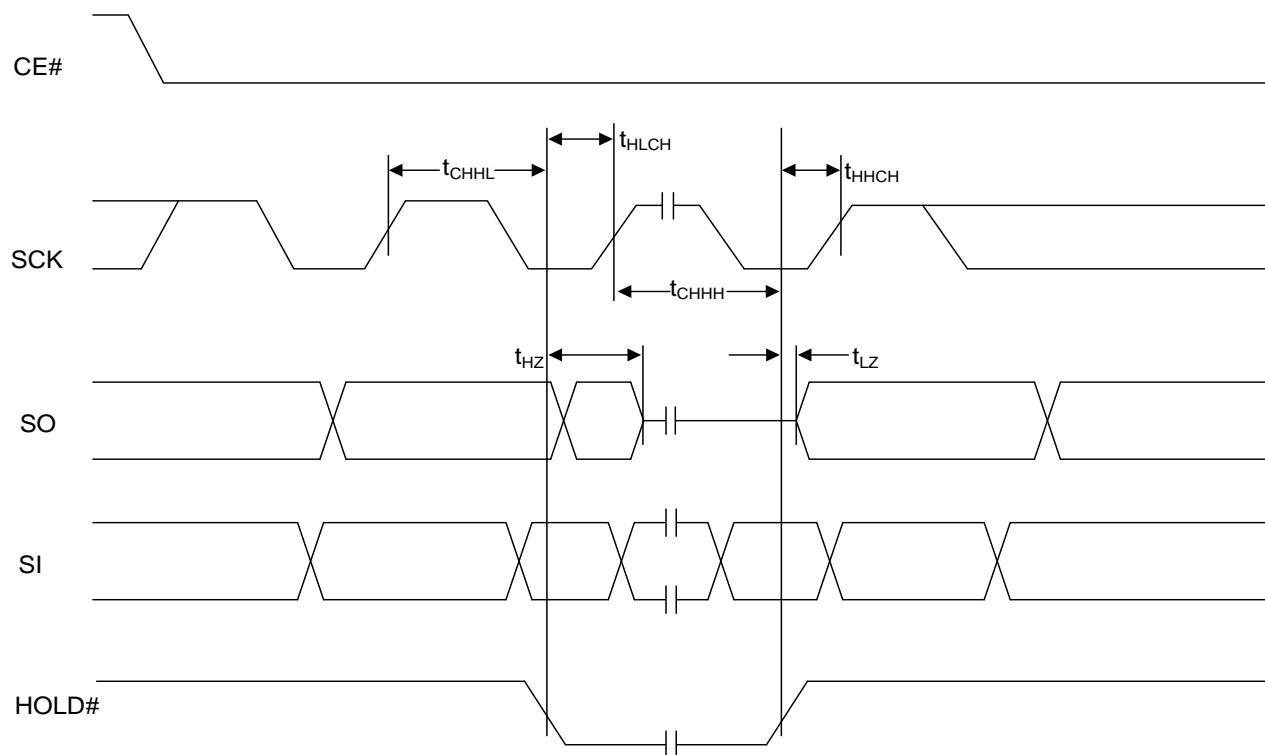
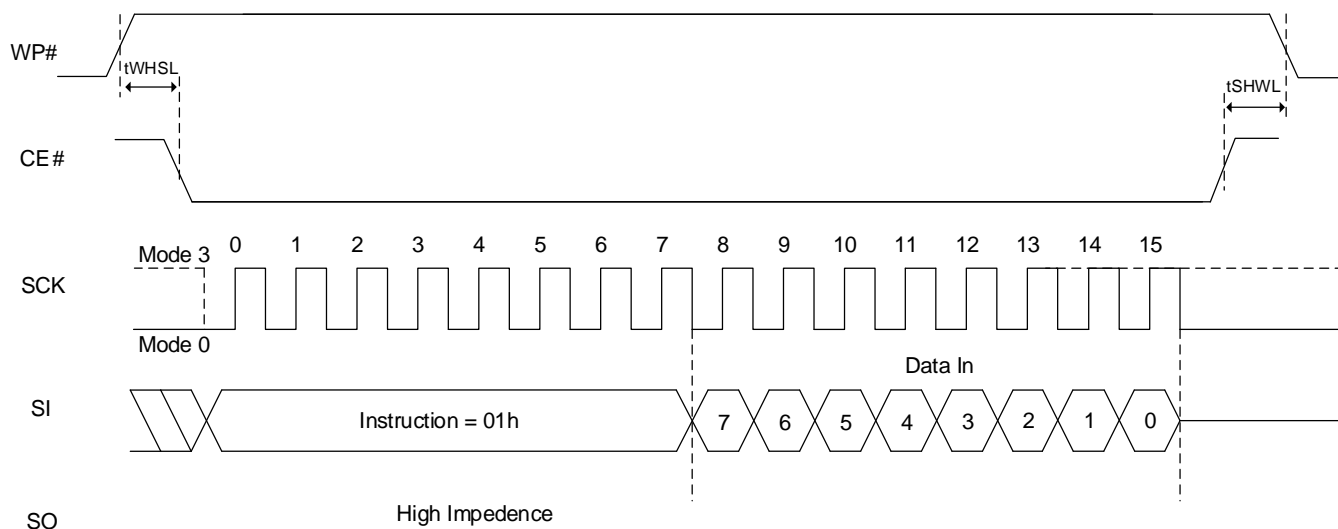


Figure 9.5 WRITE PROTECT SETUP AND HOLD TIMING DURING WRITE STATUS REGISTER (SRWD=1)

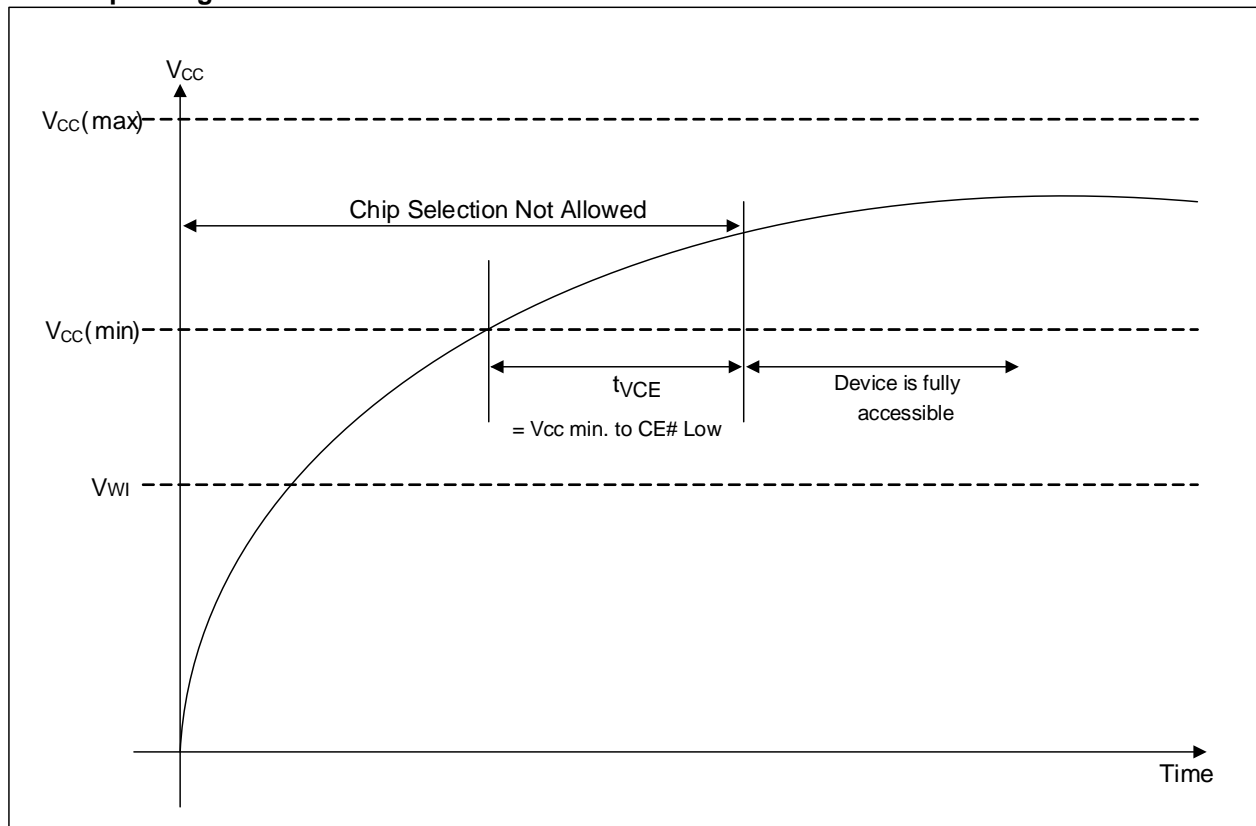


Note: WP# must be kept high until the embedded operation finish.

## 9.7 POWER-UP AND POWER-DOWN

At Power-up and Power-down, the device must be NOT SELECTED until  $V_{CC}$  reaches at the right level. (Adding a simple pull-up resistor on CE# is recommended.)

### Power up timing



Symbol	Parameter		Min.	Max	Unit
$t_{VCE}^{(1)}$	$V_{CC(min)}$ to CE# Low		300		us
$V_{WI}^{(1)}$	Write Inhibit Voltage	IS25RLP		2.1	V
		IS25RWP		1.5	

**Note:** These parameters are characterized and not 100% tested.

## 9.8 PROGRAM/ERASE PERFORMANCE

Parameter		Typ	Max	Unit
Sector Erase Time (4Kbyte)		100	300	ms
Block Erase Time (32Kbyte)		0.14	0.5	s
Block Erase Time (64Kbyte)		0.17	1.0	s
Block Erase Time (256Kbyte)		0.68	4.0	s
Chip Erase Time		70	180	s
Page Programming Time	256byte	0.3	1.0	ms
	512byte	0.6	2.0	
Byte Program		10	50	μs

**Note:** These parameters are characterized and not 100% tested.

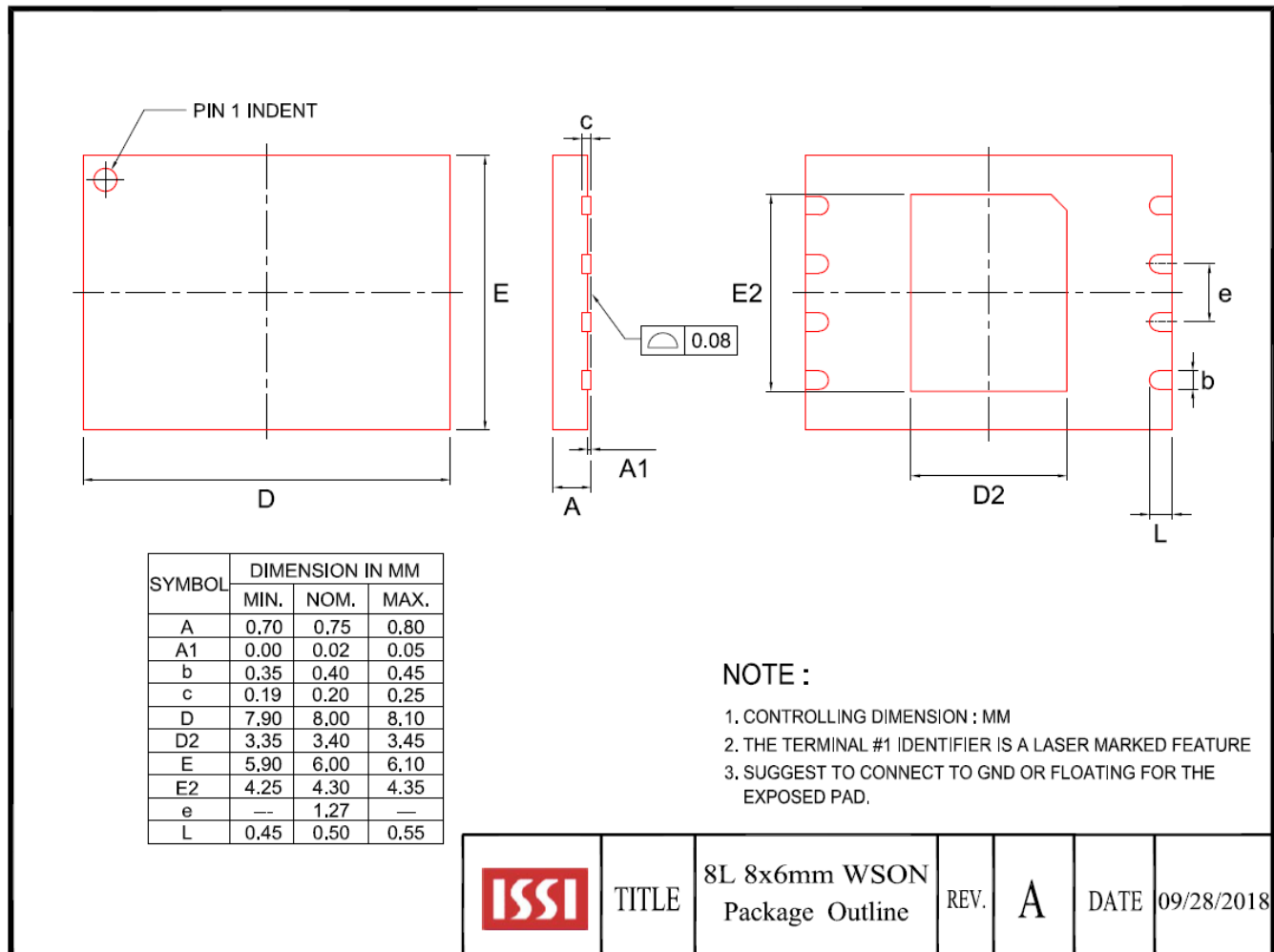
## 9.9 RELIABILITY CHARACTERISTICS

Parameter	Min	Max	Unit	Test Method
Endurance	100,000	-	Cycles	JEDEC Standard A117
Data Retention	20	-	Years	JEDEC Standard A117
Latch-Up	-100	+100	mA	JEDEC Standard 78

**Note:** These parameters are characterized and not 100% tested.

## 10. PACKAGE TYPE INFORMATION

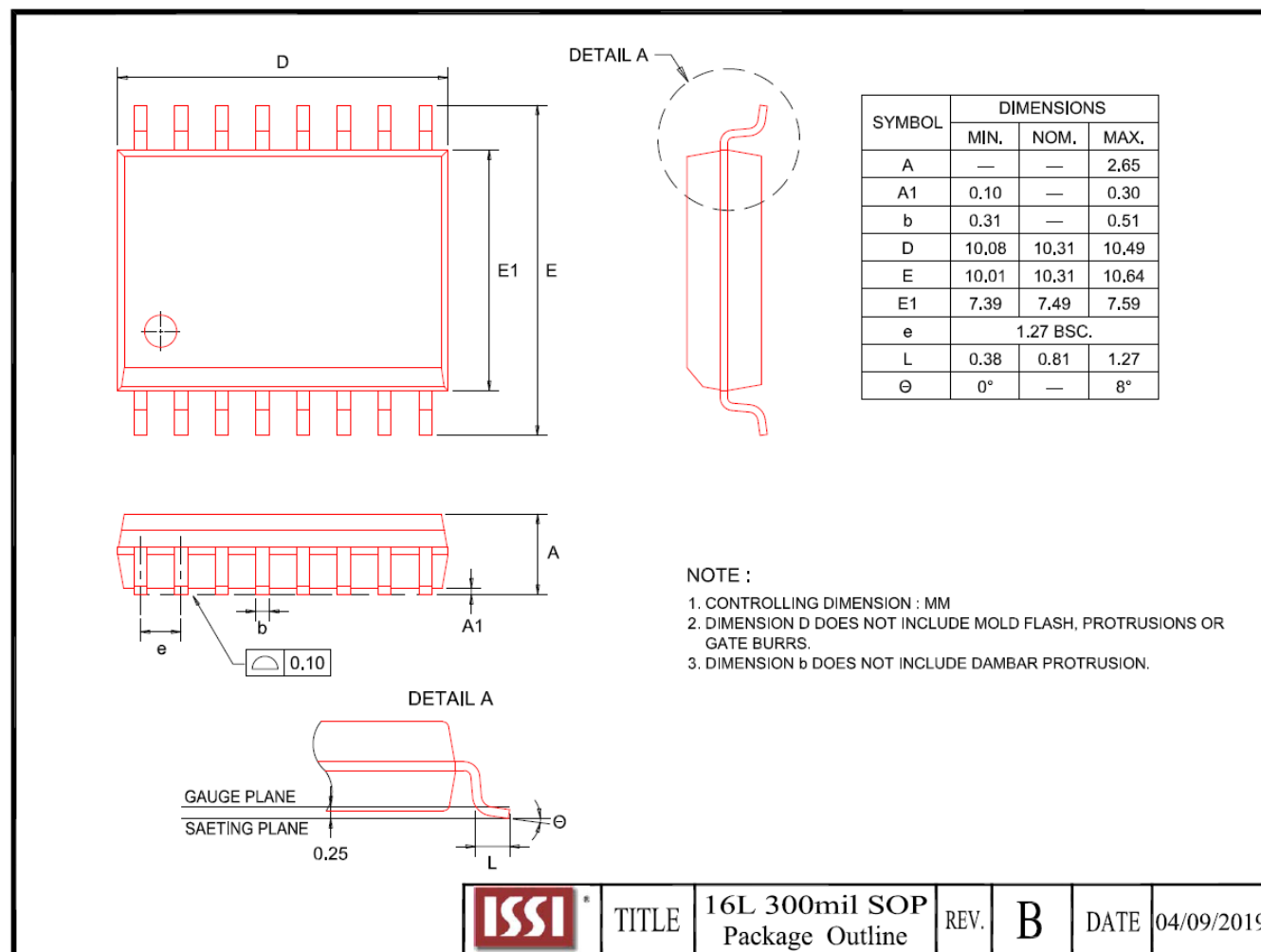
### 10.1 8-CONTACT ULTRA-THIN SMALL OUTLINE NO-LEAD (WSON) PACKAGE 8X6MM (J)



**Notes:**

1. Please [click here](#) to refer to Application Note (AN25D011, Thin USON/WSON/XSON package handling precautions) for assembly guidelines.
2. Pad Open size is 3.4mm x 4.3mm

**10.2 16-LEAD PLASTIC SMALL OUTLINE PACKAGE (300 MILS BODY WIDTH) (M)**





## 11. ORDERING INFORMATION – Valid Part Numbers

IS25RLP 256 E - P M L I

### TEMPERATURE RANGE

I = Industrial (-40°C to +85°C)

### PACKAGING CONTENT

L = RoHS compliant

### PACKAGE Type

J = 8-contact WSON (8x6mm)

M = 16-pin SOIC 300mil (Call Factory)

### Option

J = Standard (Call Factory)

Q = QE bit set to 1

P = QE bit set to 1 + Dedicated RESET# (Call Factory)

### Die Revision

E = Revision E

### Density

256 = 256 Megabit

### BASE PART NUMBER

**IS = Integrated Silicon Solution Inc.**

25RLP = FLASH, 2.30V ~ 3.60V, QPI & RPMC

25RWP = FLASH, 1.70V ~ 1.95V, QPI & RPMC (Call Factory)

#### Note:

1. Call Factory for other package options available

Density, Voltage	Frequency (MHz)	Order Part Number	Package
256M, 3V	STR 166, DTR 80 RPMC 104	IS25RLP256E-QJLI	8-contact WSON 8x6mm